Man's major projects have been and are realized through organization. This implies not only a hierarchy of responsibility but also a means of communicating within it. This is particularly true in a modern military organization.

In this day of missiles and supersonic aircraft, real-time information is essential for rapid, effective response. Our large, complex, globally dispersed defense forces generate enough information to overwhelm unaided men; hence computers, large memories, and sophisticated displays are needed to automate analysis and to aid decision making. A vast communications network utilizing almost every part of the spectrum is needed to tie the national command centers to land combat forces, ships and aircraft. Satellites have become a significant part of this link.

This issue highlights modern concepts in hardware, software and systems for Command, Control, and Communications for our defense forces.

Dr. H. J. Woll
Division Vice President
Government Engineering
Government and Commercial Systems
Moorestown, NJ

Our cover
Emphasizes command, control, and communications (C3) systems. The front cover symbolizes the remotely piloted vehicle concept (see Shore, p. 13); the back cover shows communications satellites and tactical ground forces. In all C3 systems, the human is central to the decision-making processes (see Ireland, p. 8). Cover concept and design: Ed Burke, Missile and Surface Radar Division, Moorestown, NJ.
### Contents

#### Editorial input
- Editorial input for authors only

#### Engineer and the Corporation
- Command, control, and communications—an overview
  - D. Shore

#### Command, Control, Communications

<table>
<thead>
<tr>
<th>Topic</th>
<th>Authors</th>
</tr>
</thead>
<tbody>
<tr>
<td>Human element in command and control systems</td>
<td>Dr. F. H. Ireland</td>
</tr>
<tr>
<td>Remotely piloted vehicles—command and control</td>
<td>D. Shore</td>
</tr>
<tr>
<td>Managing computer program development</td>
<td>Dr. S. A. Steele, R. A. Dupell, A. M. Fleishman, E. T. Hatcher</td>
</tr>
<tr>
<td>Computer program reliability</td>
<td>Dr. P. G. Anderson, L. H. Crandon</td>
</tr>
<tr>
<td>Impact of hardware/software tradeoffs in command and control systems</td>
<td>J. C. Kulp</td>
</tr>
<tr>
<td>Computer selection for command and control systems</td>
<td>H. J. Hurtado, Dr. S. A. Steele</td>
</tr>
<tr>
<td>Software verification and validation for command and control systems</td>
<td>A. R. Chandler</td>
</tr>
<tr>
<td>Simulation of a command and control data processing and display system</td>
<td>C. S. Warren</td>
</tr>
<tr>
<td>AEGIS engineering model command and control system</td>
<td>F. Bernstein, J. Strip</td>
</tr>
<tr>
<td>Command and data handling subsystem for Atmosphere Explorer</td>
<td>W. V. Fuldner</td>
</tr>
<tr>
<td>Performance monitoring and fault diagnostics of command and control systems</td>
<td>T. Taylor, M. LeVarn, J. O'Connell</td>
</tr>
</tbody>
</table>

#### General Interest

<table>
<thead>
<tr>
<th>Topic</th>
<th>Authors</th>
</tr>
</thead>
<tbody>
<tr>
<td>The use of computer-aided design in the manufacturing of welded-wire circuit boards</td>
<td>M. A. Eastwood, A. S. Baran</td>
</tr>
<tr>
<td>Self-calibration of land-based pulsed radar from satellite track</td>
<td>J. J. O'Connor, R. R. Rowe</td>
</tr>
<tr>
<td>Videovoice—a review</td>
<td>S. N. Friedman</td>
</tr>
<tr>
<td>Simplifying microcomputer architecture</td>
<td>J. Weisbecker</td>
</tr>
<tr>
<td>Recent developments in photomultipliers for liquid scintillation counting</td>
<td>D. E. Persyk, T. T. Lewis</td>
</tr>
</tbody>
</table>

#### Departments
- Pen and Podium
- Dates and Deadlines
- Patents Granted
- News and Highlights

#### Special
- IEEE Fellows
If you think the title of this "editorial input" excludes you, you're wrong.

Each year, RCA engineers and scientists put their names to about 2500 formal documents, including published papers and presentations, reports, and patent disclosures. This amounts to more than one document each year for every two engineers. Add in the plethora of "everyday" writing—memos, letters, proposal and report inputs, engineering notebooks—and the number could easily double.

The statistic is at once encouraging and disturbing. On the one hand, 2500 formal documents is an impressive output. Yet, one has only to glance at several RCA Engineer listings of published papers, presentations, and patents granted to see that a rather small percentage of the engineers and scientists produce most of these published documents.

But this imbalance makes sense: for some members of the RCA technical staff, the technical report or paper is the product. In product design and development areas, however, primary emphasis is, rightfully, on the hardware. Before the first unit is designed, built, and shipped, the next five are needed, and there is little time to document any of the solid, often ingenious, engineering work that is the backbone of every successful product line.

There is no easy solution to the problem, except to ask busy engineers to become busier and to offer some time-saving techniques.

Unfortunately, most formal courses in technical writing (and most books about the subject) tend to over-emphasize grammar, style, and elegance of expression, almost to the exclusion of such important considerations as planning, organization, and content. Correct grammar and clear style are important, but they often become unnecessary psychological obstacles to getting a paper written.

There are many places an engineer can find help with these details after his first draft is complete: Senior engineers or managers who have done it before; engineering editors or writers who do it every day; or coworkers who may simply be able to read a paper with a fresh outlook.

The really important part of the writing process—producing a complete first draft—can be done only by the engineer. And there are ways to do it quickly.

Many writers waste time laboring over the introduction, title, section headings, and conclusions. Usually it is better to work from the "inside out"—to write the detailed text first; the words for the title and introduction will spring from this effort.

Another time-saver is to make use of the vast quantities of material hidden in drawers and files already in the form of status reports, memos, and notebook entries. Still another source may be proposals, contract reports, or instruction manuals.

Yet another aid in getting a paper done is to discuss the topic with coworkers; often such two-way communication will help the author organize his effort.

The primary ingredient, however, is self-confidence. The engineer is the final authority; few other professionals are trained to think as logically or organize as well as engineers. They are—by training and tempera-
Command, control, and communications — an introduction
D. Shore

Previous issues of the RCA Engineer have treated the technology associated with military command, control, and communications — or as it is normally abbreviated, C3. We have felt that these articles covering data processing, software, displays and communications needed a systems overview. This issue of the RCA Engineer is the result.

COMMAND, control, and communications is, indeed, a systems area of great interest to RCA engineers and to the Corporation. First, it is an extremely large business area (see Table I) which closely matches our corporate abilities. Second, military C3 systems and technologies have become the precursor of civilian communication and control systems and hardware — another major RCA business area. Communication satellites, automatic switching and checkout, operator consoles and displays, and the latest handheld or mobile radios are typical examples of this technology synergism. Finally, it is an area that appeals to the creative engineer for the technical challenges it offers and the opportunity to contribute to meaningful progress in our national security and economic and social welfare.

Military C3 in all its forms is so complex and diversified, encompassing so many technical areas, that it takes multi-volume reports to describe all of its ramifications. This brief overview attempts to summarize what it is all about, how the US military C3 system is organized, some of the problems it is grappling with, and to point out technology areas where continued progress will lead to further gains in effectiveness. Subsequent papers in this issue will deal with specific C3 system elements and technologies.

Because there is an oversupply of C3 definitions, it is advisable to agree on this matter before proceeding further. Without arguing the merits of the many definitions offered in the past by individuals and agencies,1 it is probably safest to choose the following latest definitions issued by our highest military authority as follows:2

1. **Command and control system** — The facilities, equipment, communications, procedures, and personnel essential to a commander for planning, directing, and controlling operations of assigned forces pursuant to the missions assigned.

2. **Communications** — A method or means of conveying information of any kind from one person or place to another. "

3. **Telecommunications** — Any transmission, emission, or reception of signs, signals, writing, images, and sounds or information of any nature by wire, radio, visual, or other electromagnetic systems.

A more pithy and all-embracing definition attributed to an unidentified individual is, command and control is running the show.3

---

**Table I** — DoD C3 expenditures ($million) for hardware and software. Defense Communications Agency expenditures are allocated to the military service which did the contracting, mostly the Air Force.

<table>
<thead>
<tr>
<th>Service</th>
<th>FY 1946-72</th>
<th>FY 1973-77</th>
</tr>
</thead>
<tbody>
<tr>
<td>Army</td>
<td>11,000</td>
<td>2,720</td>
</tr>
<tr>
<td>Navy</td>
<td>19,000</td>
<td>3,470</td>
</tr>
<tr>
<td>Air Force</td>
<td>43,500</td>
<td>5,530</td>
</tr>
<tr>
<td></td>
<td>73,500</td>
<td>11,720</td>
</tr>
</tbody>
</table>

---

**Fig. 1** — The military operational process.
Military operational process

The military operational process (Fig. 1) is the indispensable common element of the C³ function for every type of military force: land, sea or air, at every command level. It is the essential antecedent of every combat or support action of whatever nature or scale. The four sequential elements of the process may be briefly described as follows:

sense — Gathering of all that information of the environment and friendly and enemy forces necessary to conduct a military operation.

evaluate — Determining the credibility and significance of information gathered and its correlation into a situational picture.
decide — Choosing the best course of action to take in order to accomplish the assigned mission.

act — Issuing orders and the subsequent actions taken by the commanded forces to execute the decision.

While the fundamental process has a universal application to all military situations and operations, there is wide diversity in its forms and actual employment. It is obvious that there will be great differences in the “what” and “how” of the process between the Joint-Chiefs-of-Staff level and an infantry battalion in combat in Vietnam, between anti-submarine operations in mid-Atlantic and air defense operations in NATO, and even between forces having the same basic mission as in the case of land or sea based strategic offensive missiles and aircraft. The nature and methods of the process will also vary for the combat and the support (supply, maintenance, personnel, movement, etc.) operations at any given command level and type of force.

The men, hardware, and software of the military C³ systems are the physical manifestation of the sense-evaluate-decide-act process. It is in comparing our present-day capabilities to execute this never-ceasing process with modern electronics to the historic past that we begin to realize the scope and magnitude of the C³ revolution.

For millenia, the sense-evaluate-decide elements of the process were entirely human and usually embodied in one man — the military leader on horseback or the naval leader on the quarterdeck of a ship. In strategic planning and operations, he made his evaluation and decision based on information reaching him from a variety of distant sources. In tactical planning and operations, his own visual sensing of the scene of operations provided the information to be evaluated to arrive at a decision. The military action resulting from the decision was again basically human, with whatever assists the mechanical technology of the age could provide. Communication was wholly non-electronic and line-of-sight.

Now, our strategic decision makers are

David Shore Division Vice President, Government Plans and Systems Development, Moorestown, N.J., was promoted to his present title in May, 1971. Formerly Manager of Government Plans and Systems Development since 1969, Mr. Shore continues to direct the planning activities for G&CS and development of major new weapons systems concepts. He is also responsible for supporting the major program efforts of the five operating divisions of G&CS. Mr. Shore received the BS in Aeronautical Engineering from the University of Michigan, 1941; and the MS in Physics from Ohio State University, 1950. Before joining RCA, Mr. Shore was with the United States Air Force, from 1941 to 1954, in various positions at Wright Air Development Center, Wright Patterson AFB, Ohio. As Civilian Chief of the Systems Liaison Office from 1950-1954, Mr. Shore was responsible for conceiving new aircraft and guided missile weapon systems for the Air Force. Mr. Shore joined RCA in 1954 and became Manager of Systems Synthesis in the Missile and Surface Radar Department. In 1961-2, Mr. Shore was made Chief Systems Engineer and Manager of SEER (Systems Engineering, Evaluation and Research) located in Moorestown, N. J. He was appointed Chief Engineer, Communications Systems Division (now the Government Communications Systems) in August 1965. In December 1966, Mr. Shore was promoted to Chief Defense Engineer of Defense Electronic Products. In this capacity he was responsible for the management of the DEP IR&D Program, establishment of engineering policy for the engineering activities, DEP divisions, and the direct management of Advanced Technology, Central Engineering, Defense Microelectronics, Systems Engineering, Evaluation and Research activities. His memberships include: American Institute of Aeronautics and Astronautics, American Ordinance Association, Aerospace Industries Association, Institute of Electrical and Electronic Engineers, Air Force Association, Armed Forces Communications and Electronics Association, Association of the U.S. Army, Army Aviation Association of America, and the National Aviation Club. Mr. Shore holds a Professional Engineering License for the state of New Jersey.
served by multi-billion-dollar global electronic sensing and communications systems and large, complex command and control centers in which the most advanced electronic technology plays a central role in the evaluation-decision process (Figs. 2,3). The commanders of land, sea and air tactical forces are also served by electronic sensing and communication systems no less advanced than those of the strategic forces. In some instances their evaluation-decision process is also based on electronic sub-systems and in others remains a human-oriented action, as I will discuss later. (Figs. 4,5,6).

US military C^3 systems

The total US defense C^3 facilities and systems may be broadly divided as follows:

- Those required by the National Command Authorities (Fig. 7).
- Those required by the Department of Defense for the military operating and support forces (Fig. 8).

The C^3 system of the military forces may again be divided into two broad categories with differing requirements for communications and the sense-evaluate-decide-act process.

- Those required by strategic offense and defense forces.
- Those required by general-purpose (tactical) forces.

Both the strategic and general purpose forces are further divided into operating and support elements whose mode of operation and operating medium again introduce unique requirements. These divisions usually are:

- Land forces
- Air forces
- Sea forces
- Joint forces

This division and subdivision of our combat and support forces into operating systems tailored to solve specific military problems (which in themselves keep increasing in number) has led to multiplicity of C^3 systems. This creates an interoperability problem in trying to integrate all of the required functional systems into a total overall system. While we still have a long way to go to reach the ideal, much progress has been, and continues to be, made.

The whole military C^3 structure is hierarchical (with some exceptions, as might be expected). While sensing (information seeking) goes up, down, and laterally in the hierarchy, decisions (orders to act) always flow down an ever-broadening chain of command, receiving amplifying instructions on the way. Thus, the simple World War II order of the Combined Chiefs of Staff to Gen. Eisenhower to land on the coast of France in the Spring of 1944 and defeat the German forces eventually reached an infantry company commander as an order to land on a specific area of Omaha beach, advance and capture a designated piece of terrain; it also reached a fighter bomber pilot as an order to be on station at a definite time to give the company commander close air support; and the captain of a ship received an order to be in position to provide naval gunfire support.

The growing contribution made by electronic technology to the evaluate-decide progress was noted previously. This application of technology is most advanced in those combat operations where the inputs from the sensing function can be quantified and digitized for computer evaluation and decision-making and the result symbolized for presentation to human controllers, and also where the allowable time interval between sense and act is too short to permit human evaluation and decision making. Thus, we find a high degree of evaluate-decide automation in our strategic offense and defense systems and in those tactical systems (such as air defense) which pose similar operational problems. The AEGIS fleet air defense system developed by the Missile and Surface Radar Division at Moorestown is a splendid example of the great gains in combat effectiveness that can be achieved when the most advanced
electronic technology is applied with skill and ingenuity to automation of the sense-evaluate-decide-act process.

Support operations have also proven amenable to a process automation because numbers are the essence of this function: quantities of people or things needed or on hand, weight and size, consumption rates, transportation modes, capacity and speed, locations and distances, time to service, repair, load, etc. Consequently, this was an early field for C$^3$ automation and one in which continuing refinements are still being made.

Likewise, automation of both military and civil telecommunications has made significant progress utilizing the increasing bandwidth available from the use of higher frequencies and satellites.

In some areas of general-purpose-forces (limited warfare) operations, most notably land combat, there has been little progress in developing C$^3$ systems comparable in task execution and performance to those now employed by strategic forces. While the past decade has seen great progress in this area in the development and use of electronic devices for sensing and communications, the evaluate and decide process present problems for which a solution is yet to be found.

There are a number of reasons why we have failed thus far to automate the land combat evaluate-decide process satisfactorily. One fundamental reason is the nature of the problem as stated below:

"When the game he (the commander) is playing is continually subject to unpredictable change, there is simply no information technology that can be called upon to cast the commander's problem in the form of a set of binary choices."

"The crux of the matter is that our commanders—and, derivatively their information systems—must deal with contingencies. Each contingency is a choice-point, and the possible paths ahead at any moment ramify so rapidly that detailed pre-planning that is adjusted to contingent events is literally impossible." 4

This inability to define with precision the problem to be solved leads to a computational requirement of the awesome magnitude described below:

"Dr. Nicholas Smith of the Research Analysis Corporation, in his pamphlet "Operations Research in the Next Twenty Years: A Technological Forecast" (1964), estimated that the large number of variables—men, weapons, enemy, terrain, and weather—involved in the actions of a battalion-size unit in the attack would require $10^{45}$ computations to determine the best tactic by examining the effect of all possible combinations of factors. This number is one trillion multiplied by itself over 41 times. Even after discounting the irrelevant and trivial possibilities, the number of conceivable combinations, remains unmanageably large."

At best, C$^3$ can provide this tactical commander with the latest data, in the clearest form, at his greatest convenience. The total automated solution is not in sight.

A lively current issue in C$^3$ system design, particularly in those systems intended for tactical forces, is whether they should be centralized or decentralized. At the beginning of the era of digitized automation of C$^3$, all systems were centralized, both by necessity and by inclination. The space required by the available general-purpose computers and their peripheral equipment and the controlled environment necessary for their operation called for large permanent installations. At the same time, there was the traditional desire of a commander to retain the greatest possible measure of control over all of his forces and their operations. Automation of C$^3$ systems seemed to make this possible in a degree never before attained. The development of strategic C$^3$ systems thus went forward rapidly because these two conditions were not contrary to any system requirement.

Tactical C$^3$ lagged strategic C$^3$ because mobility was essential. As technology reduced the size of command and control elements, tactical C$^3$ began to evolve. The aforementioned multiparametric problems forced the solutions into specific categories of operations such as fire control.

Full dependence on a central tactical C$^3$ system is unlikely even should technology ever permit it. This central C$^3$ would provide too lucrative a target to the enemy. Technology has, however, opened up the alternative of a federation of smaller systems netted together by communications. As an example, the RCA Space Ultra-reliable Modular Computer (SUMC) being developed by G&CS Advanced Technology Laboratories for NASA (Fig. 9) replaces a commercial machine that is 15% larger, 25% heavier, and 40% more power consuming.

C$^3$ technology

The continued advances now forecast in electronic technology will upgrade the effectiveness of our present C$^3$ systems and very likely lead to new advanced systems. Space does not permit a thorough treatment of the probable impact of new technology on C$^3$ systems. Therefore, I will limit myself to some of the key areas where expected technical progress will have a high payoff or where dedicated effort is needed to overcome present C$^3$ system problems.

Data processing — hardware and software

The revolution in electronics—starting with tubes, achieving the large-scale integration of solid state devices which has made the SUMC computer possible, and predicted to yield a computer-on-a-chip in this decade—is not over. But it will now be focused on ameliorating the software problem.

Software is the most troublesome current problem in C$^3$ system design and development. It is the primary reason why some systems in existence or under development have not achieved their performance goals. It is now also the most costly part of a C$^3$ system, and software costs continue to grow while hardware costs continue to decline. In some cases, software now accounts for as much as 80% of a command-control system cost.

The normal procurement cycle for a new C$^3$ system requires early estimation of data processing capacity required. Despite the improving ability to make such estimates and the provision of reserve capacity of 50 to 100%, there are
an astonishing number of instances where the computer procured is found inadequate when the software is finally developed. This has caused expensive reprogramming to utilize the computer most efficiently, adding mainframe and memory capacity, and/or adding a separate system to share the load.

Much research is under way to cope with the software dilemma. (see several articles appearing in this issue).9,10,11

A new viewpoint growing in the community is the idea of "software first". This means software is developed on a large, general-purpose computer and then a C3 computer is selected or developed with good assurance its capacity will be adequate. Software can be translated through a higher order language in the case of an existing computer. With the rapid evolution of design automation and large-scale integrated circuits, it is also possible to develop a C3 computer organized to use the software directly—yet at a reasonable cost and schedule.

Technology is also expected to relieve the software problem. First, mini- or microcomputers will perform parts of the C3 function; software then can be broken down into a number of smaller, more manageable units. Second, very large solid-state memories will be made economical by LSI technology. This permits compartmenting the memory into different program elements and easier programming.

The commander, today, is fairly remote from the software problem. Worse, he feels isolated from his automated system because programmers handle all software. The future may eventually see the commander able to address his system by voice and personally control its actions.

Communications

The last "C" in C3 is communications. Without it the commander, of course, cannot command. The human runners employed by the ancients are largely replaced by electromagnetic messengers covering a frequency domain from a few hertz to light. As the bandwidth demands rose to permit computer-to-computer communications, there has fortunately evolved satellite systems which will eventually provide global, wideband services.

Military C3 communications make use of every other applicable civil system as well. However, the military must strive for security against enemy snooping or disruption. A number of dedicated, secure modes have been developed for this purpose. It is in the never-ending game of measure vs. countermeasure that technology challenges remain. Frequency allocations are difficult to obtain because of the many demands on the total spectrum from a host of international claimants. As a result, frequencies in the K-band and above are being opened. These are less desirable for tactical C3 use because of weather attenuation.

It should be mentioned that voice communications are still needed. Most commanders derive important situational clues by listening to how their subordinates sound in their verbal reports.9

One author comments on this problem:

"We have spent billions of dollars increasing the number of channels, speeding the bit-rates, adding memory capacity, and improving the signal-to-noise ratio; but we have paid scant attention to the problem of how we are transferring meaning from one mind to another."4

Displays

The displays in C3 fall into two categories: small for individuals, and large displays for groups.

The individual displays are normally TV consoles with black-and-white being replaced by color. Computer-driven displays are now capable of showing video images, graphics, and/or alphanumeric information.

Large-scale displays, such as in Fig. 2, have depended on projection techniques and suffer from inadequate resolution and light intensity requiring the room to be dimmed. Research is progressing on a number of new concepts such as plasma and liquid crystals which should correct the problems of today. In addition, they could be computer-driven to display dynamic, real-time situations.

In addition to voice and displays, hard copy is usually prepared for record keeping. Even in the case of automated fire control systems for the artillery, hard copy is specified to permit later review of results compared to intent.

System diagnosis and tests

All C3 systems of whatever size, nature, or purpose must have built-in automatic diagnosis and test of all subsystems and elements under central computer control. The notable achievements of the AEGIS system in this regard must become standard in all future systems.

Summary

Hopefully, this broad overview provides a frame of reference to the reader of this issue for the articles that follow. Some key points to keep in mind are:

—C3 exists to assist the commander.
—Commanders vary from a Patton to a Montgomery. C3 must adapt to the commander, not vice versa.
—Technology is improving the ability of the C3 system to assist the commander, not replace him.
—Software is the pacing, riskiest part of C3.
—The improvement of C3 appears to be a never-ending challenge to systems and equipment designers.

References

3. Parson and Perry, op. cit. p73.
Human element in command and control systems

Dr. F. H. Ireland

Man's role as decision maker is central to the concept of any command and control system. How much man participates in the operation of such a system and how much relevant data the system provides him in making his decision depends on the system performance requirements, the cost and complexity of the system, and the backup needs. Thus, in considering man's role, the human factors engineer must weigh these tradeoffs as well as the design of the man-machine interface. This paper discusses research that has attempted to clarify the human decision-making process and provides some insight into where further work is needed.

In attempting to discuss man's roles, functions, and problems in command and control systems it would be well to first indicate what is meant by "command and control". Originally the label was attached to a wide variety of military systems. Most of these were intended to facilitate operational control of forces (e.g., NMCS, SAC, TACS, CCIS-70, NTDS or MTDS); but others were designed for such diverse purposes as resources management (473L), missile launchings (BMEWS), space traffic (Space Track) and electromagnetic intelligence (466L). More recently, the Department of Defense approved a definition of command and control system as the facilities, equipment, communications, procedures, and personnel essential to a commander for planning, directing and controlling operations of assigned forces pursuant to the missions assigned.

As implied by its name, a command and control system performs two distinct functions:

Command deals with broad problems and policies and with the setting of objectives. Control, on the other hand, refers to the directing of activities toward objectives that are already clearly defined.
The command function predominates in the systems intended for the higher military echelons; at lower levels the emphasis tends to shift more toward control.

Automation and the changing role of the human operator

Automation has had a considerable impact on command and control systems and has led to changes in the roles played by the human operators. The effects are particularly pronounced in systems in which the control functions are emphasized. In these systems, automated equipment has largely taken over such previously human activities as target detection, acquisition, correlation, tracking, identification, threat evaluation, intercept prediction, weapons assignment, and guidance.

The basic command function, on the other hand, is still exercised by a human commander, who is responsible for the final decisions. He, in turn, is supported by other men who form his staff and are responsible for:

1) Acquiring information on the status, disposition and actions of their own and hostile forces, on logistics, on weather, and on the operability and availability of weapons and communication facilities;
2) Evaluating, analyzing and interpreting the information in terms of assigned missions and objectives;
3) Developing hypotheses about enemy intent, planning response alternatives, and estimating attendant effects and costs; and
4) Submitting action recommendations to the commander for his decision.

In some of these activities, the staff members can be, and have been, supported by automated equipment.

In fact, automation of the information acquisition, manipulation, storage, and retrieval processes has already made considerable inroads at the higher command levels where data processing equipment and automated displays are now in common use. Furthermore, due to the rapid development of large-scale integrated circuitry and the concomitant reductions in size, weight, and power—miniature general-purpose battlefield computers can now be packaged together with keyboards, hardcopy printers, and "soft copy" displays in suitcase-size containers. When connected to automatic data links, these devices can process messages automatically and their computational capabilities make them useful for the solution of operations control problems. It appears quite possible that such units, augmented with external memory, could be used for the processing of large amounts of intelligence data to support command decision-making in the field. Someday we may even see such equipment playing a more direct role in the formulation of higher-level decisions.

Machines, on the other hand, are sensitive to many stimuli that man is unable to perceive; they can respond more quickly to control signals, and they are not plagued by boredom or fatigue when performing routine, repetitive tasks. They can store and recall incredibly large amounts of data and perform complex computations rapidly and with perfect accuracy. The intriguing question of whether, and to what extent, a machine could or should eventually participate in formulating command decisions has often been debated. Some of the relevant experimental work which has been done to date will be described later.

Cost tradeoffs

Economic considerations also enter into the man-machine choice. If the performance of a function, which could be accomplished by either a man or a piece of equipment, were to represent a full-time job for a man, it might well require a four- or five-man complement for around-the-clock manning. The recruiting, training, and maintenance
costs for such manpower is not trivial. On the other hand, if a function which could be performed either by man or machine, can be assigned to a man already in the system without overloading that man, it may be desirable to forego automation of that function.

System back-up

The capability of a system to sustain equipment damage or malfunction and still be able to function, albeit with reduced effectiveness, is another important consideration when determining the desirable degree of automation. So-called graceful degradation can often only be accomplished by using men to back up the automatic operation. To be instantly useful and effective in emergencies, these men must not only be available on a standby basis, they must also be given all the up-to-the-minute data needed to assume the operational responsibilities. Further, since time for a “cold start” is usually not available, the operators must monitor the system while it is operating in the automatic mode. Maintaining the proficiency and morale of back-up personnel—who may only rarely, if ever, be called upon to perform operationally—is a challenging and frequently frustrating human factors problem.

Interface design

Equipment designers are interested in an early assessment of the numbers and types of operating positions in the system because each operator usually requires a console or other man-machine interface equipment. Therefore, once the human functions have been defined, it is necessary to determine the number and types of operators needed to carry them out. Computer-based simulation programs have been developed to help human factors specialists gauge anticipated operator workloads and manpower requirements. In some cases, several logically or procedurally related human functions can be allocated to a single operator. In other cases, it may take several operators, working in parallel, to carry the workload imposed by a single system function.

Defining information requirements

The next problem is to determine exactly what information each operator needs to perform his tasks. The data needed for control purposes is much easier to define than that for command. After all, the thing to be controlled (such as a radar antenna, an aircraft or an infantry battalion) is clearly defined as is the objective to be attained. Information needed for control can be fairly accurately determined by a detailed task analysis. During that process, each human function is broken down into sequences of discrete task elements. The information items needed for each element are then specified as are the action requirements. It is questionable whether this analytic approach can be used successfully to determine information needs for command decisions as well. At first thought, it seems relatively simple to outline the range of decisions that a commander must make and then to determine the minimum set of information items required for these decisions. However, in the typical strategic or tactical command situation, the objective is usually not clearly defined. There is often considerable uncertainty concerning the significance of data, the intent of the enemy, or the likely cost and effectiveness of various action options. The real difficulty in defining information needs for command decisions stems from the fact that the human decision-making process is still very poorly understood.

Although this process is yet a mystery, we do know that good decisions require reliable, relevant and up-to-date information; lacking such information, a commander might just as well flip a coin. What information should we then provide for a commander? If we simply ask him what information he thinks desirable he will probably answer something like “all of it”. What he really means is that he cannot say beforehand what data he will want in order to evaluate a new threat, and that he would rather be faced with the problem of having to filter the irrelevant data than to chance the possibility of its insufficiency.

Results of experimental work

Suppose it were possible somehow to determine all information items relevant to a decision, as we can more or less do in an experimental situation. Would it really be beneficial to present all of them to the decision maker? This question was investigated by Hayes in 1962. He found that as the number of relevant information items increased, the time required to make a decision also increased, but the quality of the decision did not improve. Furthermore, when the time available for making the decisions was held constant, the decisions actually deteriorated as more and more relevant information was supplied. Thus, it would appear that too much information, even if it is of the right kind, can be a bad thing. However, if it was really the quantity of information that created the problem for the decision makers in the Hayes experiment, predigesting and summarizing the data prior to display should improve matters.

The merits of this approach were investigated in 1964 by Merifield and Erickson. In their experiments, the decision-making subjects were required to discover an enemy’s strategy underlying his missile attacks on various cities. After several cities were reported hit, a subject was asked to judge which of his nation’s resources the enemy was attempting to destroy and to predict which cities would be attacked next. Before any attacks occurred, the subjects studied a large data matrix showing the values of fourteen different resources located at each of forty-nine cities. One group of subjects also had access to a display which presented computer-generated statistical summaries of the information which were based on correlation, factor analysis, groupings and variance ratios, which another group did not. As one might suspect, the group aided by the statistical summary outperformed the other.

To really find out what information a commander needs in the field, it may be necessary to observe his decision-making process—at least under realistically simulated conditions such as during maneuvers—and have the commanders attempt to describe their procedure in detail. Despite the fact that many decision-makers find it difficult to verbalize the mental steps involved, such an introspective method applied to the few who can do so may yield more worthwhile data than analytic techniques or so-called objective behavioral studies. To the extent that the decision-making procedures educed in this fashion can be reduced to algorithmic form, they can be modeled and validated. Conceivably, they might find use as decision-making aids, perhaps in the form of recommended actions.
Other human factors constraints

The remaining human factors tasks involved in designing a command and control system are essentially the same as for any man-machine system, and adequate techniques for accomplishing them exist. These tasks include design and specifications of

1) Man-machine interface equipments such as the display and control panels and entire operator consoles;
2) Display formats;
3) Operations and maintenance area arrangements including consideration of operator environmental factors such as lighting, noise levels, and air conditioning;
4) Operating procedures and special job aids; and
5) Training, exercising and proficiency evaluation means.

To verify the designs and to discover possible shortcomings, human factors specialists should also play a key role in the system tests and evaluations.

User participation during system development

The participation of the end-user is virtually essential during the development of a command and control system. The designers can gain an appreciation of the user's real needs and concerns, and the chances for ultimate user acceptance of the system are increased.

The human factors specialist can often precipitate or substantially enhance such user involvement in the design process by providing early mockups of proposed consoles and displays. To be useful design tools, mockups should be full-scale, arranged in an area that simulates, as closely as practical, the environment in which the equipments and operators will ultimately have to work. During the early stages of system design, inexpensive static mockups made of pressboard or similar material can be used. These can be fitted with artwork panels showing the layout of displays and controls. Such displays should show the proposed information content and format. Later, when design approaches are better defined, more substantial mockups with front panel hardware should be substituted. At that stage, partial instrumentation of the mockups and dynamic simulation of displays may prove valuable.

The mockup area should also contain up-to-date system flow diagrams, artist concepts and scale models of equipment, a word-picture or a voice recording describing a typical operation and other audio visual material for briefing future users and other visitors. Such mockup and briefing facilities, if constantly updated, can also serve as a valuable catalyst for the integration of the many concurrent, but disparate in-house design efforts.

Research on human decision-making

Human decision-making was previously pointed out as a persistent problem area that has been surrounded by much confusion. The main reason for the confusion is that the term decision-making, though often used, is very poorly defined. It has been applied to behavior which ranges all the way from determining which of two buttons to press in response to a simple light stimulus to a highest level judgment on whether to engage the nation in a war.

In the command and control system context, it is useful to distinguish between decisions made for purposes of control and those involved in command. The process underlying the former is much better understood, and analytic methods are available for determining the decision-maker's information requirements. In fact, algorithms for control decisions can often be clearly enough defined to permit automation of the entire process. On the other hand, much less is known as yet about how a commander makes the higher-level decisions when confronted by considerable uncertainty about his information and the possible consequences of his acts.

Questions about the basic process involved in command decision-making are indeed among the most challenging faced by human factors investigators who have already devoted considerable effort to this area. For example, a team at Ohio State University spent several years on attempts to determine how human subjects made certain types of diagnostic decisions and to see whether computers might be used to aid or replace men for such purposes. A Bayesian decision model was used. For each of several potential tactics available to an enemy, the experimental subject was given probabilities attached to specific observable actions which an enemy might take in order to implement the tactics. Thus, for a pincer attack, the subject might be told the probability of enemy troop concentrations in certain areas, the probability that he would use tanks or observation planes, etc. The specific probabilities assigned to each observable action were different for the various possible enemy tactics. The problem for the subject was to induce composite conditional probabilities for the various possible enemy tactics from the known conditional probabilities of the discrete enemy acts reported to him.

As might be expected, the results showed that men usually do not extract all of the information inherent in the data. Surprising, however, at least to this writer, was the fact that a computer only performed about 10% better than men in aggregating the various conditional probabilities. Furthermore, as the men became more accustomed to the somewhat foreign task and were able to check their estimates against those produced by the machine, their performance began to approach that of the automatic equipment.

An obvious problem with this approach is that, in a real-world situation, it would be very difficult to obtain the original conditional probabilities for observable enemy acts as a function of alternative tactics. It might be possible to derive some probabilities from historical data, but most of the input probabilities would have to be based on expert opinion.

Some have argued that threat evaluation does not really constitute decision making at all since it does not involve action selection. In any event it is clear that command decision-making involves considerably more than merely the diagnosis of a situation.

Many other experiments have, in fact, been conducted in attempts to discover how commanders make decisions not only about the threat situation but also about their own tactics, the selection of weapons, the timing of countermoves, and conservation of their resources. Perhaps typical of these is a series of investigations begun in 1959 by the Operations Application Laboratory of the Air Force Systems Command at an annual cost of $250,000. In a simulated air defense environment, military subjects with prior air defense experience played the role of a commander who was
required to evaluate the threat posed by attacking aircraft or missiles and select appropriate defensive actions.

Among other things the experiments showed that

1) While subjects were able to devise successful strategies, they were generally unable to clearly verbalize them;
2) Their performance did not level off as expected with increased decision-making load (i.e., performance rate increased with load), and there was no clear cut breaking point at which performance suddenly deteriorated;
3) There was a tendency to use up and even squander available weapons when the supply appeared to exceed the demand; and
4) Subjects changed about one half of the previously made decisions and often unwittingly those which they had made themselves.

But while investigations of this type undoubtedly have yielded interesting and useful information, none of these have, thus far, been able to provide any clear insight into the higher-level human decision-making processes involved in command.

An approach that may some day prove to be valuable is to provide the decision maker with a computerized simulation tool by means of which he can, in an accelerated time frame, test the probable effect of various potential courses of action before issuing his commands. The idea is similar to one for an automated aid to a chess player which would permit the latter to look several moves ahead, both with regard to his own as well as his opponent's possible actions. This is, of course, difficult enough to do for a chess game which represents a clearly defined artificial battlefield with rules that regulate the movement of each piece. Simulation of real battlefield situations, with far less rigorously circumscribed action potentials on either side, poses an enormously greater challenge.

Suppose a commander had a computer available which could automatically determine the appropriate counteractions at least for a limited set of threat situations. Would he accept and use such a device—would he, in effect, let it make decisions for him? A series of experiments which were initiated in 1962 by the Applied Physics Laboratory (APL) of the Johns Hopkins University and which lasted four years, attempted to answer this question. The setting was a simulated shipboard combat information center (CIC) and the task again was to counteract airborne enemy threats. Eighty-one naval officers, ranging in rank from captain to lieutenant participated as commander subjects. In addition to the customary situation and status displays found in a CIC, the commander was provided with a display showing the computer-recommended actions. Although the subjects were never aware of the fact, no actual computing equipment was used. Instead, one of the experimenters, trained to work in accordance with computer program rules, simulated the computer and provided problem solutions which would have been produced by such equipment.

With regard to the "computer" solutions, a commander had two options:

1) By doing nothing he could accept the computer recommendation, in which case the recommended action would be executed automatically, or
2) He could reject the recommendation within a certain time period and substitute action commands based on his own judgment.

During and following each experimental run, qualitative data were gathered from subject remarks and personal observations of the experimenters. In general, the subjects considered the experiments themselves quite educational and felt that the experience would help them in their operational tasks. But, while the commander welcomed the idea of computer aiding, the experiments revealed a wide range of individual differences as to the extent they actually accepted the computer-generated tactical action recommendations. Manual override often occurred because a commander felt that he had too little time to evaluate the recommendation. Nevertheless it appeared that a commander could comprehend and evaluate up to five computer recommendations within a half a minute and that most officers approved of the manual override method adopted as opposed to the possibility of requiring a positive action in order to accept a recommendation. It was felt that the latter approach would let an indecisive officer delay action too long.

Need for further work

Obviously much research on decision-making remains to be done. In particular, human engineering requires a set of basic principles in this area on which to base its recommendations for specific command and control system designs. Approach as an engineering problem, development of some useful knowledge, at least with regard to command information presentations, should not require such extensive experimentation as that described above. For example, assuming that a commander will want to know the current status of his own resources as well as those of his enemy, how should that information be formatted on a display so that it can be easily assimilated? Should it be done graphically or in the more conventional tabular fashion? If the graphic approach is taken, which data should be shown in the form of curves, which as bar-charts, or in some other pictorial manner? A cataloguing of such engineering related questions and problem areas is needed so that the required research can be programmed effectively.

References

Remotely piloted vehicles —
command and control

D. Shore

The latest war in the Middle-East demonstrated the lethality of air defense against manned aircraft. At the time of this writing, the cost in lives and in aircraft are not known in detail, however, it is clear that these losses were significant. As a result, increased emphasis is being paid in this country to the utilization of the unmanned remotely piloted vehicle, or as it is more generally known, the RPV. This paper discusses some of the milestones of RPV development and highlights significant RCA work in progress. We plan to provide more information in future issues on the RPV and RCA's technical contribution to this challenging new field.

The radio-controlled aircraft is not new; indeed, Hugo Gernsbeck in the March/April 1931 issue of TV News published an article on the use of radio-controlled aircraft with a TV sensor in the nose for military purposes. At that time, both radio control and television were laboratory curiosities.

In World War II, war-weary B-17 bombers equipped with RCA television and a radio relay were used in "kamikaze"-type, unmanned attacks by the U.S. Air Force against the heavily defended Helgoland fortifications of the Germans. In the post-World War II era, military developments were concentrated on guided missiles in which man participated in mission planning but had little or no "hands on" control of the missiles once launched. There continued to be experiments utilizing obsolescent manned aircraft in an unmanned mode with television in the nose of the airplane and control accomplished by a pilot on the ground via a two-way data link. These experiments, while successful, were not exploited, primarily because guided missiles do not require communications from the ground that could be jammed.

The war in Vietnam forced reconsideration of the RPV as an important part of the weapon arsenal. The Soviet surface-to-air missile systems, the SA-2 and SA-3, coupled with radar-directed anti-aircraft guns took a heavy toll of the manned aircraft used by our Services. The Services then began to utilize drones for reconnaissance. These drones were adaptations of target drones which had been developed to exercise our own air defense systems. In the final phases of the war and in the negotiation for peace, the prisoner-of-war issue further stimulated the search for a replacement of manned aircraft penetration over highly defended enemy territory. The RPV was recognized as a logical solution.

RCA began internal study of the RPV in 1970. It was clear to us from the first that the key issues to be overcome before the RPV could be considered a viable combat weapon were:

1) Secure, anti-jam, two-way communications.
2) Multiple RPV command and control.

Furthermore, solutions to these problems had to be achieved at a low cost commensurate with the limited life expectancy of these vehicles in combat. Our efforts to date have been recognized by the Air Force in three significant contracts covering the design of a viable command, control, and communication system capable of handling up to twenty RPVs per system. We have been able to develop cost-effective solutions to the issues mentioned above and are in the process of demonstrating their practicality through the medium of a hardware system demonstrator (Fig. 1).

As one might expect, the man-machine interface problem for the RPV is, indeed, a challenging one. The pilot sitting in front of a display console is expected to fly the RPV using television eyes located in the aircraft. He must recognize the target, maneuver the vehicle to an aimpoint, and direct his weapon with high accuracy against a heavily defended target. The RCA demonstrator (shown in Fig. 1) is being built to provide us with an experimental tool to enhance man's ability to operate the RPVs. This demonstrator will also be useful in training RPV pilots, through the use of simulation, without the risks of actual flight. This is another cost saving byproduct of the RPV concept.

---

Reprint RE-19-5-25
Final manuscript received January 4, 1974
Mr. Shore is Division Vice President, Government Plans and Systems Development, Moorestown, NJ. His biography and photograph appear with his other paper in this issue.
Managing computer program development
S.A. Steele | R.A. Dupell | A. Fleishman | E.T. Hatcher

Traditionally, the computer software industry has tended to perpetuate basic errors in computer program development, carrying them over from one project to another. Studies have shown that problems of software productivity on medium or large projects are frequently attributed to inadequacy of system design and project management. In recent years, MSRD management has recognized the need for reevaluation of the traditional approach to computer program development and management, and has researched the field extensively to determine the most effective approaches and techniques. A task force of experienced personnel was formed to develop a disciplined, total system of problem definition and resolution in software systems development. The computer program development and management system (CPDAMS) evolved from this effort and is a workable system applicable to a wide range of computer program development projects.

Dr. Stuart A. Steele, Mgr, Command and Control Systems Engineering, MSRD, Moorestown, N.J., received the BEE from Bucknell University in 1958, and the MS and PhD in electrical engineering from the Pennsylvania State University in 1961 and 1965 respectively. From 1959 to 1961 he was involved in digital data acquisition system design with Electronic Associates, Inc., and from 1960 to 1966 he served as an assistant professor of electrical engineering at the Pennsylvania State University, in charge of the digital computer facility. In 1966 he joined the General Electric Space Technology Center, where he served in various capacities through 1969. His principal activity involved evaluation of large-scale digital simulations with man in the control loop. He also participated in various spacecraft control studies and project development. Dr. Steele joined RCA in 1970, with responsibility for a variety of development tasks involving real-time software systems. In his present position he is responsible for the division's computer center, application programming, and all software engineering. Prior to this he was involved in the development of major software proposals at MSRD and has been active in the development of MSRD's technological base in the computer and software field. Dr. Steele is a member of IEEE, American Society of Engineering and Education, American Association of University Professors, and Sigma Xi. He has published more than 20 papers in the areas of computer and control systems, and has served for the past several years as a part-time lecturer in control systems and computer organization at the Pennsylvania State University Graduate Center in King of Prussia, Pennsylvania.

Allen M. Fleishman, Program Management Staff, Command and Control Systems Engineering, MSRD, Moorestown, N.J., received the AB from Bridgewater College and the MS in 1961 and 1965 from Virginia Polytechnic Institute and State University. He joined RCA in 1964 and has been active in the development and management of software systems for the Navy and Army. He has served in a number of capacities with the Navy and Army, including as a software systems representative and was responsible for computer systems software installation and application guidance for the RCA's first general purpose computer customer, the Army Tank and Automatic Command. From 1967 through 1969 Mr. Fleishman held various computer software development and computer systems marketing management positions. In 1982 he assumed responsibility for the computer program development of the AFAR Global Communications System, at that time the world's largest on-line data communications message switch. He was also responsible for the computer program development of the worldwide RCA Global Communications Electronic Telegraph System in 1984. From 1986 through 1971 he served in a number of computer management positions within the RCA Data Processing Division. During 1973 Mr. Fleishman was a consultant with Equimatics, Inc., providing communication consulting services to the life insurance industry. Since rejoining RCA in 1973, he has been engaged in several large data communications studies.

Raymond R. Dupell, Ldr, Software Management, Command and Control Systems Engineering, MSRD, Moorestown, N.J., received the BS from the University of New Hampshire in 1942 and the EdM from Boston University in 1951. He joined RCA in 1973, involved in the development of the Computer Program Development and Management System (CPDAMS). He conducted a series of workshops to develop the system, and published the CPDAMS Standard which has had extensive distribution within RCA and to a number of agencies of the defense establishment. Subsequently, Mr. Dupell developed the AFAR Project Management System, an adaptation of CPDAMS, which is currently being used in the development and management of the AFAR Program. From 1969 to 1972 Mr. Dupell worked for Honeywell Information Systems as a Senior Staff Analyst and Group Manager of Computer Resources. He also managed Honeywell's Boston Computer Operations to provide continuing support for both software and hardware development. Mr. Dupell retired as a Colonel from the Air Force in 1969. During his Air Force career, he had extensive operational experience in navigation, airborne radar, electronic countermeasures, and strategic planning. He culminated his career as Chief of the Planning Branch, with responsibility for all software used in the Strategic Air Command Control System.

Although the software-hardware cost ratio appears disproportionate now, the imbalance is forecast to increase rapidly in the years ahead as hardware gets cheaper and software (people) costs continue to go up. The Air Force conducted a comprehensive study to estimate costs of software and hardware for command and control systems through 1985. Fig. 1 shows the estimate for software expenditures in the Air Force going to over 90% of the total automatic data processing system costs by 1985.

The infancy of the software development process is emphasized by the number of computer program projects that have gone out of control, resulting in time slippages and cost overruns. Study after study has shown that the software industry has not profited from previous experience and tends to propagate the same errors in the development and management of computer programs. These points were emphasized by W. Boehm:

“The CCIP-85 study found that the problems of software productivity on medium or large projects are largely problems of management: of thorough organization, good contingency planning, thoughtful establishment of measurable project milestones, continuous monitoring on whether the milestones are properly passed, and prompt investigation and corrective action in case they are not. In the software management area, one of the major difficulties is the transfer of experience from one project to the next. For example, many of the lessons learned as far back as SAGE are often ignored in today's software developments.”

The management problem was also stated well by J. Aron at the 1969 Second NATO Conference on Software Engineering:

SOFTWARE is big business. Total costs for computer software in the United States exceed $10 billion annually, more than 1% of the gross national product. Major technological advancements have resulted in significant decreases in the cost of computer hardware, but related software costs continue to escalate rapidly. The recent World-Wide Military Command and Control System (WWCCS) computer procurement was estimated to involve expenditures of $50 to $100 million for hardware and $722 million for software. An estimate for NASA was an annual expenditure of $100 million for hardware and $200 million for software.
“We made a study of about a dozen projects, though not in a very formal manner. However, our results were convincing enough to us to set up a course on programming systems management.

“The nature of the study was ‘Why do our projects succeed or fail?’ We took as ‘successful’ a project that met its requirements on schedule within the budgeted dollars and satisfied the customer. On this basis out of 10 or 12 projects that we examined, we had one success and a whole lot of failures.

“We analyzed the reasons for failure, as given to us by the project managers, and by the people who had performed previous examination of the projects. They gave various reasons behind the failure of the projects, virtually all of which were essentially management failures. We ran into problems because we didn’t know how to manage what we had, not because we lacked the techniques themselves.”

Recognizing the problem

The Missile and Surface Radar Division (MSRD) has had a wide range of experience in real-time computer programs. As a result of this and comprehensive investigations, MSRD identified the following areas in software development which require extensive evaluation:

- Cost estimation procedures, including all checks and balances.
- Omissions in the specifications and what these entail in the end product.
- Approaches to changes in the baseline.
- Quality of the design including reliability, “good code,” clean interfaces, etc.
- User consideration in the requirements.
- The software-build process.
- Version control (multiple version consideration, language conversions).
- Testing levels.
- Performance monitoring techniques.

Effectively, MSRD recognized that the concept of programming quality/standards is critical in turning the general magic of software development into a controlled software engineering process. Therefore, a task force was formed of managers, systems engineers, and senior staff analysts with extensive experience in software system development on projects conducted by RCA, other companies, and the military services. The assignment of this task force was to design and implement a total system for managing and developing computer programs.

Birth of CPDAMS

Using the collective experience of its members and the analyzing the techniques espoused by experts in the industry, the task force established that a disciplined, total system of problem definition and resolution is needed. Conceptually, MSRD’s Computer Program Development and Management System (CPDAMS) evolved, and workshops were conducted to develop detailed specifications for the system to make it viable and applicable to the development and management of any computer program project.

CPDAMS provides a disciplined and scheduled implementation effort of a well-defined problem solution. This is achieved by:

1. A top-down approach to problem definition and resolution.
2. Organization of the system into logical development phases.
3. Definition of required inputs and outputs for each phase.
4. Establishment of organizational responsibilities and interfaces.
5. Implementation of effective management controls.
6. Incorporation of scheduled reviews and approvals.
7. Development of documentation as an integral part of the system.

The top-down approach in CPDAMS gives total project visibility by covering the spectrum from initial system requirements definition through system operation and maintenance. An integrated team of systems engineers and systems analysts is identified and given the responsibility for all system requirements activity. This group maintains continuity through subsequent computer program requirements, design, implementation, and integration phases of development. By this means, the performance and design specifications are improved significantly, thereby establishing an efficient baseline for the system. Quality of design in the initial document-production phases of development eliminates many of the system deficiencies and reduces the iterations back to design during integration and system testing. The CPDAMS chart (Fig. 2) summarizes project outputs, technical activity, project management, and reviews.

Project management tools

CPDAMS provides effective project management tools which assure orderly progression from system definition through implementation, integration,
and system testing. A brief discussion of the principal tools follows:

Adaptability to any computer program development project. CPDAMS is a standard system of computer program development and management. Therefore, provisions are made to identify unique requirements of a specific contract and make the necessary modifications to the system.

Input and output requirements are clearly established throughout the system. The hierarchical structure of documentation provides continuity from the initial top-level documentation through the detailed specifications.

Checklists have been prepared for all of the project output documents to provide guidelines for total system design. Additionally, the checklists are valuable for conducting document audits to verify completeness and during technical review to evaluate the quality of the design.

Primary and support responsibilities are specified for all development, implementation, integration, testing, review, and management activities. The delineation of responsibilities across the entire system establishes organizational interfaces, integrated activities, logical transitions, and effective channels of communication.

Management plans are used for project management, training, and computer program testing and validation. Each plan is developed incrementally to ensure that instructions, procedures, and associated documents are completed in a timely manner prior to implementation.

An Operations Library is established for the control, storage, protection, and accountability of master tapes/disk and computer program releases.

A Configuration Management Plan is developed to control all changes to the project output documents. Through this plan, all changes to the system must be approved and approved changes must be communicated through use of formal change notices.

Quality Control/Quality Assurance are established as functional requirements throughout all stages of development. Effective implementation results in assured quality from initial system design through system operation.

Flexibility to implement advancements in computer programming techniques is built into CPDAMS. MSRD personnel are scheduled to receive training in structured programming, a top-down system in which major programs are broken into small programs to provide ease of coding, testing, maintenance, and modification. Structured programming represents a new technical standard which permits better enforcement of design quality for programs. Use of this programming method in CPDAMS will contribute to the effectiveness of the computer program design, reduce coding errors during implementation, and provide better computer program releases for integration and system testing.

Project monitoring is accomplished by scheduling several types of reviews to evaluate the quality of project outputs and management plans and to effect smooth transition from one phase of development to the next.

Supporting documentation is required to provide detailed information on the full scope of CPDAMS. A series of standards has been identified to control input on project output documents, management plans, reviews, audits, and levels of testing. MSRD will develop these standards on a priority basis during the next several months.

**CPDAMS Development Phases**

With the repertoire of project management tools provided by CPDAMS, the software manager is prepared to address the total project and organize it in the following eight logical development phases:

- **Phase 1: System Requirements Definition.** A team of systems engineers and systems analysts is formed to accomplish the systems requirements activity through system synthesis, system analysis, and analysis of computing system requirements. The following project output documents are produced: System Performance Specification, Functional Flow Diagram and Description (Tiers 0,1), Operational Scenario, System Interface Document, and Computing System Specification. Production of these documents within the guidelines of CPDAMS provides a quality definition of system requirements to serve as the overall guidelines for subsequent development of computer program system requirements and design.

- **Phase 2: Computer Program System Performance Requirements.** The integrated team accomplishes computer program performance definition and computing system analysis as reflected in the project output documents: Computer Program Performance Specification, Functional Flow Diagram and Description (Tier 2), and Computer Program Interface Document. A quality definition of computer program system performance requirements is accomplished to provide the guidelines for development of the computer program design.

- **Phase 3: Computer Program System Design.** The integrated team develops the computer program architecture, determines the computing system configuration, designs modules, and prepares module-level working documents. Primary project outputs are the Computer Program Design Specification and Functional Flow Diagram and Descriptions (Tiers 3,4) which provide the guidelines for a total system approach to coding and computer program implementation.

- **Phase 4: Computer Program Implementation.** Systems analysts who were part of the integrated team during the previous three phases lead the implementation effort. Primary technical activity involves: module data base build and test, unit testing, preparation of the module build, and internal training for the programming team. The objective of this phase is to develop quality module code and data base build to produce the releases required for computer program integration and validation.

- **Phase 5: Computer Program Integration and Validation.** The systems analysts and programmers accomplish computer program functional, performance, and reliability testing of the data base. Modifications to the computer programs and data base are accomplished to effect timely resolution of problems until effective integration and validation are accomplished. The system program and data base releases required for equipment and computer program integration are produced as outputs of this phase. Internal training of the computer programming team is completed and training is provided for the program manager, quality assurance, and system integration. Development is initiated on program and user manuals required by contract. At the end of the phase, the computer program system is held off to the program manager for use by system integration during the subsequent testing phases.

- **Phase 6: Equipment and Computer Program Integration.** System integration, supported by the developing departments, accomplishes computer program functional, performance, and reliability testing using simulation and subsets of the equipment configuration as detailed in the Test and Validation Plan. Timely response to resolution of problems is provided by the developing departments to assist in the completion of effective integration. Updated computer program and data base releases are produced for system testing. Work continues on the development of program and user manuals. Internal training is completed and plans are finalized for conducting customer training.

- **Phase 7: System Testing and Acceptance.** System Integration conducts the system functional, performance, and reliability testing until validation against the System Performance Specification is achieved. Developing departments continue to provide timely response to problem resolution. The program manager, system integration, and quality assurance conduct the tests for customer acceptance of the system. Final computer program and data base releases are prepared for delivery to the customer as well as program and user manuals required by contract. Customer training is provided during this phase, as required.

- **Phase 8: System Operation and Maintenance.** This phase is primarily the responsibility of the program manager who provides operational assistance to the customer, as required, and responsive resolution of problems. Developing departments provide the changes required to resolve customer problems in accordance with established system change control procedures. Customer training may continue during this phase if required by contract.

**CPDAMS in Action**

CPDAMS is currently being used for the
development and management of the computer programs for AFAR, the advanced phased-array radar system RCA is developing for the Advanced Ballistic Missile Defense Agency of the Department of the Army. In its application to this project, the system is called the AFAR Project Management System. Modifications from the standard CPDAMS are primarily in the nomenclature of project output documents required by the customer and the extent of customer involvement in project reviews. A standard on the AFAR Project Management System has been published and distributed primarily for use of personnel involved in project development and management. Other standards which have been published and distributed are: AFAR Software Documentation Standards, AFAR Phase Reviews Standard, and AFAR Document Checklists.

MSRD management has selected AFAR as the pilot project for implementation of CPDAMS. With the experience gained, MSRD will be evaluating the system's applicability to other developmental projects. MSRD is optimistic that effective implementation of CPDAMS will result in the capability to do a better job in less time at reduced cost, thereby achieving customer recognition of RCA's ability to manage computer program developmental projects efficiently.

Acknowledgments

The authors would like to thank many members of the engineering staff at MSRD for their contributions to this project. Much of the early research work was developed and led by E. A. Behrens. L. H. Crandon added greatly to the structure and format of CPDAMS.

References


Bibliography

Chief Programmer Team: Principles and Techniques, PSC. 71-5108 (IBM Federal Systems Division, Gaithersburg, Maryland).


Fig. 2 — Summary of computer program development and management system; a formal CPDAMS chart is available from the authors. (Asterisks * show areas of customer involvement; bullets • indicate configuration management).
Computer program reliability
Dr. P. G. Anderson | L. H. Crandon

To introduce the problem, some "horror stories" in computer program development experiences are given. Then the underlying source of unreliable behavior is discussed, followed by a comparison of computer program reliability with existing reliability disciplines. Features and functions that can be incorporated into computer programs to increase reliability are shown, together with design principles that can be applied to increase reliability. Measures of and cost tradeoffs for reliability engineering are discussed and two design approaches (structured programming and proof techniques) are described in detail.

Lawrence H. Crandon, Mgr, Hardware/Software Systems, Command and Control Systems Engineering, MSRD, Moorestown, N.J., received the BS in 1945 from the City College of New York, and the ME degree in 1947 from the Polytechnic Institute of Brooklyn. He has taken additional graduate work in mathematics and computer sciences at the University of Pittsburgh and the University of California at Los Angeles. Upon graduation he joined the Westinghouse Research Laboratory, developing electronic circuitry for fm-cw radar; he was subsequently involved in system analysis of missile guidance, automatic landing of aircraft, radar data processing, and radar countermeasures at Gillilan Inc. In 1953 he joined Autonetics, Lockheed Missile and Space Division, where he designed and developed computing and data control systems for satellite projects, and participated in selection of hardware and software for satellite tracking stations. From 1959 to 1970, Mr. Crandon served in a variety of supervisory and staff positions in the Hughes Aircraft Ground Systems group, where he planned, performed, and guided study, analysis, design, development, and test/evaluation efforts in air defense and other command and control systems for both the US and foreign countries. Mr. Crandon joined RCA MSRD in 1971 in the Command and Control activity, where he has been engaged in structural studies and in directing development of real-time computer programs for control of advanced phased array radar equipment. He is a senior member of the IEEE, and a member of the ACM. He has been awarded seven patents.

Dr. Peter G. Anderson, Command and Control Systems Engineering, MSRD, Moorestown, N.J., received the BS in Mathematics in 1962 and the PhD in Mathematics (Algebraic Topology) in 1964, both from MIT. He spent a year at Princeton University as an instructor before joining RCA Computer Division's Systems Programming Department to work on the development of the Fortran system for the Spectra/70 computer line. His special interest in that and subsequent compiler development activities was object code optimization. In September 1971, Dr. Anderson joined the faculty of Newark College of Engineering as an Associate Professor in the Department of Computer Science. He is currently working with MSRD's Command and Control Systems Engineering Department in addition to his work with NCE.

Authors Crandon (left) and Anderson.

How can one specify or measure total system reliability using only one of the contributing factors—equipment reliability? The computer program, residing in a computer which is an imbedded, integral part of the system, must also take its place in the reliability equation. This remains true despite the difficulty in defining "reliability" of a software product. But with the immense—and growing—national investment in computer software, computer program reliability is being, and must continue to be, reduced to a structured methodology subject to prediction and measurement techniques with effective design principles and techniques.

Horror stories

From a report on a radar-array testing program:

"An apparent phenomena which evidenced itself during the testing of — was that of a computer program which had been declared operational, 'gradually' deteriorating. The particular programs had been in operation for several weeks, and were evidently 'sick' and getting 'sicker'."

"The sickness would demonstrate itself by the frequency of program failure or program error. With the — program, the output, under normal operation, was supposed to be the number and location of the inoperative driver modules. In the deteriorated condition of the system, the number would come out wildly high, or else the program would 'hang-up' in an intermediate routine. As the system got sicker, the frequency of 'hang-up' would increase. With a second program, errors associated with the radar peripheral interfaces, —, would demonstrate its inoperability by the higher rate of error while copying from disk to tape."

From a report by IBM Corporation on its Apollo programming support:

"...For example, during one of the missions, they found an error involving a module not releasing core store when it should. It just kept on claiming more and more core, and this was not noticed in the simulations. As a result, when the system was running, after a few days it got more and more sluggish. It was working, but non-essential functions were being dropped, because of lack of core space..."

Comments from military personnel maintaining an operational, formally accepted, large-scale computer-based defense system:

Reprint RE-19-5-21
Final manuscript received November 10, 1973.
The programs get tired, sluggish. They become full of crud. "We have to clear the computers and reload each day."

These examples show the effects and time-dependence of unreliability in computer programs in three diverse applications: a radar testing program, an Apollo mission, and a defense system. Similar examples could be found in computer programs controlling the manufacture of medicines, regulating life-support systems for hospital patients, governing airport activity, and other commercial activities.

The reliability of computer programs needs to be guaranteed; currently, it cannot even be measured. But the situation is even worse than that: sometimes it seems impossible even to get the large computer programs to perform at all. The final phase of development of a program—acceptance testing—is often concluded only by agreeing to modify the original specifications.

Source of unreliable computer program behavior

The source of unreliable behavior in computer programs is, of course, errors introduced during the initial program requirements specification, the design process, the implementation (coding), or subsequent modifications. These errors can be blamed on such things as insufﬁcient skill or training or lapses of attention by designers and programmers, and on faulty information (documentation) provided to them. But complexity is the key. Complexity can be found in all phases of computer program development. "It has reached the point where users find it easier to understand how a computer works than how an operating system works."3 The more complex a computer program, the more difficult it is to understand and to test, measure, or modify—consequently the less likely it is to be reliable.

Complexity is not a simple, one-dimensional quantity and it seems most resistant to quantitative treatment. However, complexity does depend on program size—especially the sizes of the individual program modules—and on the extent and nature of the interrelations of the individual components. Miller4 has suggested that the number of entities a human can mentally grasp at one time is about seven; if the number of constraints that a line in a program must meet to be right exceeds this number, then perhaps the program's complexity has gotten out of hand. This problem of program complexity has been discussed widely.5,6,7,8

One approach to reducing the extreme complexity is a programming discipline—popularly called structured programming—which will be discussed in some detail later in this paper.

Contrast with existing reliability disciplines

Other engineering disciplines, such as electronics or mechanics, have developed reliability methodologies that permit them to measure (and guarantee) the failure rates, mean-time-between-failures, expected life, and so forth, for their systems. They can design systems to conform to previously specified constraints and can advertise mean-time-to-failure in multiples of years.

The reliability engineers produce bathtub curves (Fig. 1a) showing the failure rates for their devices; the first part shows infant mortality, followed by constant failure rate, and finally wear-out. We are able to construct a similar curve for computer programs (Fig. 1b) showing an initially high error rate for a young, bug-ridden program, then a momentary calm for the programmers (who now think their program works rather well), followed by a surge of exposed bugs at the time the program is put into actual use, and eventual stable performance as the bugs that cause the errors are removed. There might be later humps on the curve when a bug fix has disastrous consequences for the rest of the program, when new improved versions are released, or when new customers and uses enter the picture.

Such curves have been investigated and mathematical models have been developed, but only to a very preliminary stage.5-11 The organic structure of a computer program is unlike anything else studied thus far. Computer programs are not homogeneous; typical programs, for instance, spend 95% of their time exercising 5% of their instructions.12 Error models—based on such metaphors as drawing balls from urns and replacing defective ones with good ones—may be suggestive but fall far short of representing computer programs.

Nonetheless, one can often improve a situation before specifying all its aspects rigorously. Consequently, a computer program reliability study must start with the various factors within the disciplines of computer program design, implementation, testing, and measurement that bear on the determinations and effectiveness of computer program reliability.

Computer program features and functions for increased reliability

There are many features and functions that can be incorporated into a computer program which will increase its reliability. Convenient groupings of these are discussed below.

Effective error management

Separate computer program modules called device managers are generally assigned the responsibility of controlling the operation of devices ancillary to the computer (e.g., tape devices, disks, printers). These computer program device managers normally have a formal interface with the program modules that utilize these devices, and also a formal interface with the executive program both for operational purposes and for conveying device and service failure information (see Fig. 2).

To enhance the reliability of a program these device managers should have the responsibility for detecting and reporting device failures — both total and partial — where service can still be maintained selectively. A device manager should also be equipped to monitor its own performance to detect failures such as might result from queue overloads.
The design of computer program device managers will vary within the system according to the system's equipment and functional needs; however, the basic architecture of a device manager will follow a standard design, at least for classes of devices. Some of the elements of the standard design may be eliminated or reduced by parametric adjustments for each situation. The computer program device manager can be equipped with failure-avoidance programming which acts to:
1) Bypass transient errors—achievable by retrying the device, recopying data on failed memory, etc.
2) Bypass permanent errors—achievable by reallocating storage to bypass bad storage areas (e.g., on disk, tape).
3) Continuing service at degraded performance levels—e.g., by formally rejecting additional user requests until queues are emptied to working levels.
4) Discontinuing operation of non-essential failed services by formally rejecting user requests.
5) On-line error repair—tolerating errors by utilizing redundancy, error correcting codes, averaging (damping/filtering) methods.

Automatic error logging

The notion of device managers that monitor the performance of hardware resources and (possibly) their own performance can be extended to many other system resources. Some possibilities are data bases and their handlers, math subroutine libraries, and schedulers. For large systems, all these monitoring facilities can feed into a standard error manager (see Fig. 2) which would maintain a program-error data base. This data base would contain the error descriptions furnished it by the resource managers as well as items such as time-of-occurrence and user. Such records would then give the data necessary to develop reliability figures for the system as well as aiding program maintenance work.

Program self-protection

Errors should be isolated so that a failed program module does not damage another module or its data sets. A list of special considerations is shown below.

1) The highest priority is safety. Programs affecting safety and security should be designed for the highest reliability. These include a) the minimum executive (kernel) program which should be designed for control of the computer when the rest of the executive program fails and b) the input/output programs (kernel) which manage equipment requiring safe shutdown (e.g., missile and manager program).
2) Use memory protection devices to protect instructions from being modified, to keep instructions from being read, data from being executed (except in very special cases), and one module from affecting the resources belonging to other modules except along the lines of specified interfaces. A specific suggestion for achieving this latter protection (in addition to hardware "locks" that prevent inadvertent modifications) is to design interfaces between data bases and user modules along functional lines rather than direct accessing. This means that information is input to, and extracted from, a data base via a single interface routine, where the interface routine is the only routine privileged to access the data directly. This yields a benefit of flexibility; the data base structures can be thoroughly revised without affecting the user programs.
3) Use program self-diagnosis tools, such as check sums, memory checks, and evaluations of assertions of the form used in proving or explaining programs (see the section below on proof procedures for some examples).
4) Use alert/abort and recovery procedures so that calls for scheduled maintenance or for immediate error diagnosis can be handled without complete or unsafe termination of service.

Design principles for increased reliability

In addition to the effect of program architecture, the way a large program is put together affects reliability. Design principles to increase reliability fall into three major categories: 1) error-management documentation, 2) tools for error-free implementation, and 3) error-management and program-change procedures. These are discussed in the following paragraphs.

Error-management documentation

Any large-scale computer program should be considered from the viewpoint of its error-management structure and capabilities, and should be systematically described so that a comprehensive view of error handling may be available in a single document. This document should be made suitable for assessing the computing system's error-detection capability, error vulnerability, recovery/reconfiguration potentials, and overall reliability.

Tools for error-free implementation

The following principles lead to a reduction in errors of program origin:

1) The less the programmer has to do toward the end of program development the less likely is error to occur; the fewer instructions, design, and operating procedures the better. Also, the fewer languages he needs and the more syntax-error-tolerant languages (with strong editing and debug capabilities, etc.) the better.
2) Existing well-documented, time-tested programs should be used, and modified for program parameters.
3) The design should take into account requirements for easy and complete testing.
4) The designer should use such tools as data dictionaries, data directories, and data preparation systems; standardized interface formats; and expressive programming...
languages.  

5) Modular structures conducive to reliability should be used. Modular design can minimize testing requirements when program modules are short (simple), with well-defined functions to perform. When programs are partitioned to match modules to functions, then loss of a module minimizes the number of lost system functions. If modules are designed independent of each other, then loss of one does not bring others down. Modularly designed programs also allow modification over a time period in order to minimize errors due specifically to design structure, over a time period, thereby improving computer program reliability over time.

6) Modules should be designed so that their modifiable parts (generally data sets) and non-modifiable parts (instructions and static data) are segregated. This will facilitate application of reliability enhancing measures.

Error management and program change procedure

Computer program error management and reliability must be maintained throughout program change procedures. The principal danger is that program changes to a "tested" computer system may negate previous tests. This can occur if a change alters the design or the previously tested computer program execution paths and data structures. Compilers and other program generators must be examined to ascertain the extent to which a tested program has had its earlier tests invalidated by compilation with either single-parameter changes or extensive program changes. When patches are made at the lowest language level, and recompilation is not involved, methods should be used to assure that previously conducted tests are not invalidated.

Reliability engineering

How does one measure computer program reliability and what influences cost tradeoffs? The testing procedures and the performance monitoring for "working" computer programs should be designed not only to detect bugs to be repaired but also to measure the program reliability. In particular, the errors should be categorized according to their severity, their origin (i.e., environment at time of occurrence), the modules/functions to which they pertain, and what was done about them (fixed or symptomatically removed). This log should yield, by extrapolation along curves of reliability modules, some quantifications for program reliability. Where program modules exhibit a network structure, methods similar to those provided by Kim, Case, and Ghare can be applied to compute system reliability. Where the structure of modules is hierarchical, other methods need to be considered.

A particularly suggestive notion is that of using program exposure as indicative of reliability:

\[
\text{Reliability} \geq \frac{(\text{use of program})(\text{possible use})}{(\text{exposure})}
\]

where the \text{use of program} is some weighting of the tests the program has passed, along with the successful user experience. The ratio, \text{exposure}, is based on some estimate of what the program hasn't yet been required to do. A naive approach to this might take exposure as the fraction of program instructions that have been exercised. Better, but much more complicated measures can be envisioned, but of course "combinatorial explosion" does take effect.

Factors in reliability-cost tradeoff are online error management programs vs. computer use, program testing, computer program maintenance service costs, and equipment vs. computer program alternatives.

Two approaches to computer program reliability

Two of the techniques mentioned earlier are rather new and promising: structured programming and program proof techniques. These are discussed below in terms of their potential for meaningful enhancement of program reliability.

Structured programming

Structured programming is a discipline for proceeding in a systematic way from a program specification to a computer program. Structured programming is a process of successive refinements (a top-down process). First, the overall program is specified in "broad brush" style; that is, the programmer specifies what is to be accomplished and in what order, but he leaves the detailed methods unspecified. This "broad brush" description is written as a computer program; it's not just a narrative description of a program.

Each component of the program is either an operation that the computer can perform directly, such as \text{add} 1 to \text{x}, or the name of a more complicated operation that will have to be specified in greater detail, such as \text{delete stale entries from radar track table}. Operations of the second type are programs that have to be written. These are written the same way the main program was written. This process continues repetitively until there are no processes whose details are left unspecified, at which time the program is written.

The building blocks for structured programming are shown in Fig. 3. Notice that each of these three constructions is surrounded by a dotted rectangle with one entering arrow and one exiting arrow. These larger rectangles can be used as the components of larger sequences, alternatives, or loops.

Conventional programming languages have the facility to specify these three constructions, and to indicate steps to be specified later as subroutine calls or macro's for which there are library maintenance routines and linkage editors in the operating system.

Structured programming has often been referred to as "go-to-free" coding, but that is not entirely accurate. The undisciplined use of \text{go to} statements does increase the complexity of a computer program and consequently decreases the reliability. However, Fortran does not have the \text{if-then-else} construction nor the
begin-end and do-end nesting structures of PL/1 and Algol; but the same program control can be coded in Fortran and even of PL/I and Algol; but the same program statements in a disciplined manner. Additional stylistic techniques such as comments, spacing, indentations, and so on, then yield very readable, understandable programs.\textsuperscript{30-33} Other language features—if they are available—are permitted within structured programs. Such features would include the full PL/1 do-loop facility, a case statement (see Fig. 4) or decision tables.

Experience indicates that the necessary tools for expressing algorithms are available within the rules of structured programming. Bohm and Jacopini proved that every "ordinary" computer program can be rewritten as a structured program.\textsuperscript{34} More to the point for the individual programmer is the observation that structured programming restrictions are not felt to be confining. When one has an overwhelming urge to construct a jump into a loop or a leg of an if, then it is time to reconsider the design of the algorithm; one has probably fallen into the trap of "bottom-up" coding which results in designing a program after detailed decisions have been made.

After learning to use structured programming, the programmer quickly finds that, far from being a strait jacket, this discipline provides him the means of organizing and managing the mass of details in his job. For instance, he can modularize the program so that each program module can easily be held down to a single printed page (i.e., under 50 lines of text) or even smaller to keep the complexities from becoming overpowering.

**Proof techniques**

One answer to the issue of guaranteeing the correctness of a computer program is to prove, as a mathematical theorem, that the program is correct. Programs constructed according to the discipline of structured programming are especially amenable to this technique. In structured programming, one expresses a large process in terms of several smaller processes—subroutines, macro's, etc. To prove that the large process is correct, one takes as "lemmas" that the component processes are correct, then applies rules of inference similar to those described in Fig. 5 to conclude that the larger process is correct. The verification that the component processes are correct may proceed along the same lines, or they may be obviously correct or otherwise inspire confidence.\textsuperscript{35,36}

The formal statement of what a program is supposed to do is expressed as an implication: "if a certain relationship holds among the variables when the process starts, then the process will terminate, and when it does, another specified relationship will hold among the variables." Such relationships or assertions about the variables can be interleaved among the instructions or components of a program or attached to the arrows of the program flowchart. Fig. 5 shows some of the ways assertions can be attached to flowcharts and indicates some elementary rules of inference. Fig. 6 shows how these can fit together to prove that a program calculates what it is supposed to.

The given proof does not include the verification that the program eventually terminates; that is left as an exercise for the reader.

As in classical mathematics, there are techniques developed for proving that programs with certain patterns of construction are correct.\textsuperscript{37-51} For instance our example relies on a key *invariance*

\[ ra = 1 - b(1-a) \]

The details of such proofs may seem to be the same kind as those that cause programmer errors in the first place, but consider these two points: assertions and inferences provide the perspective the programmer needs to jog him out of his mental lapses, and eventually we expect computer programs to take over the mass of details involved in the verification process.\textsuperscript{52,53}

**Conclusion**

Unlike classical reliability engineers, we do not know how to model or measure satisfactorily the failure mechanisms in computer programs. However, we can follow their lead, noting where programs fail and tracing the chain of "why?" back to the source. The advances in electronics reliability came with learning how to simplify construction and work in a "clean room." Top-down programming is our simplification and the three construc-
Quantification will be an elusive solution for a long time. Much that goes wrong is attributable to human failure where more than size—is the overwhelming room."

References

Impact of hardware/software tradeoffs in a command and control system

J. C. Kulp

Within the last ten years, computers have become a critical link in large command and control systems. For example, in modern radar systems the phased array is totally controlled by computers, whereas the older mechanical radars closed the search and track loops by hardware. In the tradeoff decisions, however, hardware advantages still receive greater emphasis. This is because the software practitioner remains the "new boy on the block." He finds himself explaining his software interface problems in terms of hardware terminology, usually without a one-to-one correspondence. Likewise, the senior system engineer's background has more often been in hardware than in software, with resulting hardware orientation in the tradeoff decisions. One intent here is to stress the software viewpoint. This paper is divided into two sections. The first discusses areas of systems that are ideal candidates for hardware/software tradeoffs, and provides basic rules governing the decision process. The second section discusses hardware/software interface design approaches and their impact on the software design and system development.

Today's command and control systems are constantly evolving to meet new and changing requirements. This evolutionary process, quite naturally, exerts substantial influence on hardware/software tradeoff decisions, particularly as current hardware technology becomes outdated and designers turn more to software modifications as the most cost-effective method of making system changes. Unquestionably, a major design requirement of any command and control system evolving today must be ease and reliability of modification.

Moreover, the basis for tradeoff decisions between hardware and software will continue to change as computer capabilities increase, processing speeds multiply, and multiprocessing technology becomes more widespread. Some of the more complicated functions normally associated with equipment will be available as part of the computer (e.g., FFT, coordinate conversions). Hence, evolution in computer capabilities as well as in hardware technology will continue to alter the outcome of hardware/software tradeoffs.

Identifying hardware/software tradeoff areas

A prerequisite for identifying hardware/software tradeoff areas is a clear understanding of the system objectives and functional flow. From this point, an initial functional design can be developed as the basis for deriving the software elements for a tradeoff decision: estimates of bulk data storage, functional execution requirements (both mathematical and logical), execution rates, data flow rates, data block sizes, and initial definition of the operator interface requirements.

With an initial functional layout in hand, it is practical to identify areas that are candidates for hardware/software tradeoff studies. These areas can be classified in five categories:

1) Bulk data handling with fast processing cycles
2) Independent algorithms with fast completion cycles
3) Incompatible processing loops
4) Real-time (ms to μs) control algorithms
5) Mathematical models with complex computations

Bulk data handling with fast processing cycle

Systems with large data bases and many elements that must be processed within a short period of time are candidates for hardware/software tradeoffs, because a total software implementation approach could easily consume an excessive amount of the computer execution time. An example is the identification of known satellites in space track systems. Satellite positions are calculated for specific times using their orbital elements, and these elements must be correlated with those of unrecognized detections within milliseconds to determine if an object is a new or a known satellite. This process of deriving correlation candidates for several thousand satellites each minute of the day is a bulk-data-handling problem. Performing a correlation with each of these candidates for a single search detection (which occurs several times per second) represents a fast processing cycle.

Another example is the problem of keeping track of railroad rolling stock location in a fast transportation system. The track system must be organized in several thousand sections, and each train or car in the system must be associated with a given section. This bulk data base is used...
as the basis for the software to control traffic and ensure a "fail safe" system. When a train enters a given segment, its physical location is fed into the computer to update the car/track location file. Because the location detector (hardware) may fail, alternative methods in the software are used to predict the actual location, such as last reported car speed. The car/track location file must be updated many times a second to provide a "failsafe" system. Hence, there is a large data file of track sections vs car locations that must be updated at a high rate (processing cycle).

These types of algorithms are candidates for hardware/software tradeoffs because of the large storage requirement and the percent of computer execution time required to perform the algorithms. A partial hardware implementation approach becomes feasible when these algorithms are characterized by a well-defined set of input and output parameters and the decision logic has few special conditional requirements. Special conditional decisions or mode changes are areas that are modified during system integration and system evolution. Hence, the decision process should be kept in the computer and the decision passed to the hardware to maintain system flexibility.

In the first example, the satellite candidates could be placed in an auxiliary storage unit along with special processing logic built that would accept search detection parameters and the correlation cell size. The hardware logic could sequence through the candidates and correlate each candidate with the search detection. Those candidates that are within the specified correlation cell size would be identified and sent to the computer. Upon completion, the computer would have a small number of potential satellites detected by the search function. From here, the computer could perform finer correlation logic to resolve any further conflicts.

The total software solution to the same example is to organize the data into multilayers of various parameters (e.g., time, x, y, z, v, etc.). The search detection would be categorized to a specified multilayer area which would reduce the number of candidates requiring fine correlation processing. The computer time required to build the multilayer file would require continuous updating as satellites move, taking a large percentage of computer time.

The advantage of hardware implementation is that 1) the satellite candidate file structure can be simple, requiring less computer time to build the file; and 2) the computer can start the correlation process and go about some other task until the results are fed back from the hardware.

The disadvantage of the hardware approach is that the computer time to drive the hardware (i.e., set up the file in the auxiliary memory, prepare the inputs, start the hardware, safe-store data before starting another task, retrieve the safe-stored data to continue the correlation processing) may add up to be the same or more execution time than if the entire correlation process were performed in the computer.

Hence: The computer execution time required to set up, control, and effectively use a hardware-implemented algorithm may equal or exceed the computer time required if totally performed in the computer.

The algorithm design with hardware is entirely different than the design with software only. If the tradeoff is made using the hardware design for both hardware and software implementation, the hardware would be the best approach. The same could be said for the software implementation approach.

Hence: The hardware/software tradeoff must compare the hardware design algorithm with the software design algorithm and not with an algorithm as implemented by hardware, then software.

Independent algorithms with fast completion cycles

Algorithms that are independent of other software algorithms and have a fast completion rate are also candidates for hardware/software tradeoffs. Again, these algorithms, if implemented partially or entirely by hardware, would relieve the computer of an execution burden.

An example is the fast transportation system described above. The car/train location file is the main data base for the entire software system; many algorithms access this file. One of these algorithms is the file update algorithm. It is special in that while it is performing an update, the file is in a transient state and other algorithms cannot be allowed access. Also, it is closely integrated with the "fail safe" check algorithm since the car location and car/track status information is packed and sent together by the hardware to the computer whenever a car enters a new track section. The integrated algorithms would check for dangerous situations and immediately enter the "fail safe" recovery algorithm before completing the update. Hence, the file update algorithm checks for the "fail safe" situations and updates the car/track file.

The software implementation approach

James C. Kulp, Mgr. Real-Time Computing Systems, Command and Control, Systems Engineering, MSRD, Moorestown, N.J. received the BS in Mathematics in 1958 from Kansas State University and the MS in Operations Research and Management Science in 1970 from Johns Hopkins University. Mr. Kulp's early experience with U.S. Army Signal Corps from 1958 to 1961 was primarily in the area of design and development and functional specification of special purpose computer systems. In 1961 he joined the Bendix Corporation where he became involved in development and testing of the radar driver program for the AN/FPS-85 radar data processor. From 1969 to 1971 he was manager of the programming group at Bendix, responsible for technical software improvements on the AN/FPS-85 systems as well as technical direction of new software design for operational programs and analysis of special radar algorithms. He joined MSRD in June 1971 in the Command and Control section where he became involved in subcontract management of the design, development, and testing of the AEGIS AN/SPY-1 radar control program effort. He was responsible for development of the AEGIS radar simulator and a member of the team that integrated the AEGIS system in the Moorestown Land Based Test Site in 1973. Mr. Kulp currently directs the real-time programming group within the Command and Control Systems Engineering organization, with overall responsibility for all MSRD radar control computer programs.
would be to develop an _update preparation_ algorithm integrated with the "fail safe" _check_ algorithm. The _update preparation_ algorithm would update one section of the track at a time. It would read the section into a work area, update the work area, then notify an _update_ algorithm to complete the update. The _update_ algorithm would use a one-command transfer (storage-to-storage) to update the car/track file section. This would eliminate overhead in having to place algorithms in the wait status while the file was in a transient state.

A partial hardware implementation approach would be to build special-purpose equipment that would perform the car/track file-update by storing the actual car location and status information directly in memory. The advantage would be that the computer need not use execution time in processing interrupts, accepting car/status data, then storing them in the car/track file. The disadvantage would be that the update would not be independent of the "fail safe" _check_ algorithm. By the time the computer sequenced through the car/track file, identified an error and took some "fail safe" recovery action, any saving in computer time would be eliminated.

Hence: _a clean, independent algorithm lends itself to hardware/software tradeoffs_. Algorithms which have functional interactions with other software algorithms become questionable for implementation by hardware.

Incompatible processing loops

One area of software design that consumes storage and time is that involving various algorithm loops that merge but have different or incompatible cycling criteria. For example, in phased-array radar space-track systems, there are two basic processing loops. One is based on time (the satellite correlation candidates file); the satellite-data-collection algorithm is performed based on a time loop. In this case the "satellites in coverage" files are generated every 24 hours, satellite candidates are computed every minute for 30 minutes in advance of real time, and the active satellite candidate file is updated once a minute.

The second basic processing loop is event-oriented (i.e., based on the track rate). When a target is scheduled for a track point, the software track loop schedules a radar transmission, receives and processes the return, updates the track list, then waits for the next time to schedule a track point. Note the loop is triggered by time, but the software starts processing when data are available.

Mode switching in the track loop is a critical software design area where the hardware/software interface with the radar has a marked influence. Mode switching requires additional data gathering in order to calculate or derive all the necessary parameters for the new mode. Hence, in mode switching there is usually a significant transient period where unique logic paths are required in the software to prepare to enter the new mode.

Hence: When defining the interface between the hardware and software, if each mode is considered separately, the transient processing loops become complex, resulting in increased software timing and integration testing problems. If the hardware/software interface is defined with respect to the basic functions, the software design will be able to implement a simple mode change design and be less sensitive to modes and mode switches.

Real-time control algorithms

Real-time control algorithms that control activities down to the μs level are excellent candidates for hardware/software tradeoff.

An example is the interface between a phased-array radar and the computer. The computer's radar control loop normally cycles once every 20 to 50 ms. It prepares a radar control list which is interpreted and executed by the radar interface control (RIC) unit. The purpose of the RIC unit is to synchronize the various radar units, both digital and analog, on a μs-to-ts timing basis that results in a radar transmission or return.

This interface has been made at various levels of timing control. The computer can establish timing controls that are to be executed on a μs-to-ts basis but cannot implement the synchronization since the computer executes too slowly. Normally, real-time radar control algorithms consist of three loops.

1) The time loop at which the software cycles,
2) The time loop at which the radar schedules transmit and receive (called the PRF),
3) The time loop at which a target is tracked.

The software-cycle loop is sometimes based on the PRF cycle, permitting one common block format that the computer generates for each pulse repetition in the radar interface control unit.

The PRF size is usually from 50 to 20 ms and is based on hardware considerations. As the PRF becomes shorter, the software timing loop is increasingly critical and becomes a major influence in the software design.

Another approach in choosing the software-cycle loop is to base it on the minimum track rate. A typical minimum track rate is from 50 μs to seconds. This approach minimizes the software timing problems and requires a different type of hardware/software interface. For example, the interface buffer containing the radar control list would be larger and would cover more than one PRF. Having the software cycle based on track rate permits the software to vary the rate according to the specific targets under track. Therefore, the software design would have an additional degree of design flexibility. One could go farther and let the software control the PRF size by adding it as a function of the radar control list. This would provide the software with still another degree of design flexibility with respect to timing.

Hence: _The hardware/software tradeoff interface of a real-time control system must consider the basis for the software cycle to minimize software timing problems and provide needed design flexibility._

Mathematical models with complex computations

Mathematical models with complex computations are yet another class of candidates for hardware/software tradeoffs. Complex computations such as integration, solution to differential equations, and FFTs take a large percent of computer time and storage.

Approaches to keeping complex mathematical computations in the software include converting the computations to curve-fit or table lookups, changing data spacing from exact to nominal, etc. If these computation modifications affect the results of the model, another approach is to perform...
There are two types of interface flow and the software has to be designed faces or software-controlled interfaces. A methodology: hardware-controlled interface will require the software design to stop and restart hardware interfaces should be controlled by the software. Also, if the interface design can stop and restart the hardware, then it can start it with nearly the same logic.

Hence: The design of a hardware/software interface must consider how the system is started, stopped, and restarted. Software control of hardware/software interfaces minimizes software timing problems.

A hardware/software interface design has two functional interfaces: data/control flow and status flow. These two functional interfaces may have different timing requirements. Data/control flow usually has a constant rate, as in a radar interface. The status flow does not necessarily have a constant flow rate; it may be activated only when a change in status is generated. The difference between data/control-flow errors and status-flow errors is that data/control-flow errors are solved by retries, whereas status-flow errors require some reinitialization and restart actions which usually terminate and restart the algorithm.

Design questions arise, such as:
1) How is "real time" entered and distributed throughout the system?
2) What does "real time" mean and how is it used in various parts of the system? (Real time means different things to different parts of the system.)
3) How is "real time" checked for errors, and if an error is detected how is "real time" recovered?

Failing to resolve the system design problem of implementing the real-time clock may result in an excessive number of clocks being used throughout the system. Later in the integration phase, the problem becomes a complex one of keeping them all in synchronization. This is usually a very costly situation, and sometimes becomes literally impossible.

Implementation of functional or mathematical models in hardware or software to ensure a required accuracy can become a critical hardware/software interface design area. In some software systems, implementation of a complex, exact mathematical model may take an excessive amount of time and storage. The model could be modified by converting an integral computation to a curve fit or collecting data requiring nominal spacing rather than exact spacing. Such modifications could reduce the accuracy or change the model to a point where it becomes ineffective. In this case, a solution is to perform a hardware/software tradeoff to provide for performing some of the more complex mathematical computations in the hardware, with the results fed to the computer.

Conclusion

Hardware/software tradeoff decisions affect every part of a system design, including the software. The careful, informed system designer can identify key system areas as candidates for hardware/software tradeoff implementation decisions, and in this way enhance system performance and cost effectiveness.

The hardware/software interface design can change a good hardware/software tradeoff decision into a bad one. A number of specific interface design areas must be considered, since they have a direct impact on integrating and delivering the system on time, on schedule, and at reasonable cost.
Computer selection for command and control systems

H.J. Hurtado | Dr. S.A. Steele

In early command and control systems, a computer was often selected somewhat arbitrarily and then the system designed around the hardware. Today, however, systems have grown and customer needs have become more clearly defined. Moreover, existing tactical environments often dictate severe limitations on the physical size of the computer, thereby further complicating the selection process. This paper discusses the procedures for selecting a computer to satisfy the various real-time requirements of a command and control system. The evaluation-and-selection process is described, with some detail devoted to the use of benchmarks. The basic considerations include hardware and software architecture, languages, support software, simulation requirements, I/O capability, and use of existing software — and how these considerations influence the cost/performance ratio. The emphasis is on a single or basic computer system in a multi-programming environment as opposed to a distributed and/or multi-processing system. Minicomputers are not addressed, since their selection is another topic.

At the outset, the pacing requirements on any computer hardware configuration must be identified. Then, as the system is designed, these pacing requirements can be updated and the impact of changes on the computer hardware can be evaluated. Fig. 1 summarizes the computer evaluation and selection process. Each block is discussed briefly below, and salient considerations are highlighted. Each of these blocks contains certain steps that are prerequisites to an intelligent and cost-effective computer selection.

Requirement analysis

In establishing system specifications, an overall functional flow must be developed that specifies what the system must do. In addition, a detailed typical operational scenario must be generated showing how the computing system must interact with the man and with other hardware to make the total system operate properly. System constraints and data rates must also be identified.

With this system information, it is then possible to start defining the software requirements. This includes software functional flow, hardware/software partitioning, and a functional break-out of the software modules. A load model must be developed to be used in estimating the software real-time processing load during worst-case, average, and nominal situations. Where time permits, computer simulation is recommended. The model must account for all concurrent high-priority real-time tasks, including resource scheduling, special I/O processing, system-error recovery, as well as all low-priority processing tasks.

With the load model completed, it becomes possible to estimate the size of each computer-program functional module in terms of core memory and real-time utilization. Core estimates should be based on code written in assembly

---

Henry J. Hurtado, Ldr. Advanced Systems Design, Command and Control Systems Engineering, MSRD, Moorestown, N. J., received the BA from Fordham University in 1954 and MBA from New York University in 1960. He has eighteen years of professional experience in real-time computer systems, including switching and communication systems, the Polaris/Posiedon project, Air Force command and control systems, and management information systems. He has served as consultant to corporate management in the areas of feasibility studies, computer selection, and software management. Since joining RCA in 1971, he has had responsibility during the system definition phase for the computer selection and software architectural design of AFAR, a computer-controlled phased-array radar for ABMADA. He was responsible for RCA's software systems engineering concepts utilized in an Air Force study program for a multi-computer complex to control remotely piloted vehicles (RPVs) during tactical missions. Presently he is managing a software development subcontract related to the Integrated Radio Room for the Trident submarine.

Reprint RE-19-5-19
Final manuscript received January 23, 1974.
language. The total number of instructions, local data, global data and execution rates should be tabulated. From such tables, percent of real times during average load and peak load can be generated for a given instruction mix and specified computer times. Thorough and careful analysis is required to develop a representative instruction mix for a particular application.

With the proper instruction mix and the basic machine execution rates specified, processor speed can be estimated in number of instructions per second. For a particular application, the “effective” number of instructions per second is much lower than the number of operations per second quoted by computer manufacturers; compiler inefficiencies for memory and timing must be added to the basic estimates. With this information, a basic computer configuration may be put together. It is important to keep track of the margin in both time and memory in order to maintain good design practice. The effect of inadequate design margin can be seen in Fig. 2., taken from the CCIP-85 report.\(^1\)

Unfortunately, in many hardware procurements, decisions have been made as if the folklore curve were true. Typically, after a software job is sized, the hardware is procured with only about 15% capacity over that determined by the sizing, presenting the software developers with a machine that is 85% saturated at the outset.

Other major factors that must be considered are the word length, cycle time, core size, and features such as mass memory, real-time clock, memory protect, direct memory-access channels, peak transfer rate, hardware floating-point arithmetic, power fail safe, adequate availability, and MTBF. With this information specified and assessed, a preliminary computer configuration is generated. Evaluation factors structured as in Fig. 3 are used to analyze system performance and cost tradeoffs and to assess the risks. Each factor is scored (from 0 to 10) and weighted before it is combined with other factors in developing a total score for each candidate computer system. An existing time-sharing program can then be used to evaluate the overall score in terms of its sensitivity to each factor. Another step in evaluating software architecture is the development of an overall timing diagram showing critical timing loops.

**Benchmark utilization and formulation**

Benchmarks provide an indication of how well the vendor's computer performs on the required instruction mix. The system designer should provide the benchmark (or instruction mix) to the computer vendor and insist on its use. A vendor's benchmark should not be used since it does not reflect the required instruction mix; as a result, there would be no uniform way of evaluating various computers.

The benchmark provides a standard measure which reflects the application to a specific set of requirements. In a radar real-time system, for example, the benchmark may be how fast a Kalman filter executes to meet the system requirements of tracking \(n\) targets per second. In a computer network, the emphasis may be how fast and how many on-line terminals the computer can poll. In essence, the system designer must identify the critical items of his system.

It should be emphasized that although a preliminary screening of possible vendors narrows the eligibility list, an analysis of vendor specifications without benchmark data is insufficient. There are several reasons:

1) One vendor’s specifications are not always analytically comparable to those of another vendor.
2) Today's third-generation hardware architecture is highly complex, and the selection of various hardware options can radically affect the computer’s performance. Examples of options that affect performance include additional high-speed general-purpose registers, multipoint memories, mass-storage units with faster access, increased core memory, or a hardware black box for a time-critical function (e.g., a fast fourier transform).
3) The high interrelationship between the hardware and operating-system software can affect the net throughput rate of a computer system differently, depending on the application.

Simplicity of benchmark design will save many man-weeks of evaluation. The questions to be asked are:

- What does one do with the benchmark results?
- What do the results indicate?

If the answers are not obvious, the

![Fig. 2 — How do hardware constraints affect software productivity?](image-url)

![Fig. 3 — Evaluation factors.](image-url)
benchmark is not properly designed.

Among the several technical factors essential for an adequate benchmark, four major considerations affect the computer configuration:

1) Input/output throughput: the rate at which the computer channels can get data in and out of the machine. Is the data rate so high that cycle stealing is inadequate and requires a multipport memory?
2) Processor logic and arithmetic: the computational capability and range of processing. For example, is floating-point double-precision required? Can software routines substitute for hardware multiply/divide? Is emulation required?
3) Data storage: the type and volume of data storage. Is random access required or can magnetic storage be adequate?
4) Operating-system software: an integrating benchmark is not properly designed.

Unless the system designer plans to develop a specially tailored operating system for an end use, the benchmark should include operating-system performance. In past real-time systems, it has been necessary for computer users to write parts of their operating system. In these situations, a system designer may consider modifying the vendor's operating system rather than developing a new one. The development of a new operating system is expensive and generally requires a long-lead implementation schedule.

A product of any benchmark is evaluation of the compiler. An inefficient compiler may cost excessive computer time and additional core storage. A compiler must be evaluated both from a development and production standpoint: that is, compile time and execute time.

Two benchmark models have been useful in computer evaluation; these models address the technical factors and application considerations mentioned above.

Benchmark 1

The first model (Benchmark 1) provides a standard program that can be used by all computer vendors. It is written in a compiler language (FORTRAN or COBOL) and addresses the highly critical function(s), which may be a worst-case situation or a representative example of a major processing load. In one radar system, for example, this program was a fully coupled Kalman filter, written in FORTRAN. This model provides the following advantages:

1) It is written in FORTRAN, giving all vendors a standard program.
2) It demonstrates the compiler capabilities and whether the vendor's FORTRAN compiler is at a specified standard level.
3) It outlines the computer configuration necessary to support that program.
4) It provides core estimates and compilation times.
5) It allows a demonstration of debugging aids and their capability in software development.
6) It provides core estimates and running time in executing the program.

Benchmark 2

The second model (Benchmark 2) is a program specification that permits the vendor to demonstrate the capability of his computer hardware and the overhead associated with the operating system. The model has two elements: the first is a basic processor task (no I/O) and the second is the execution of the same basic processor task under increasing rates of I/O and interrupt activity.

The basic processor task(s) may be to exercise certain highly used instructions (e.g., shift instructions, move instructions, bit/byte manipulation, double-precision). The processor task should be written in the language that will ultimately be used. In real-time processing, the task might be written in assembly language.

The processor task may be short but can be executed repetitively and long enough to get an accurate time measurement. For example, a routine that can execute in 600 µs is too short if the internal clock has only a millisecond resolution. However, if it is executed 1000 times, then the measurement of 600 ms ± clock variance is meaningful. Timing should be provided for each task in the case where several tasks are used as the processing baseline.

The second part of this model performs the same task(s) but is interrupted by I/O and hardware/software interrupts. There are several stages, where each stage runs the processor tasks but interrupts at a certain rate (e.g., every 100ms, 50ms and 1ms). This allows the user to see the amount of operating system overhead required to service I/O and interrupt requests. As the rate increases, a good operating system will show no appreciable degradation.

Variations of this model can be developed to measure the operating effectiveness of sequential files, random files, and index-sequential files. The model can also be set up to show the effect of multiple I/O requests to the same data base.

The advantages of this second model are as follows:

1) The vendor can write a program task utilizing the best features of his system.
2) The user is given definite time statistics on system overhead.
3) The user can see whether cycle-stealing for I/O can severely affect processing time.
4) The points where interrupts can degrade the system are highlighted.

Comparative analysis

It is interesting to note the credibility of these benchmarks compared with the vendor's specification and other benchmarks which had the same candidate computers.

The first comparative evaluation used for radar application is on three processors supplied by different computer vendors. Table I shows some of the basic computer characteristics of three machines. At first glance, processor A looks more powerful; however, the benchmark results of the two models shown in Table II do not substantiate first impressions. Benchmark 1 which was the FORTRAN Kalman filter (basically a number crunching routine) proved processor C to be twice as fast as A. In benchmark 2, processor B showed its low operating system overhead as the interrupt rate increased. In fact, the 1-µs interrupt caused processor C processing time to rapidly increase, and processor A failed to provide any timing data.

Comparing these results against other benchmark studies, we find somewhat similar results among the same family line of computers. The Philco-Ford Study 3 by J. Corsiglia also shows processor B as the most cost-effective machine over processor A and processor C.

RFP preparation and evaluation

Referring to the process outlined in Fig. 1, with the basic configuration and benchmarks defined, it is now possible to screen and identify various computer candidates and to develop specific
The conference, which all interested computer vendors are invited to attend. This provides a vehicle for identifying response schedules and a procedure for During the period between release of the proposal stage.

proposals, software definition and detail specification continue. In certain cases, it addenda to all vendors, where such information affects the configuration that they will propose. Information dissemination must be equitable to all vendors.

Vendor response to the RFP should contain a technical proposal, benchmark results, a cost proposal, and contract information relating to delivery, maintenance, support, and relevant past experience. It is important to assess the level to which the computer manufacturer assumes responsibility for his computer hardware and software items (component, subsystem, or system).

The vendor's technical proposal should identify the hardware configuration proposed (and options offered), deliverable software, and the level of software support. It may be necessary to clarify or modify the vendor's configuration because of cost tradeoffs, and this process may be iterative.

From the standpoint of benchmark results, the following questions should be answered:

1) Were the benchmarks performed on the same configuration as the one proposed? If not, are there substantial differences? How can these differences be reconciled?
2) Were all listings and test data output verified to ensure that the benchmark was actually performed according to user specifications?
3) Were the test results consistent with vendor specification? If not, why not? Can the vendor provide an adequate explanation?
4) Are the test results consistent with previous benchmarks performed by other users? If not, was the user benchmark addressing certain characteristics that the other benchmarks did not possess or stress?
5) Does any vendor have to rerun the benchmark either because he did not follow the specification or wants the addition of some feature in the computer hardware that could make a substantial difference?

When the vendor material has been reviewed, each candidate system should be evaluated against the criteria established before the release of the RFP. This evaluation starts with a tabulation of total hardware and software costs. Other considerations such as maintenance and reliability of equipment also must be weighted in the evaluation. Finally, the candidate systems are ranked; management review is conducted; and the top two or three candidates are selected for the final round. Best and final offers are then obtained from each selected vendor.

Conclusion

In almost all command and control systems, the selection of the basic computer system has been plagued with problems. In some cases, excellent hardware has been selected but the operating system and support software have been inadequate. In other cases, the computer has been undersized for the job. To avoid such problems, it is important to specify completely all the requirements that must be met before the selection process and to develop a checklist to be used in the final selection and evaluation. Accurate and ample design margins for time and memory are important considerations. Intelligent applications of benchmarks to evaluate the actual computing power of the hardware (with appropriate software) can produce meaningful information which often differs from that implied in computer vendor specifications. In summary, the computer hardware must be evaluated with the operating system, application, and support software to select the best computing system in terms of a cost/performance ratio.

References

Software verification and validation for command and control systems

A. R. Chandler

Software — the computer program subsystem — is the heart of a command and control system, in that software determines the performance of the system. Consequently, the quality of the verification and validation of command and control system computer programs during system development is a major determinant of the quality of the delivered system. Proper provision for the necessary testing, measurement, and evaluation techniques from the earliest planning and design stages greatly reduces the cost while increasing the effectiveness of software verification and validation throughout the development period.

When a military command acquires a command and control system, it does so to meet an operational requirement—it buys a capability. Thus, what the system does is more important than what it is. This fact, plus the pivotal role that software plays in determining how a total command and control system operates, makes software verification and validation extremely important.

Although some imprecision exists in the application of the terms verification and validation individually, collectively they refer to those activities that ensure the software operates as specified and meets its design goals. Validation generally is used more broadly than verification, and alludes to meeting performance or mission objectives. Verification addresses the manner in which the development is executed:

Do the detail design specifications satisfy the higher-level performance and development specifications from which they were derived?

Does the software, as built, physically and functionally match its detailed design?

The subject of software verification and validation merits far more attention than it generally receives in publications covering computer programming techniques, and more discussion than space permits here. The reader is referred to a book Program Test Methods containing an excellent collection of papers presented at a symposium held in June 1972 at the University of North Carolina. Not only is this book informative on a number of aspects of the subject, but it also provides a comprehensive bibliography on related subjects. Having thus referred the reader to a broad range of topics related to software testing, this paper now focuses on a single overriding theme: that software verification and validation must begin with the earliest phases of computer program design and must be an integral part of every aspect of the design and development activity.

There is a high payoff in concentrating from the beginning of a system project on designing for testability. Typically, nearly half of the total effort expended in the development of software for a large system—from preliminary program design through system integration—is spent in verification and validation, not including module debugging (i.e., unit testing). Experience also shows that the poorer the system definition and design, the larger will be the effort expended in verification and validation to uncover and correct design weaknesses. Moreover, since the correction of problems and errors (including the testing of fixes) is less expensive earlier in the development process than later, when more parts of the system have been integrated, provision for the early identification of errors must be made. Clearly, considerations of software verification and validation must enter the command and control system picture when the software performance and design requirements are first documented.

Documentation

The specification of the software components of a command and control system.
modules and integrating the major software elements, including at least a portion of the total system integration process. Sometimes this plan is called the "build" plan.

The test plan provides the basis for the achievement of the milestones of the build plan, in the context of the software quality assurance process. This plan cannot be implemented prior to software implementation, since the design specifications have been completed. In most command and control systems, this is the "executive" program, responsible for keeping track of what is going on, what devices are in use or available for use, and what modules are scheduled for operation. There are also service routines, controlled by software for attached equipment and for performing standard mathematical and data-management functions.

Each of these many components is designed to accept certain specified inputs—variables and parameters—and to generate certain outputs. These inputs and outputs may be passed directly between modules or may be held in buffers or data files for later processing. It is here that verification and validation are applied to the test and verification processes, built in a controlled and updated as a part of that process.

As the development progresses and the elements of the software system are designed, the test planning can be accomplished. It is not possible to anticipate unit test requirements, since the performance and design specifications do not provide detailed module (unit) design data. However, since it is essential that the test process proceed in a controlled fashion, assurance must be given that the fashion, assurance must be given that the software modules have been completely verified prior to their aggregation into larger packages.

Software design

The design of command and control systems is a process of organizing the systems is a process of organizing the system hierarchy of functions that the system is to perform in satisfying the prescribed operational requirements. From the major functions, subfunctions are derived which together constitute the higher-level connections, common tasks, etc., being defined at each level.

The level of this functional analysis at which the tasks are allocated to the prime elements of the system—equipment, software, people—is the basis for the software performance specifications. These in turn become the basis for the design requirements and the design itself.

Before coding of the software can begin, an effective software structure must have been derived, and the detailed design of each software element (program modules, data files, special subroutines, executive programs) must have been accomplished. In most command and control system programs, there is an "executive" program, responsible for keeping track of what is going on, what devices are in use or available for use, and what modules are scheduled for operation. There are also service routines, shared by many modules, for controlling attached equipment and for performing standard mathematical and data-management functions.

Test planning

Planning for software verification and validation is intimately related to development planning, and proceeds in parallel with it. The software development plan covers the entire process from coding and debugging the modules or "units" through building packages of
major processing flow of his module, calling on each next lower-level task as needed, but by name only, without coding it. He then codes these next lower-level tasks in the same manner, and each lower level in succession until the lowest-level procedure has been coded. All procedures coded by this technique are self-contained, receiving prescribed inputs when called and returning control to the point in the next higher-level procedure from which they were called, delivering the prescribed results. Clearly each procedure can be separately defined and individually tested in terms of its prescribed inputs and outputs.

Software “build” process

The foregoing discussion stresses the point that the design process, from command and control system functional allocation to software module design (and including test planning), proceeds from the top level down—from the general to the specific. On the other hand, the implementation process, including verification and validation, generally proceeds in the opposite direction, aggregating system components until the entire system is built. Fig. 1 illustrates this composite “top-down/bottom-up” process.

The software-build process includes integrating, verifying, and validating ever-increasing sets (builds) of modules, testing each against subsets of the command and control system functions as prescribed by the test plans. It is an orderly process, designed to detect, isolate, and diagnose problems easily, by adding new pieces and functions to each existing build only after the build has been thoroughly tested. Thus whenever an error is uncovered, it can reasonably be presumed to have been introduced by the new addition. This presumption, of course, depends on the completeness of the test plans, the effectiveness of the module testing (unit testing, debugging), and the proper separation of functions during initial design, including precise definitions of module inputs and outputs.

As suggested by Fig. 2, the build process can begin before the individual modules have been completely coded and debugged. The build plan, together with the system test plan, schedules the coding process so that the procedures required to build the functions are verified first (the hatched areas on the figure). Later, on completion of module coding and debugging, the complete module is integrated in place of the partial one. The build schedule is thus telescoped with the programming schedule, shortening the overall software implementation period. Moreover, the programmer gets additional assistance in testing his modules, and the test teams get more experience operating the software system than if they waited for module completion before beginning integration.

Generally, the very first functions to be integrated in the build process are those required by the software design requirements rather than the overall command and control system performance requirements. These include such features as the executive system, inputs from and outputs to computer system peripheral devices (display consoles, tape units, etc.). Other requirements in addition to these capabilities are programs supporting the build process itself, and verification and validation. These primarily involve the generation, recording, and processing of data items, control tables, and files that provide module interconnection and operational control.

The next capabilities to be added are the major functional loops or threads which represent “front-to-back” system operation of the simplest functions in the simplest operating mode. In terms of module inputs and outputs, these consist of a selected set of data linking a basic build of related modules. Such modules might range from a “raw” external data input from a sensor, for example, to a system-generated command to an effector, as the active external system elements are sometimes called. Some examples include the receipt of track data to the issuance of missile guidance commands; entry of flight plan data to establishment of track identification; processing of raw radar data to compute track position; and conversion of raw and telemetry data to compressed and filtered reports. All of these processes could be part of the same system and thereby tested in parallel, in separate builds. These builds would themselves be integrated and linked through their shared data, executive program, and resource utilization.

The “hardware/software” integration proceeds in parallel with this software module-to-module integration process. The links between equipment and computer program elements are very much like those between modules: Items of data, coded and formatted according to specifications, are sent to the computer by external equipment for software interpretation and use, and then sent by the computer to the external equipment to control it, as directed by the software. The verification and validation of this process, then, is very much like the verification and validation of software builds, in that data is passed between system elements and constitutes the information used to verify that the elements and their operation is as specified.

Tools and techniques

Many tools and techniques exist for
carrying out software verification and validation, at all levels. The implementation and application of these tools and techniques depend on early planning and design, including allocation of adequate funds for their implementation and planning for their application within the software design itself. Some of the key tools and techniques which have proved useful are described briefly below:

1) **Data matrix**—the data matrix is a simple layout of tests (stimulus-response pairs) in a tabular form for ready and complete presentation of predicted results and check-off of satisfaction of tests. It portrays input data items on the left side, in columns, and required outputs on the right. Rows, or sets of rows, contain test values for each data item, and constitute individual tests. The data matrix is particularly effective in module debugging, verification, and qualification for the integration or build stage.

2) **Traps**—trapping is a technique in which the logic flow of a program is interrupted for the purpose of setting aside interim results for test measurement or for performing special test functions—computations, file modifications, file printouts, scripted actions, recording, etc.

3) **Data extraction and recording**—software "instrumentation" can be implemented and installed as part of the software system for recording test data for verification and validation of software builds, up to and including the entire command and control system. Such instrumentation can be extended into the operational phase to support such functions as recording the operational history data required by some users, or extracting diagnostic data for system maintenance.

4) **Data report generation**—the format for data files, tables, and messages is described in terms of scaling, units of measurement, mnemonic codes, *etc.*, and the layout of these values in binary computer "words." Programs are coded to handle this data as formatted, but the binary equivalent is difficult for the programmer to read. A data report generator can reconvert data into the external forms—units, mnemonic codes, *etc.*—simply by knowing which file, table, or message it is processing.

5) **Data reduction and analysis software**—at the validation stage, data extracted during system test runs must be subjected to statistical analysis. Noise, trends, tolerances, and other performance statistics are more of a factor than individual values yielded at a particular point in time. This statistical software is used after test runs to process data recorded during those runs.

6) **Simulators**—simulation is a technique used for modeling the performance of a device not yet delivered for system integration, simulating only its software interfaces, modeling the environment, simulating only the processing time of each module against sample loads, *etc.* Simulation techniques provide test repeatability, a quality not generally realizable with "live" data.

7) **Test drivers, scripts, data generators**—to run tests in a controlled manner, especially system tests where a situation develops over a long time, it is necessary to work within the framework of a "scenario"—a description of a dynamic situation. For this it is necessary to load the input data files for the system with data values representing the test situation or events to yield recorded data to evaluate against expected results. The tools for this permit the relatively easy generation of data in external form to be entered automatically into the system at the proper time.

These are but a few of the available tools and techniques for the verification and validation of software for a command and control system, but enough to illustrate the point that these tools are as important to software implementation as the design, development, and acquisition of tools, jigs, fixtures, and instruments are to equipment development and fabrication.

**Conclusions**

The command and control system is built and tested as a hierarchy of elements—from the most basic procedure or hardware component to the entire system. Each element or aggregation of elements making up the system has a prescribed job of converting a specified set of input data to a set of specified results under certain conditions. At each step of the building process, the verification and validation activity consists of setting up representative and critical conditions and data inputs, predicting the results (as part of test planning), and comparing the results of actual operation of the elements against the predictions.

**References**

Simulation of a command and control data processing and display system

C. S. Warren

System simulation can be used to verify that a proposed hardware/software configuration for a command and control system will have the needed throughput and response times at peak load. Recently, a system simulator was developed to be used in determining the throughput and response-time capabilities of a data processing and display system for a proposed Tactical Air Control Center Automation. The simulator specifically addresses the hardware and software of the data processing subsystem. It shows the utilization of various hardware and software components of the data processing subsystem and keeps track of the lengths of all the subsystem queues, it exhibits the critical paths in the flow of data. The simulator also provides sensitivity tests to such parameters as CPU speed, assignment of priority, and various software features, thus allowing the system architect to evaluate alternative subsystem designs. The simulator, therefore, a tool for use in system optimization as well as a means of performance verification.

THE SYSTEM being modeled consists of a central processor, a mass memory subsystem, a controller for all communication and display interfaces, the resident portions of the TDOS-IPS/70 software, an operational data base, and application software. The TDOS-IPS/70 software is a data-base management software system that was offered commercially by RCA Computer Systems. The simulation model is driven by input data supplied by the user in the form of operational scenarios. Each operational event in a scenario is subdivided into transactions, one transaction for each operator action or input message. The individual transactions, scheduled in accordance with the scenario (allowing for human reaction-time to responses from the system), are the actual input to the simulator.

The outputs of the simulator consist of reports of the response time to each event in the scenario, statistical reports on response times to individual transactions, and statistical reports of the utilization of hardware and software components of the data processing and display (DP&D) system. All queues in the subsystem are monitored and their length statistics are reported. The statistics, which are reported periodically throughout the simulation, consist of the mean, the standard deviation, the maximum value, and the minimum value of each variable.

The simulator is coded in Fortran and is based on the general-purpose simulation program—GASP II. Most of the critical system parameters are provided as input data, and so can be varied without recompilation. The use of Fortran allows structural changes in the model (to investigate potential structural simplification in the DP&D system) with manageable effort.

Hardware

The hardware elements included in the model are shown in Fig. 1. The central processor unit (CPU) is characterized by its speed (instructions/sec) and memory capacity. It could be any processor with simultaneous compute and input/output (I/O) capability. The I/O devices modeled include the mass memory and interface controller. The mass memory was assumed to interface the CPU thru a dual-channel controller permitting access to two records simultaneously. Drums were assumed in the model, but only modest modification would permit multiple disk files in place of the drums. The interface controller is the interface for all communication and display subsystems. The unit record peripherals (printers, card readers, etc.) normally associated...
with a data-base management system were not included since such equipments represent an insignificant load on the DP&D system.

The communications and display subsystems serve only as sources and destinations of the messages which define the beginnings and endings of activities in the scenario. They are represented by a message queue which feeds the DP&D system through the interface controller. Five display processors and two communications processors are identified as sources and destinations. The display hardware furnishes message-size parameters for display input and output messages.

Software being simulated

The software represented in the model consists of the resident portion of TCOS, Version 2, and the applications software appropriate to the scenario. The software in the communications and display subsystems and in the interface controller is not modeled in detail.

Operating system

The operating system modeled (TCOS) is a direct evolution from commercial software known as TDOS and IPS-70, which are standard operating system products that run on the Spectra 70 computer family. TCOS operates as a multi-tasking system driven by transactions generated in the communications and display subsystems and by internally generated (periodic) transactions. It provides scheduling and resource-management facilities for a large number of online users, with data-management facilities enabling these users to share a data base.

In addition to the inter-computer link (ICL), TCOS contains the transaction control program (TCP), the data control program (DCP), and the executive control system (ECS). This latter system provides general executive services, including handling, resource management, physical and logical 1/O control, and console control. It is tailored by system generation to the particular configuration required.

TCOS is a transaction-oriented system in which data-processing activity is performed by transaction user programs (TUPs) activated by input messages from local or remote terminals. The transaction control program (TCP) is that module of TCOS which examines messages coming into the system, recognizes legal transaction codes in message headers, identifies the correct transaction user program for processing, delivers input messages to transactions and accepts their output, and supervises the termination of transaction user programs and links their successors. The TCP also schedules periodic, self-generated transactions for such purposes as self-test and supports a system administrator who receives performance monitoring information and controls error recovery and system reconfiguration through a command language.

The remaining module of TCOS, the data control program (DCP), controls and facilitates access to the common data base. Because the DCP maintains a single, comprehensive directory of all files in the system, file descriptions are created by the system administrator; TUPs need not repeat these descriptions in order to use the files. Data integrity and privacy are assured by controlling access privileges for both reading and writing. Sequential, indexed-sequential, and direct-access methods are available, along with a hierarchical data-base language which can be used with any of these access methods. The direct access method is serially re-entrant and can queue several physical access requests to the data base at one time.

Of the six executive control positions normally available to applications software when TCOS is used in the batch mode, TCP and DCP each occupy one. The inter-computer link (ICL) uses a system control position. The remaining four user positions are available for applications software in the on-line mode used operationally in the system and represented by the model. They may be occupied by TUPs, which are scheduled by TCP and are the normal vehicle for applications processing in an on-line environment; alternatively the console operator may use one or more positions for "local" programs (batch programs which run directly under ECS and for which TCOS may provide services). TCP has two own-code options available for additional applications software; one of these options allows input messages to be pre-processed, altered, or rejected before they generate transactions, while the other provides for post-processing of output messages before they are transmitted to the communications or display subsystem.

In addition to its normal function of providing an interface between TCP and the communications and display sub-
systems, the intercomputer link (ICL) carries out functional processing in connection with track-message reception and frag-order dissemination.

Prominent software modules

The software modules most significant in the operation of the simulation model are the Loader, the TCP, the TUP, the ICL, and the executive scheduler. Other functions of the ECS are accounted for by overheads added to the processing times of the programs supported by the ECS; the execution time of the DCP is similarly treated.

Loader: the Loader module handles the relocatable loading of the transaction user programs (TUPs) when the transaction control program (TCP) converts incoming messages into transactions using a terminal directory list prepared by the system administrator. It determines which TUP in the program library should be the initial TUP to resolve the transaction request. If an ECS control position (usually known as a slot) is available, the TCP assigns it at once and immediately initiates the Loader. Otherwise the transaction is placed on the transaction queue for processing when resources become available. In either case, the TCP stores the message on the primary mass-storage medium for later retrieval by the TCP or its successors. The TCP also processes the program-to-program linkage in the chain of TUPs used to resolve a single transaction and forwards output messages to the communications and display subsystems.

Transaction control program: the Transaction Control Program (TCP) converts incoming messages into transactions using a terminal directory list prepared by the system administrator. It determines which TUP in the program library should be the initial TUP to resolve the transaction request. If an ECS control position (usually known as a slot) is available, the TCP assigns it at once and immediately initiates the Loader. Otherwise the transaction is placed on the transaction queue for processing when resources become available. In either case, the TCP stores the message on the primary mass-storage medium for later retrieval by the TCP or its successors. The TCP also processes the program-to-program linkage in the chain of TUPs used to resolve a single transaction and forwards output messages to the communications and display subsystems.

Transaction user programs: transaction user programs (TUPs) are not simulated in detail. Their sole function in the model is to provide load for the system. Each TUP is represented by the number of instructions it executes, its memory requirements, its input-output activity, and its initiating and terminating messages.

Scheduler: the executive control system (ECS) Scheduler selects, from among the TUPs and system functions already allocated to slots and loaded into memory, which program or function is next to have control of the CPU. A slot may be in a wait condition: waiting for I/O termination, waiting for a program to be loaded, not in use, active, and ready to run, or active and running. When an interrupt has been processed by ECS, the Scheduler makes a slot selection (to allocate the CPU) until the next interrupt. Among those which are ready to run, perhaps including the one which was running at the time of the interrupt. An active slot that has been running may be deactivated if a higher-priority slot becomes ready as a result of processing the interrupt; it may also be deactivated if it was put into a wait status by the interrupt processing. A low-priority slot is activated when all higher-priority slots are either unused or waiting for I/O or ECS services.

Inter-computer link: the inter-computer link (ICL) manages buffers in the CPU which provide temporary storage for data exchanged between the CPU and the communications and display subsystems; it also performs functional processing on track messages and frag-order dissemination. In the model, the ICL is represented explicitly for track message processing and implicitly for buffer management. Frag-order dissemination does not occur during the period of time selected for the scenario.

Simulation model

The simulator is an events-type simulation program running in conjunction with a simulator support program. The simulator support program is an RCA-enhanced version of GASPII, a general-purpose simulation aid suitable for use on all types of problems in which response time, utilization of system resources, and queueing effects are of paramount importance.

System representation

The system being modeled is represented in a GASPI simulation by events, processes, and entities. The fundamental elements in the simulation are the events, each of which represents a change in the state of the system model. An event may be generated externally; such an event is an input to the simulation and is the means by which the simulated environment affects the simulated system. In the DP&D simulator, typical external events are the arrival of messages in the communications and display subsystems and artificial events used to create interim and final reports of the progress of the simulation. Internal events represent the results of simulated functions being performed by the simulated system. In the DP&D simulator, the termination interrupt occurring at the end of mass-memory access or the supervisor call that initiates loading of the successor to a TUP that has finished processing are typical examples of internal events.

The entities in the simulation are the model representations of the hardware and software modules of the system being simulated. In the DP&D simulator, typical entities are the CPU (hardware) and the ECS slot occupied by the TCP (software).

Processes are performed by entities and occur between events. The parameters of entities determine the durations of the processes, which in turn determine the scheduling of internally generated events. For example, the primary mass memory parameters (latency and data rate) determine the duration of the process of satisfying an I/O request; the simulator therefore schedules the corresponding termination interrupt (an internal event) at the proper time after the I/O request was made. As a result of the termination interrupt, another process (the execution of instructions by the program that initiated the I/O request) begins. A sequence of such processes and events, beginning with an external event (the arrival of an input message) simulates the progress of a transaction through the system and the delivery of the resulting output message to a display or communication line.

GASP maintains a set of files for the use of the simulation program. For example, the event file stores all events, both external and internal.

Events are removed from the event file in the sequence in which they are scheduled, and the removal of each event advances simulated time to the scheduled time for the event removed. Other files are used in the simulator for system queues, a typical example being the transaction queue, which is an exact representation of the transaction queue actually maintained by TCOS.

Simulation program structure

The simulation program running in conjunction with GASP consists of event routines and process routines. Event routines define the logical structure of the model. Process routines combine the parameters of the entities (simulated representations of the system hardware and software modules) and the previously scheduled events (e.g., length of an input
message) to derive the durations of processes.

As each event is removed from the event file, GASP calls the events subroutine of the simulation program and furnishes it with the attributes of the particular event removed from the file. The events subroutine then calls the event routine appropriate to the type of event in question. This routine defines the response of the simulated system to the specific event.

The event routine executed as a result of the removal of an event from the event file sets up the parameters to be interpreted by the process routine to be executed after the event routine is complete. In the DP&D simulator, a process routine may be set up by an event routine, but its execution may be delayed by the scheduler. This procedure represents the behavior of TCOS which may, for example, load a TUP and set it ready to execute but may then schedule a program in a higher priority slot for immediate execution.

GASP services

In addition to removing events from the event file and calling the simulation program, GASP provides commands for manipulating its files and maintains two kinds of statistics at the direction of the simulation program. On command GASP collects the current value of any variable designated by the simulation program and maintains the mean, standard deviation, maximum, minimum, and number of occurrences for that variable; up to 30 distinct variables may be maintained. Histograms may also be prepared on command from the simulation program. Through a separate command, GASP collects and maintains time-generated statistics, for which the mean and standard deviation are time-averages rather than averages over the number of occurrences, for variables designated by the simulation program. GASP also automatically maintains time-generated statistics for the length of each of its internal files. All these statistics, plus the current content of all files, are printed on command.

Elements of the Model

The specific events, processes, entities, and queues maintained by the simulator are as follows:

External events

The external events are:

1) The arrival of a message in a communications processor or display controller; unless it is a track-data message, this type of event initiates a transaction.
2) A periodic event simulating the real-time clock of the RCA 200, such as the initiation of the process of updating situation displays from stored track data.
3) A command for a GASP report of statistics and file contents.
4) A command for the end of a simulation run.

Supervisor calls

The following supervisor calls to the TCOS executive control system are modeled as internal events, generated by the simulation program:

1) Request for mass-memory access.
2) Initiation of output message transmission.
3) Request by a TUP to load a successor TUP for the same transaction.
4) Termination of the last TUP of a transaction.
5) End of processing for a track message.
6) End of processing by the Loader.
7) End of processing by the TCP, when a transaction has been generated from an input message.

Termination interrupts

The following termination interrupts are modeled as internal events:

1) Completion of access to the mass memory.
2) Completion of transmission of an input message from a communications processor or display controller through the interface controller to the CPU.
3) Completion of transmission of an output message from the CPU through the interface controller to a communications processor or display controller.

Computational processes

Computational processes are modeled by calculating the number of instruction executions from software parameters—either parameters of TCOS modules obtained from analysis of TCOS or TUP parameters forming part of the input data to the simulation—and then by multiplying by the average execution time per instruction of the CPU. A supervisor call is scheduled to occur at the end of the calculated time. The computational processes modeled are as follows:

Input-output processes

Input-output processes are modeled by calculating the time required for each I/O service from the parameters of the I/O device and the parameters of the record being transferred. A termination interrupt is scheduled to occur at the end of the calculated time. The input-output processes modeled are as follows:

1) Primary mass-memory access.
2) Input message transmission.
3) Output message transmission.

Scheduler

The Scheduler of the simulation model incorporates the logic of the TCOS executive exit module, so that the priorities of the simulated processes are the same as in the system being simulated. The computation time consumed by the TCOS in allocating resources to the various computational processes is included with the simulated time for each of the processes.

Entities

Utilization statistics are collected on the following hardware and software entities:

1) The CPU.
2) The primary mass-memory channels to the CPU.
3) The two-way channel between the interface controller and the CPU.
4) All ECS slots available to TUPs, collectively.
5) The main memory, in terms of the number of bytes occupied by TUPs as a function of time.
6) TUPs waiting for update access to locked files.
7) Each of the ECS slots, individually—namely, the slots for the Loader, the ICL, the TCP, the DCP, the TUPs (four slots), and the idle condition; the DCP slot is not actually modeled and so always appears with zero utilization.
8) The number of data-base files not locked.
9) The number of updating TUPs active and running.
10) The number of transactions whose input messages have been received and which are still in process.
Queues

The simulator maintains the following queues:

1) **PCU message queue**: a model of the dynamic input queue maintained in each communications processor and display controller. In the Simulator program these dynamic queues are maintained together in a single GASP file, but they are treated independently.

2) **Program loader queue**: a queue of programs which have been assigned system resources and are waiting for the Loader to be available.

3) **Mass-memory queue**: a queue of I/O requests to the primary mass memory; this queue models the availability of any location in the primary mass memory to any program through either non-busy channel.

4) **TADIL-B process queue**: a queue of track messages received by the CPU and awaiting the availability of the ICL module of TCOS to update the resident track data file.

5) **Transaction queue**: a model of the transaction queue maintained by the TCOS transaction messages for which the correct initial TUP has been identified but for which system resources are not available. The transaction queue is ranked by the priority of the transactions.

6) **TCP queue**: a queue of processing requests to the transaction control program; this queue contains requests to transfer input messages from the message buffers to the primary mass memory and to generate the appropriate transactions and requests that initiate the transmission of output messages.

**Operation of the simulator**

The flow of transactions through the queues and servers are shown in the simulator block diagram of Fig. 2.

**Tracing a transaction**

The history of a simulated transaction begins with the arrival of a transaction input message in a communications processor or display controller. The procedure for scheduling the arrival of input messages is described later in this paragraph. At the scheduled arrival time, the simulator places the input message on the PCU message queue; since this event takes place in the communications processor or display controller in the real system, the simulator models it with zero elapsed time for the main data-processing system. If the model buffer is available, that corresponds to the display controller or communications processor in which the message arrived, the simulator begins transmitting the message to the CPU; otherwise transmission is initiated when the buffer is available, unless the PCU message queue contains an earlier message that arrived in the same PCU.

The arrival of an input message in the CPU is scheduled on the basis of the length of the message and the activity level of the interface controller. The total data rate of the interface controller, a variable parameter of the model, is divided equally among the messages currently in transit to compute the effective number of bytes per second being transmitted for each message. One message at a time can be transmitted for each display controller and communications processor; this message can be either input or output.

The arrival of an input message at the CPU creates a simulated terminal interrupt requesting the service of the TCP to put the message on the primary mass-storage medium and to schedule a transaction. If the TCP is available, the simulator sets up the parameters of the message and marks the TCP as ready. The message parameters are its identity (corresponding to "message type" in the real system) and are used to determine 1) the sequence of TUPs required to resolve the generated transaction, 2) its length, 3) the fact that it is an input message, and 4) the identity of the processor or controller from which it came. If the TCP is not available, the interrupt is enqueued, along with the message parameters; the message remains in the buffer and the buffer is not available for other input or output message traffic. At the conclusion of termination interrupt processing, the simulator calls the scheduler, which assigns the CPU resource to the highest priority processing routine that is ready to run.

The simulated TCP processing module is third in priority, behind the Loader and the ICL, and is activated by the Scheduler.
whenever the Loader and the ICL are both “not ready”. The TCP transfers the input message from the buffer to the mass memory and interprets the message type into a transaction. It identifies the first TUP required to resolve the transaction, and attempts to assign one of the four ECS slots designated for TUPs. If a slot is available, the parameters of the TUP are set up to simulate TUP execution, and the TCP attempts to initiate loading. If the Loader is available, the necessary TUP parameters are set up for the Loader slot, and the Loader slot is designated as ready to run. If the Loader is not available, the load request is enqueued. If the attempt to find an ECS slot for the TUP was unsuccessful, the availability of the Loader is tested. Instead the transaction is enqueued on the transaction queue, from which it will later be dequeued in priority sequence when a TUP slot becomes available.

These actions are modeled by a sequence of I/O simulations interspersed with computation simulations. When the process is complete, the simulator searches for additional work for the TCP on the TCP queue, dequeues a TCP request if the queue is not empty, and sets up the TCP slot.

The Loader simulates loading of a TUP by scheduling a sufficient number of mass-memory accesses to retrieve the amount of code specified in the TUP description, which is part of the input to the Simulator. The loading process is broken down so that each mass-memory access corresponds to the amount of code that could be retrieved in a single access in the real system. Separate accesses are included for the program modifier blocks. When the Loader terminates, it sets the TUP slot ready to run.

The TUP process is the same for all TUPs and all four TUP slots. The TUP specifications contain the number of file updates, the number of file retrievals, the number of instructions to be executed, and a specification of the output message, if any. The number of mass-memory accesses per file update is a variable parameter, as is the number of accesses per file retrieval. The total number of instruction executions is divided by the total number of file accesses, so as to make equal parcels of instruction executions interspersed with mass-memory activity. Additional mass-memory accesses are scheduled for output messages, plus an access to retrieve the input message. The process concludes with a simulated supervisor call to request termination of the TUP, either to load a successor or to designate the completion of the transaction.

The transaction specifications contained in the input data to the simulator determine whether a particular TUP has a successor or is the last in the sequence required to process a transaction. If there is a successor, the simulator identifies it from the input data and either sets up the Loader to simulate loading it or enqueues it if the Loader is busy. The successor is assigned to the same ECS slot as the terminating TUP. If the terminating TUP completes the resolution of a transaction, the simulator sets the TUP slot free and attempts to assign a new transaction from the transaction queue. If the transaction queue is empty, the TUP slot remains available.

If a TUP has an output message, the simulated supervisor call that occurs at TUP termination attempts to assign the buffer corresponding to the destination communications processor or display controller. If the buffer is not available, the message is enqueued on the TCP queue. If the buffer is available, a mass-memory action to move the message from the mass memory to the buffer is simulated; then the message is sent on its way to its destination, simulated by an increase in the interface-controller activity level. A termination interrupt, scheduled in accordance with the specified length of the message and the activity level of the interface controller, indicates that the transmission of the message is complete. The buffer is declared free, the interface controller activity level is adjusted, and the simulator seeks work for the buffer on the PCU message queue and the TCP queue.

If a terminating TUP is the final TUP of a transaction, the simulator collects statistics on its response time. Transactions are grouped into “threads”, each of which represents the behavior of one display terminal operator or communication line in the input scenario. If the terminating transaction is not the last one of its thread, the simulator retrieves the specifications of the next transaction from the input data. One of these input parameters for each transaction (except the first of a thread) is the “think time” of a human operator; the simulator schedules the arrival of the input message of the next transaction of the thread at the conclusion of the think time. If the terminating transaction is the last one of its thread, the simulator calculates the total simulated elapsed time of the entire thread, including think times, and compares it with the time allowed in the DP&D system specification. If the last TUP of a transaction has an output message, the computation of response times and the generation of the next transaction take place after the output message has been transmitted.

### Tracing a track message

Individual track messages are not included in the input data because of their high frequency of occurrence. Instead, a statistical distribution of track message length and arrival rate is determined from variable parameters of the model. Each time a simulated track message arrives in a communications processor, the simulator generates another track message from the distribution function and schedules its arrival in the event file.

The transmission of track messages from the communications processor to the CPU is simulated in exactly the same way as the transmission of other input messages. When a track message arrives, it is identified and the simulator turns it over to the simulated ICL module of TCOS. The ICL module is a pure computational process, updating the memory-resident track data file directly from the message as it appears in the input buffer. The process of scheduling the ICL module is similar to the process for scheduling the TCP. When the simulated ICL processing is complete, the buffer is declared empty and work for it is sought on the PCU message queue and the TCP queue. The simulator maintains an inter-computer link (ICL) queue for track messages that find the ICL module busy. When ICL processing is complete, the simulator collects statistics on the response time.

### Initialization

The input to the simulator, consisting of thread specifications, transaction specifications, TUP specifications, and variable parameters of the model, is presented on punched cards. At the beginning of each simulation run, the simulator reads all the input cards and stores the input data in memory. The initial conditions of the model are established—for instance, the CPU is
## Sample simulator output:

<table>
<thead>
<tr>
<th>Code</th>
<th>Mean</th>
<th>SD.</th>
<th>Max.</th>
<th>Obs.</th>
<th>Total Time</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0.0068</td>
<td>0.0136</td>
<td>0.0045</td>
<td>0.0262</td>
<td>1.5516</td>
</tr>
<tr>
<td>1</td>
<td>2.0320</td>
<td>2.0490</td>
<td>1.2960</td>
<td>1.4960</td>
<td>27.7982</td>
</tr>
<tr>
<td>2</td>
<td>6.0878</td>
<td>6.9828</td>
<td>4.4680</td>
<td>5.0813</td>
<td>83.8015</td>
</tr>
<tr>
<td>3</td>
<td>5.0878</td>
<td>5.9828</td>
<td>4.4680</td>
<td>5.0813</td>
<td>83.8015</td>
</tr>
<tr>
<td>4</td>
<td>7.0878</td>
<td>7.9828</td>
<td>5.4680</td>
<td>5.0813</td>
<td>83.8015</td>
</tr>
<tr>
<td>5</td>
<td>8.0878</td>
<td>8.9828</td>
<td>6.4680</td>
<td>5.0813</td>
<td>83.8015</td>
</tr>
<tr>
<td>6</td>
<td>9.0878</td>
<td>9.9828</td>
<td>7.4680</td>
<td>5.0813</td>
<td>83.8015</td>
</tr>
<tr>
<td>7</td>
<td>10.0878</td>
<td>10.9828</td>
<td>8.4680</td>
<td>5.0813</td>
<td>83.8015</td>
</tr>
<tr>
<td>8</td>
<td>11.0878</td>
<td>11.9828</td>
<td>9.4680</td>
<td>5.0813</td>
<td>83.8015</td>
</tr>
<tr>
<td>9</td>
<td>12.0878</td>
<td>12.9828</td>
<td>10.4680</td>
<td>5.0813</td>
<td>83.8015</td>
</tr>
</tbody>
</table>

## Impact generated data**

<table>
<thead>
<tr>
<th>Code</th>
<th>Mean</th>
<th>SD.</th>
<th>Max.</th>
<th>Total Time</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0.0068</td>
<td>0.0136</td>
<td>0.0045</td>
<td>1.5516</td>
</tr>
<tr>
<td>1</td>
<td>2.0320</td>
<td>2.0490</td>
<td>1.2960</td>
<td>27.7982</td>
</tr>
<tr>
<td>2</td>
<td>6.0878</td>
<td>6.9828</td>
<td>4.4680</td>
<td>83.8015</td>
</tr>
<tr>
<td>3</td>
<td>5.0878</td>
<td>5.9828</td>
<td>4.4680</td>
<td>83.8015</td>
</tr>
<tr>
<td>4</td>
<td>7.0878</td>
<td>7.9828</td>
<td>4.4680</td>
<td>83.8015</td>
</tr>
<tr>
<td>5</td>
<td>8.0878</td>
<td>8.9828</td>
<td>4.4680</td>
<td>83.8015</td>
</tr>
<tr>
<td>6</td>
<td>9.0878</td>
<td>9.9828</td>
<td>4.4680</td>
<td>83.8015</td>
</tr>
<tr>
<td>7</td>
<td>10.0878</td>
<td>10.9828</td>
<td>4.4680</td>
<td>83.8015</td>
</tr>
<tr>
<td>8</td>
<td>11.0878</td>
<td>11.9828</td>
<td>4.4680</td>
<td>83.8015</td>
</tr>
<tr>
<td>9</td>
<td>12.0878</td>
<td>12.9828</td>
<td>4.4680</td>
<td>83.8015</td>
</tr>
</tbody>
</table>

## Summary report.

A sample output from a typical run is shown in Fig. 3. These outputs are called summary reports and can be specified at any interval or at any specified time in the simulation. The summary report can be a full report or a truncated report depending on the needs of the user. The full summary report consists of five parts. These are:

1. Input parameters
2. Response-time statistics
3. Utilization statistics
4. Histogram data on operator response times
5. Status and statistics on all internal queues

Each summary report consists of a heading showing date, run number, users name, and the simulated time of the summary report.

Several of the response-time and utilization statistics are for validating the consistency of the input data and serve no useful purpose in evaluating the system being simulated.

At the completion of each run, the simulator prints out a final summary report identical in format to the interim reports plus a thread-response-time report (Fig. 4) showing the arrival time, completion time, required response, and actual response for each thread specified in the input scenario.

Sometimes it is desirable to examine certain of the statistical data over a portion of the simulation run. For this reason, the user is given the option of accumulating the statistics over the entire run or to have the statistics reset with each summary report so each consecutive report is for the previous period only. This is a powerful tool in diagnosing system bottlenecks and analyzing response time at different periods in the scenario.
AEGIS engineering model
command and control system

F. Bernstein | J. Strip

Navy weapon systems are continually growing in sophistication and complexity, and threat improvements dictate more complex and rapid decision-making logic in the command and control (C2) elements. The AEGIS (anti-air warfare) system, currently in the engineering development phase, represents the latest step in this advancement. This paper describes the C2 system under development for the engineering model. A third-generation digital computer and a new computer-controlled general-purpose console that can change operator mode quickly add greatly improved processing capabilities, improved man/machine communications, and great flexibility to AEGIS. To provide a proper framework for this discussion, a brief history and evolution of Navy systems will first be presented and the relationship and impact of AEGIS on these systems noted.

Fred Bernstein, Ldr. AEGIS Command and Control System Design, MSRD, Moorestown, N.J., is a graduate of the Illinois Institute of Technology and has taken graduate courses at the University of Pennsylvania and Drexel University. He joined RCA in 1958 after several years at Philco in TV receiver design, and at Sperry Rand Univac Division, in computer circuits and memory design. Prior to his association with AEGIS, he participated in many MSRD programs and studies including BMEWS, SAM-D, DIR, FPS-95, SHF Antenna Computer Pointing System, and the BMD Test Bed Study. He spent one and one-half years with the New Business Programs Group, RCA Laboratories, where he was technical director of the Obstacle Detection Program for Air Cushion Vehicles (railroads) sponsored by the Department of Transportation. At the onset of the AEGIS program, he helped organize and manage the computer programming activities before concentrating in the command and control area. He has taught several after-hours courses, and has several patents and papers to his credit.

Joseph Strip, AEGIS Advanced Programs, MSRD, Moorestown, N.J., received the BSEE from Newark College of Engineering and the MSEE from Stevens Institute of Technology. He joined MSRD in 1965 after several years at the Signal Corps Engineering Laboratories. Curtis Wright, and ITT where he was engaged in the development of digital communication equipment and flight simulators. Prior to joining AEGIS he participated in a systems engineering capacity in numerous MSRD programs including Land-Based Talos, Atlas, BMEWS and SAM-D. Mr. Strip has been associated with AEGIS since the inception of this program at RCA, concentrating primarily on system definition problems with emphasis on Command and Control. He has taught several after-hours courses and holds several patents.

Reprint RE-19-5-20
Final manuscript received November 8, 1973.

A WEAPON SYSTEM may be defined as a "collection of integrated subsystems comprising material, men, and skills which perform the interrelated functions necessary to render the desired effect on the enemy."

A weapon system consists of four major functional areas: sensors, weapons, support systems, and command and control (C2). Ship's sensors typically include radars, sonars, passive listening devices and optical trackers. Depending on the ship's size and missions, weapons may include surface-to-air and surface-to-air, surface missiles, guns, torpedoes, and electronic warfare systems. Shipboard support systems perform such functions as navigation and communication. Command and Control is the functional area that integrates the sensors, weapons, and support systems into a total weapon system.

Functionally, the command functions performed by C2 pertain to the commanding officer's policy-making role with respect to the use and deployment of resources, and the general conduct of the battle. The control function of C2 deals with the implementation of the policy, and specifically deals with the detailed conduct of the battle.

A C2 system accepts inputs from its own ship sensors or from external sources. It correlates all these inputs as a basis for evaluating the tactical situation. Upon evaluating the tactical situation (and in accordance with established policy and doctrine) the C2 system decides on the action to be taken and issues orders when appropriate to prepare for engagement or actually engage targets.

Historical perspective

In the early days of Anti-Air Warfare (AAW), all C2 functions were performed manually. Targets were sensed and then reported by voice for hand plotting on vertical boards. Command evaluated the situation and assigned targets, by voice, to the weapons.

The introduction of NTDS (Naval Tactical Data System) brought the first radical technological change to CIC (Combat Information Center) operation around 1965, through the introduction of digital computers and general-purpose displays. Although the introduction of
NTDS modified the technology used in CIC's, it did not materially alter the nature of CIC operation. The sensors still provided video for the operator to use for manual-track data-point entry (although general-purpose computers were used to smooth this data, display the tactical situation, and perform computations to assist C² personnel in their decision-making). Target assignments continued to be made by operators (except for special cases), but via consoles rather than voice.

AEGIS impact on C²

AEGIS represents the latest development in the evolution of AAW combat systems. The AN/SPY-I introduces a fully automatic radar to perform search and tracking, and the use of a midcourse-commanded missile increases the firepower of the ship. The introduction of these innovations into the fleet with the attendant high target-handling capability brings a requirement for considerable expansion of automation in C² and requires that command be exercised by negation.

Originally, AEGIS was designed for installation in a 10,000-ton DLGN-38 class ship. The discussion that follows briefly describes this ship's combat system and the role of C² within it.

DLGN 38 combat system

The combat system of the DLGN-38 ship is shown in simplified form in Fig. 1. The AEGIS part of the system consists of Radar Set, AN/SPY-1; Weapon Direction System (WDS), Mk 12; Fire Control System, Mk 99 (consisting of a Tracking Illuminator, Mk 91 and Slaved Illuminator, Mk 90) the Guided Missile Launchers, Mk 26; the midcourse-command-version of the standard missile, SM-2; and finally the AEGIS portion of the C² system, Mk 130.

In addition to AAW, this ship also performs surface warfare and anti-submarine warfare missions. These missions, as well as AEGIS, are supported by the additional sensors shown in the Fig. 1, and by the gun and underwater-battery fire-control systems; the latter shares one of the launchers with AEGIS. The C² system integrates the AEGIS and non-AEGIS elements into a total combat system, and interfaces these elements with command personnel in CIC.

Functionally, the AEGIS system contains three control segments: AN/SPY-1 control, the weapon direction system, and command and control (C²). The functions performed by these three segments are shown in Fig. 2. The AN/SPY-1 control segment controls the operation of the phased-array radar in target surveillance and target and missile tracking. It also controls uplink and downlink communications with the SM-2 missile (RIM 66-C).

The baseline design of the C² segment in the AEGIS combat ship was based on the use of standard digital computers (AN/UYYK-7) and associated peripherals, and an updated general purpose display system, Display Group AN/UAYA-4. Auxiliary equipment in the C² system includes switching equipment, a digital clock, a magnetic disk and a multiplexer to allow non-C² computers to access the disk.

Operational overview of AEGIS

The operation of the AEGIS weapon system is shown in Fig. 3. It begins with a search-and-detect operation by the phased-array radar. A detected target is placed into track by the radar and passed to the command-and-control system for evaluation. Once a target is determined to be a threat by the command-and-control system, orders are furnished to the weapon-direction system to develop a fire-control solution. Missiles are loaded on the launcher by commands furnished by the weapon-direction system. The phased-array radar furnishes target track data to the fire-control system which designates the launcher firing position through the weapon-direction system.

After launch, the missile is guided by the phased-array radar until the homing phase of the flight at which time the fire-control system illuminator is slaved to target coordinates supplied by the radar, thereby furnishing the illumination on which the missile can home. Following intercept, the illuminator is available for the next threat. In the meantime, the multifunction phased-array radar, coupled with the multicomputer control system, continuously searches for new targets and simultaneously tracks targets already detected or engaged. The result is virtually instantaneous response to any new single or multiple threat that appears. These features add up to a firepower capability adequate to meet the threats of 1975 and beyond.

AEGIS EDM-1

Under the terms of the AEGIS contract, RCA is responsible for developing the AN/SPY-1 control and the WDS-Mk 12 computer programs. However, in the case of C², RCA's task was "to develop sufficient C² segment capability to permit satisfactory performance demonstration" of the other segments in the engineering models, and to provide liaison with the combat system contractor.

For convenience of discussion, functions allocated to C² can be subdivided into two categories: "AEGIS functions" consist of the set of functions that must be implemented to allow AEGIS to perform its primary mission of conducting missile engagements; "non-AEGIS functions" performed by C² are associated with non-AEGIS weapon systems or in cooperation with other ships.

The design approach adopted for the first engineering development model (EDM-1) was to establish a basic architecture for the implementation of the "AEGIS functions" of C² that would generally conform to the expected design and processing flow of the final combat system, and implementing these functions only to the degree required to satisfy EDM-1 test objectives. The major requirements allocated to EDM-1 C² are:

1) Establish operating configuration, doctrine, and modes.
2) Establish and maintain system track files.
3) Identify targets.
4) Perform threat evaluation.
5) Assign weapons.
6) Control engagements.
7) Display information.
8) Support EDM-1 testing.

Major EDM-1 functions

The major EDM-1 functions designed to satisfy these requirements are outlined in the following paragraphs:

Initialization—prepares a completely dormant system for operational use, allowing the operational computer program to be loaded from tape or disk or both with minimal operator effort.
**EDM-I system setup**—largely a tactical function which deals primarily with the overall management of EDM-I system operation. It allows the setting of modes of operation. Setup functions provided for EDM-I include:

- **Special threat setup**—permits the entry of special threat parameters into the system, and also permits the enabling or disabling of permission to engage such targets automatically.
- **Console mode setup**—controls the setup of the operational console modes for proper control and utilization of the system.

**Control and operations**—implements control over the various sensor, weapon, and communication systems. It also provides operational inputs required to utilize tactical data and perform the tactical mission. It consists of the following:

- **WDS direction**—establishes the operational doctrine under which C will control the WDS.
- **Special points and references**—permit entering non-tactical points, lines, and circles on displays to aid the operators in maintaining a tactical picture.
- **Own ship position and navigation**—computes ship's position and course, and closest point of approach.
- **Display console control**—includes C program requirements for mechanics of console data entry and output.

**Tracking**—Establishes a system track file for targets tracked by the ship's radars, making this data available to various users. For EDM-I, this function executes the processing involved in accepting raw track data from the SPY-I radar, the TI (tracking illuminator), and the PPI (plan-position indicator) console displays (manual entries); it then establishes and maintains a track file on these targets for use by the tactical functions. It includes:

- **SPY-I automatic tracking**—establishes and maintains those tracks, both active and passive, input by SPY-I. Processing includes initiation and maintenance of normal and burn-through tracks, assignment of system track numbers, track quality updating, and establishment of target category.
- **Manual tracking**—establishes and maintains tracks as in SPY-I above except that data results from the actions of operators at the PPI display consoles in the C.
- **TI tracking**—maintains a track in the track file for which the tracking illuminator (TI) is transmitting data via WDS. System track number assignment and track quality updating for this track are included.
- **Track Management**—includes: operator control actions that deal with tasks such as selection of primary sensors for a specific track, ECM actions on passive angle tracks, and designation and dropping of tracks; and computer processing related to monitoring track load, track quality, and handling of messages related to track management, such as dropped track reports.

**Tactical functions**—includes all C data processing associated with evaluation and, if required, engagement of operational tracks. Although the information flow through
UYA-4 is a new highly flexible general-purpose console which allows the computer to control the role of the console at any time in the system. The console in use is shown in Fig. 4. The representation of the console in Fig. 5 highlights the main features.

The PPI display (the large round CRT) provides the range-bearing overview of the space (air and surface) surrounding own ship (which is usually located at the center of the display). Various symbols are used to show the location of the targets as well as identify their nature (air, surface, friendly, unknown, hostile, etc), Radar video may be displayed, and an RHI (range-height indicator) form of presentation may also be selected.

The rectangular CRT in the center upper portion of the console provides amplifying alphanumeric information on a selected target or other information that may be selected by the operator. When a specific procedure is required to enter data, the CRT leads the operator through the procedure via a tutorial.

The pushbutton switches labeled AEB's (action-entry buttons) are the main operator interface with the system. Each of the 18 AEB's displays a label (a projection of a film chip) which identifies its unique meaning. Each AEB may assume any of 48 meanings, all under computer program control. Use of these AEB's is described later in the paper.

The direct projection readouts (DPRO's) at the top of the console provide fixed information to the operators. Each of the 36 DPRO's may display any of 12 labels (film chips) selected by the computer program. For EDM-I, the DPRO's on the left are used to provide weapons system and engagement status information. The DPRO's at right are used to "alert" the operator to various general conditions and data of interest and to "orders" that require action on his part.

Various fixed-entry buttons (FEB's) provide single-action fixed-meaning controls. The FEB's, directly under the AEB's, provide control of the console state (on-line or in test), or provide a means for the operator to return to a higher functional state. The FEB's located in the track ball well (lower right hand corner of Fig. 5) provide control of a ball tab symbol (cursor) used by the operator to identify targets to the C^2 computer program. The digital data entry unit is used, in conjunction with the action-entry buttons, to enter numerical values.

Console manning organization

The overall system control of AEGIS is exercised through the console operators. These operators insert doctrine, select targets for various actions, and establish modes of operation. The UYA-4 consoles described above are usually collocated in a Combat Information Center (CIC) for ease of operator communications. The EDM-I organizational structure (operators' hierarchy) is patterned after the Navy Tactical Data System (NTDS) in the DLGN-38, but is simplified in EDM-I in keeping with the limited objectives. The C^2 console-manning organization is shown in Fig. 6.

The ATC (AEGIS tactical coordinator) is the lead operator. The main functions of the ATC are to supervise the overall conduct and operation of the CIC; insert tactical doctrine, threshold values, and own ship position data; and evaluate the computed threat potential of targets and assign weapons to the MSS/EC when required.

The MSS/EC (missile system supervisor/engagement controller) combines positions normally separately manned. He provides operator control of the weapon system. His main functions are to assign launcher and illuminator and mode of operation; monitor launcher and illuminator status; monitor engagementability of selected targets, control firing and re-firing, monitor missile status, and monitor air engagements.

The SS (sensor supervisor) interacts with the detector tracker (see below) and provides a degree of track management. His main functions are to delete targets from the C^2 system if the track load becomes too great, perform manual tracking using radar video displayed on the PPI, and enable use of SPY-1's burnthrough mode.

The DT (detector tracker) reports to the sensor supervisor in the organizational structure. He has a direct counterpart in the NTDS system and does manual tracking in support of the SS.
The TD (test director) is a position unique to EDM-1 AEGIS and has, accordingly, no direct counterpart in the NTDS. He provides control of test-related items, and also a few functions that are normally tactical but because of their simplified form (and because they represent major test control features) were given to the TD. His main functions are to control test operations, initiate and control recording of test data, control special tracking use of the Tracking Illuminator, initiate and control CRT display of special data, and control identification of targets (friend, unknown, hostile, air, surface, etc).

**Operator/program interaction**

As mentioned earlier, the system operators communicate with the C² computer program via the UYA-4 console. The primary vehicle for this is the action-entry button (AEB) panel shown in Fig. 5. Since each pushbutton may assume any one of the 48 meanings, all of which are under computer control, the panel array presents a truly flexible means of operator/computer program interaction. This flexibility was channeled using the mode-state concept discussed below and depicted in Fig. 7.

An operator will select his primary mode (TD, ATC, MSS/EC, SS, or DT) at initialization time when all console AEB arrays display the mode select labels. The operator will then depress the mode label AEB assigned to him and the AEB panel will automatically change to the new set of labels which pertain to the normal operating state of that mode. Many actions are possible with the AEB's available in the normal state; in addition there are many states available to him.

Some of these states are available only in a particular mode, while others are available commonly to all or several modes. Fig. 7 shows the hierarchy of states and substates that pertain to the ATC. Those unique to the ATC are interconnected with solid lines. Those common to several operator modes are shown dotted. AEB's that cause a change of state are outlined more heavily.

When an operator is in a substate and wishes to return to a higher state or substate in the hierarchy, he may depress the fixed entry button (FEB) marked canc staté which will return him one level, or rtto nos which will return him to the normal operating state, or the mode sel which allows him to start over again and choose a new operating mode (within the rules in the program concerning multiple operator mode assignments).

In any given state, the direct-projection readouts (DPRO's) and the alphanumeric CRT provide auxiliary information to the operator. Fig. 8 portrays various displays that the AEGIS tactical coordinator (ATC) might see at some particular time in the normal operating state. The left third of the rectangular CRT shows data on a target that has been "hooked" using the track ball well (lower right hand corner of Fig. 8). The center third of the CRT shows own ship data remaining from an ATC action while he
Computer program

The nerve center of the AEGIS system is the large, sophisticated C2 computer program. It runs in an AN/U YK-7 computer, a large, fast, third-generation computer. Some highlight features of the AN/U YK-7 are:

- Up to 256k (32-bit) words of memory
- 1.5 μs read-write memory cycle
- Overlapped memory operation
- Separate I/O controller with 15 basic instructions
- 16 full-duplex NTDS-compatible I/O channels
- 130 basic whole and half-word instructions
- Direct and indirect addressing
- Two sets of 7-index and 8-base registers

At the outset, the various programming tasks involved in the AEGIS software development were formidable. The "givens" were a new computer, a new compiler still under development, and a new executive program that had yet to be developed. To facilitate development of the operational programs, several ground rules were established to ease the programming task by permitting efficient utilization of inexperienced programmers on a large concurrent basis. The main rules were:

- Use the CMS-2 compiler which is easier to learn and apply than assembly language programming. The cost in core and running time was not considered to be severe (10 to 15%). Recode the time-critical elements of the program in assembly language as necessary.
- Use the same executive program (called ATEP for AEGIS tactical executive program) for all segments. The segments had similar requirements and therefore the common usability of ATEP was high.
- Use a "modular" design. Each module was to contain all or part of a functional entity. All modules were to operate under ATEP control and communicate with the "outside" world and with other modules via ATEP. All modules were to have direct access to "common" subroutines (e.g., math routines) and to have a "common" data base. This approach permitted a large number of programmers to be employed simultaneously.

The C2 computer program was designed with these rules as the framework. As one additional architecture consideration, the C2 program was designed using a limited "overlay" concept. Succinctly, a selected group of program modules, generally those which were not tactically time-critical, were kept on disk and read into an "overlay" area of core memory when required. This approach saved core memory in that one reasonably small overlay area may be used to service many modules on disk (care must be taken to select modules for disk that are not likely to be required in core memory at the same time). This space saving was not ultimately required for EDM-1, but it was decided to complete the development of the technique.

The result of the C2 computer programming effort is a program that requires approximately 60k core memory locations, common service routines of approximately 5k locations, data base of 8k locations, temporary and dynamic
storage of approximately 12k locations, and resident core modules of approximately 26k locations. In addition, the non-core resident modules (those on disk) total an additional 16k locations (if they could all be loaded in core simultaneously). From a timing viewpoint, the program uses about 60% of available running time under severe load conditions.

Present status

Development of AEGIS EDM-I began in January 1970. System testing has been completed at the land-based test site located at Moorestown, N.J., and the effort has transferred to the U.S.S. Norton Sound for sea tests off the coast of California. Current design utilizes the SM-1 missile, with the SM-2 missile to be fully incorporated later.

The C^2 system passed its qualification tests in June 1973. The tests were conducted with a simulation program providing the stimuli for those elements external to C^2 (the ship's navigational system, the SPY-1 radar, and the weapon system). The simulation program ran in a separate AN/UYK-7 computer and was interconnected to the C^2 AN/UYK-7 computer via the same I/O channels used for the actual elements. As a unique approach to the qualification tests, Naval Training Unit and Technical Representative personnel operated the UYA-4 consoles during both the test debug stage and the actual formal test. This provided early familiarization by Navy personnel, and also provided valuable operational feedback for future improvement.

As to the future, studies conducted concurrently with EDM-I development showed that the overall AEGIS can be reduced in size. Preliminary design is proceeding currently for a new class ship, the 6000 ton DG (compared to the 10,000 ton DLGN-38) to obtain further verification for this approach.

Acknowledgments

Credit is due to G.V. Metzger and S.A. Stern of the RCA AEGIS EDM System Design Group and to S.A. Raciti of the RCA AEGIS Project Management Office for their many valuable contributions during the conduct of the program, and to the C&C programming group at Computer Sciences Corporation for their efforts in design, execution, and test of the computer program.

References

1. OP-300: Weapon System Fundamentals, "Definition of Systems"
Command and data handling for Atmosphere Explorer satellite

W. V. Fuldner

The basic mission of the Atmosphere Explorer satellite is performed by controlling a full instrument complement, either in concert or individually; a regular, periodic operational program is not adequate for this mission. All telemetry data is gathered and formatted into a single serial bit stream, which is transmitted to Earth either in real-time or after storage. The command and data handling (C&DH) subsystem provides the necessary controls for the instrumentation and telemetry, and also controls the satellite's attitude and trajectory. These requirements dictate a complex and sophisticated design. The physical limits of the satellite introduced additional constraints. The design and construction of this C&DH subsystem are described in this paper.

The mission of the Atmosphere Explorer Program is to perform a twofold examination of the lower thermosphere (120-300 km altitude) involving quantitative measurement of parameters and "in concert" experimentation. This mission will be accomplished by a number of satellite-borne instruments performing correlative measurements of the chemical constituents, and of the solar input and resulting response. Table I summarizes the mission instruments and their pertinent characteristics.

The program uses three spacecraft: AE-C, -D, and -E, launched into elliptical orbits (nominally 150 by 4000 km), each at a different inclination to permit examination of the lower thermosphere over a range of seasonal variation and geographic peculiarities. AE-C was launched on December 15, 1973, and its early mission performance has met or exceeded all objectives.

Each spacecraft has an orbit-adjust propulsion subsystem (OAPS) to permit perigee excursions to altitudes of 120 km to further facilitate experimentation and to permit apogee restoration so that the one-year mission life may be ensured despite orbit decay caused by atmospheric drag.

The nature of the atmosphere in the low-perigee regions imposes severe limitations on the physical configuration of the spacecraft, especially as it affects aerodynamic characteristics, projected frontal area, and control of center of mass. The latter is critical to both the aerodynamic behavior and the thrust effectiveness for the orbit-adjust propulsion subsystem. The configuration selected (Fig. 1) resembles a cable spool, the outer surfaces accommodating the electronics, and the toroidal center section housing the propulsion subsystem. A two-piece solar array, with apertures as required for the instruments, envelopes both sections. This physical configuration placed significant constraints on the design of the command and data handling (C&DH) subsystem, which is the topic of this paper.

### Command and data handling requirements

#### Command

The C&DH subsystem executes all command information within the spacecraft, either in real time (as received over the S-band command transmission link) or remote from the command site (as required by the orbit operation schedule). The command subsystem checks the validity of the information received in real time, independent of whether it is to be executed in real time or is to be stored for execution later in that orbit or in a subsequent orbit. The combined STADAN/MSFN network (now identified as STDN) is used for communicating with the Atmosphere Explorer. Command information to the

---

**Table I — Experiments and engineering measurements in Atmosphere Explorer Satellites C, D, and E.**

<table>
<thead>
<tr>
<th>Experiment</th>
<th>Flights</th>
<th>Detector</th>
<th>Instrument characteristics</th>
<th>Parameters</th>
</tr>
</thead>
<tbody>
<tr>
<td>EUVS: Solar EUV spectrometer</td>
<td>C, D, &amp; E</td>
<td>Channel electron multiplier</td>
<td>140Å to 1850Å</td>
<td></td>
</tr>
<tr>
<td>ESUM: Solar EUV filter photometer</td>
<td>C, D, &amp; E</td>
<td>Spiral electron multiplier (SEM) and EUV diodes</td>
<td>40Å to 1300Å</td>
<td></td>
</tr>
<tr>
<td>UVNO: UV nitric oxide</td>
<td>C &amp; D</td>
<td>Photomultiplier tubes (PMT)'s</td>
<td>2350Å, 2190Å</td>
<td></td>
</tr>
<tr>
<td>VAE</td>
<td>C &amp; D</td>
<td>PMT's</td>
<td>6300, 5577, 4278, 3371Å</td>
<td></td>
</tr>
<tr>
<td>OSS: Open-source neutral mass spectrometer</td>
<td>C, D &amp; E</td>
<td>Electron multiplier</td>
<td>1 to 46 AMU</td>
<td></td>
</tr>
<tr>
<td>NACE: Closed-source neutral mass spectrometer</td>
<td>C, D &amp; E</td>
<td>Electron multiplier</td>
<td>1 to 46 AMU</td>
<td></td>
</tr>
<tr>
<td>NATE: Neutral atm. temp. exp't.</td>
<td>C, D &amp; E</td>
<td>Electron multiplier</td>
<td>T_e, N</td>
<td></td>
</tr>
<tr>
<td>MESA: Atm. density accelerometer</td>
<td>C, D &amp; E</td>
<td>Accelerometer</td>
<td>Neutal density</td>
<td></td>
</tr>
<tr>
<td>RPA: Planetary ion trap</td>
<td>C, D &amp; E</td>
<td>Electrometer</td>
<td>T_e, N, M, drift velocity</td>
<td></td>
</tr>
<tr>
<td>CEP: Cylindrical electrostatic probe</td>
<td>C, D &amp; E</td>
<td>Electrometer</td>
<td>T_e, N, M,</td>
<td></td>
</tr>
<tr>
<td>MIMS: Magnetic ion mass spectrometer</td>
<td>C &amp; D</td>
<td>Electron multiplier</td>
<td>1 to 46 AMU</td>
<td></td>
</tr>
<tr>
<td>IMS: Positive ion mass spectrometer</td>
<td>C &amp; D</td>
<td>Electrometer</td>
<td>0.5 to 72 AMU</td>
<td></td>
</tr>
<tr>
<td>IEE: Low-energy electron exp't.</td>
<td>C &amp; D</td>
<td>SEM's</td>
<td>0.2 to 25 keV</td>
<td></td>
</tr>
<tr>
<td>PES: Photoelectron spectrometer</td>
<td>C, D &amp; E</td>
<td>Johnson electron multiplier</td>
<td>Photo-electron spectra</td>
<td></td>
</tr>
<tr>
<td>Three-axis fluxgate magnetometer*</td>
<td>D &amp; E</td>
<td>Fluxgate magnetometer</td>
<td>Magnet variations</td>
<td></td>
</tr>
<tr>
<td>Capacitance manometer</td>
<td>C, D &amp; E</td>
<td>Electrometer</td>
<td>Pressure</td>
<td></td>
</tr>
<tr>
<td>Cold-cathode ion gauge range</td>
<td>C, D &amp; E</td>
<td>Diaphragm</td>
<td>Pressure</td>
<td></td>
</tr>
<tr>
<td>Temperature alarm</td>
<td>C, D &amp; E</td>
<td>Grid wire</td>
<td>Aerodynamic heating</td>
<td></td>
</tr>
</tbody>
</table>

*Engineering measurement on 'C'

T_e: electron temperature  
N_e: electron density  
T_i: ion temperature  
N_i: ion density  
T_g: gas temperature  
M: ion mass
The spacecraft is transmitted via a 70-kHz subcarrier on an S-band uplink.

The C&DH subsystem fulfills the basic mission of the Atmosphere Explorer Satellite by controlling the operation of the full instrument complement, either together or individually, depending on the nature of the information required. Since instrument operation is related to orbit altitude or to identified peculiarities in the behavior of the upper atmosphere, (also in some cases to solar input to the atmosphere) a regular, periodic operational program is not adequate for mission performance. Furthermore, the power available and the capacity for storing instrument output data demand that a high degree of flexibility be allowed for optimum use of these two resources throughout the orbit. This demands that the C&DH subsystem contain a high degree of flexibility, further complicated by the rather complex nature of many of the instruments in terms of the various modal operating characteristics.

**Power control**

In addition to the requirements for supplying the command information, the C&DH subsystem must also apply and remove power to the various instruments in such a way that the instruments are used most effectively with minimum power consumed. In certain cases, unfavorable orbital conditions may compromise the life of these instruments; thus, independent power control of each is required.

**Attitude and orbit adjustments**

The C&DH subsystem must also control the satellite magnetic torquers. Two sets of torquers are used: one for spacecraft attitude control and the other for spacecraft momentum management. To modify the orbit as required by the mission plan, the subsystem also controls the orbit-adjust propulsion subsystem at the dictated orbital locations. The capability to do so is required both under real-time command contact or remotely from the command sites.

**Telemetry and timing**

The C&DH subsystem controls the gathering and formatting of all the telemetry data from the instruments and the spacecraft equipment into a single serial bit stream which is either transmitted to the ground in real-time or is stored by on-board tape recorders for playback at a subsequent ground station contact. All generated data, whether it be the mission science data or the diagnostic telemetry, is handled in the same fashion and is multiplexed into the same data stream. The subsystem has the capacity to assemble digital telemetry words, prior to commutation, from individual status or flag bits of data. Subsequent to assembly, these words are commutated into the serial data stream in the same fashion as the conventional digital data.

In controlling the telemetry, the C&DH subsystem also provides all of the sync and timing signals required in the spacecraft. The large majority of the sync signals required are crystal derived and phase related to the serial data bit stream, which is also related to the on-board generated spacecraft time code. The telemetry subsystem, in the process of gathering the digital telemetry, causes digital word interrogation signals (called word enables) to be generated; these are used to control the readout of the digital telemetry from the various sources. Similarly, the clocking signals for the readout of such data is phase-related to the clocking signals which the telemetry processing equipment uses to interrogate and format the data.

Other sync signals, which are not time-related to the telemetry processing provided by the C&DH subsystem, include a command clock and spacecraft nadir-related pulses. These latter signals are derived from the satellite's horizon sensors and are used to perform synchronizing functions in the instruments and in the torque commutation) which are related to the instantaneous orientation of the spacecraft.

**Operational redundancy**

The C&DH subsystem provides the above basic functions on a fully redundant basis, such that a single, realizable failure will not preclude the subsystem from performing the functions required.
The equipment complement comprising the C&DH subsystem includes: a dual command decoder (DEC); a dual PCM controller (PCMC); a dual programmer (PROG); two command memory units (CMU); two dual remote telemetry modules (RTM); two command distribution units (CDU); and two logic interface units (LIU).

Selected configuration

In the equipment comprising the C&DH subsystem, the logic fan-out is accomplished by the two logic-interface units. These units deliver logic-interface signals to the instrument complement and to the spacecraft equipment. Similarly, the power distribution is accomplished by the two command distribution units. Two remote telemetry modules are used as the receiving stations for the generated telemetry.

Signal traffic between the two baseplates must be accomplished from the perimeter of one baseplate to the other; the center column is not available. With this configuration, if signal traffic were required between worst-case locations on each baseplate, cable runs of up to fifteen feet would have been necessary and large numbers of signals would be traversing the center section, causing severe harnessing and coupled-noise problems. Accordingly, the basic approach to the signal traffic control is to develop localized distribution "stations" on each baseplate for the fan-out of the logic signals; similar separate stations for the fan-out of the switched power; and, working in the opposite direction, localized receiving stations for the fan-in of the telemetry data.

The units that generate the source information to be distributed (decoders, programmers, PCM control) are configured to source their output information serially to the two logic interface units (LIU) and command distribution units (CDU) on the other hand are not redundant with respect to each other, as each contains the fan-outs required by the equipment on its associated baseplate. Internally, however, each of these fan-out units is redundant, with circuitry associated with each redundant source device (DEC, PROG, PCMC) providing redundant functions which are or'ed at the output. Similarly, the dual remote telemetry modules are unit redundant, with one half of RTM A and one half of RTM B associated with each redundant PCM controller (PCMC).

The basic approach to signal distribution within the spacecraft is to transmit source information serially from the generating unit (such as the decoder or PCMC) to the localized fan-out units (LIU A and B). By so doing, the number of signals required to pass through the center region of the spacecraft is minimized. This permits both easier harnessing and, by restricting the volume of traffic through this region, permits use of more sophisticated and higher-power circuitry to provide the required noise immunity to accommodate the long cable runs. Additionally, localized distribution for the logic-level signals reduces the maximum cable lengths required, thereby enhancing the noise immunity and permitting simpler interface circuitry to be used for the large volume of signals on each of the baseplates.

The same improvement in noise immunity is achieved in the gathering of the telemetry data on each of the baseplates, such that the system requirement to provide an 8-bit accurate system for analog telemetry can be accomplished without major power budgeting and circuit complexity for the analog telemetry interfaces.

The system is required to interface a large volume of signal traffic between various families of logic devices (COS/MOS, TTL) which are operated at different bias voltages. To preclude the necessity of developing various interface voltages in the source units for each different interface requirement, the logic interface source circuit provides an open or saturated switch to ground corresponding to a logic 0 or logic 1, respectively.

The output of the source circuit is delivered to a receiving circuit which contains a "pull up" resistor to that logic family's operating voltage. Hence, when the circuits are interfaced (and the receiving circuit is powered), the interface signals appear as either ground or +Vbb, without the requirement to generate the appropriate Vbb in each source device. The same approach to the logic interface is used in the reverse direction for digital telemetry data from the instruments to the RTM's.

Command types

The command system delivers two types of commands to the spacecraft: major-mode commands and minor-mode commands. A major-mode command is a single pulse, 34 ms in duration, which is uniquely associated with a particular command operation. Within the system, the major-mode command can be delivered at logic levels or with sufficient power to operate relay coils. The minor-mode commands are 32-bit-long serial data words which are delivered at the logic level only. When minor-mode commands are distributed in the system, a unique major-mode command is also simultaneously delivered. This is used as a destination address signal, such that minor-mode command information may be uniquely routed to a particular user. Fig. 2 displays the timing relationships of the major-mode and minor-mode commands. Note that if a major-mode command is being generated, the timing as shown in the illustration is the same, but there is no accompanying minor-mode data. The command clock is also generated, as shown in the illustration, whenever any command is being executed within the system.

Command format

The serial command information transmitted to the spacecraft is coded in the format shown in Fig. 3. The partitioning of the bits into a sequence of 8, 16, 32, and 8 bits as shown is done to facilitate Operations Control Center (OCC) computer operations on the information. The first 7 bits of the command word are allocated by NASA on a "per-spacecraft" basis. The decoders permit, by patch plug, selection of the 7-bit address subsequent to flight qualification of the unit.

Decoder functional behavior

The decoder receives the 70-kHz command subcarrier from its uniquely associated command receiver (part of S-band transponder) and demodulates the two linearly added 1 and 2 kHz command information signals. Processing of these two signals provides the digital command
and clock input to the logic section of the decoder.

As both decoders are continuously powered throughout the spacecraft life, two command bits are allocated for selection of the decoder to be used in the processing of the real time received command. The data-bus control bits are used internal to the decoder to indicate whether or not the information contained in the 32-bit segment (bits 25 through 56) is to be delivered out of the decoder as a minor-mode command. The 9 bits indicated for the op-code are the encoded bits identifying the major-mode command to be executed. These 9 bits provide for 512 unique codes; however, the subsequent processing of the command information reduces this glossary capability to 496 unique commands. The final 7 bits of the 64-bit command frame contain a cyclic check code. This particular encoding technique provides a Hamming distance of four over the entire command information. Functionally, the check is accomplished by passing the entire 64-bit code through a feedback register configured with a particular exclusive-OR feedback pattern. For the code to be valid, after the 64th bit passes through the register, its contents will be all zeros. This will not be the case if bit errors have occurred up to the capability of the code.

Command processing

When processing a command, the decoder performs an initial check for satisfactory spacecraft address, upon achieving sync with the command bit stream (a sync pattern of 63 zeros followed by a 1 is used in the system). If achieved, the decoder address is then checked, and the selected decoder continues to process the command information while the other decoder returns to sync search. Inspection of the data-bus control bits indicates, as noted above, what is to happen with the information in the minor-mode command-bit locations. The op-code is placed in a register in the decoder and held until the 64-bit command frame is complete. Upon satisfactory completion of the cyclic code check, the 9-bit op-code is shifted out of the decoder into both LIU's.

Upon entering the LIU's, the 9-bit op-code is converted to a 32-bit-long binary word which is stored in a register. The code conversion is such that the resulting 32-bit word contains exactly 2 logic ones and 30 logic zeros.* The final decode process is then performed by hard-wire decoding of the two data ones into an and gate, the output of which is the unique major-mode command. The logic performing the code conversion and the execution (final decoding) of the command is under the control of the logic sequencer of the decoder which processes the command. The control is such that the output of the register into the hardwire and gates is permitted to exist for 34 ms (thirty-four counts of the 1-kHz command clock).

The above description indicates the flow of information for the generation of a major-mode command at the logic level. If the command is to be generated at the voltage level to provide a relay driving pulse, the process is the same, except that the final decode is not performed in the logic interface unit (LIU), but rather, the outputs of the 32-bit-long register are sent (in parallel) to the command distribution unit (CDU) on the same baseplate as the LIU where the decode process is completed (in this case at a -21-V level), with sufficient drive capability for the relay coil load. Fig. 4 displays the basic flow of information through the decoders to the LIU's and CDU's. As shown, two decoders and their associated circuitry are indicated. The circuitry shown in the LIU and CDU, which is associated with each decoder, is duplicated in the second LIU and CDU (not shown). The LIU and CDU circuitry associated with a particular decoder is powered from the dc-to-dc converter of that decoder. Thus, the LIU and CDU circuitry associated with a decoder can be thought of as an extension of that decoder, rather than part of a different unit.

Minor mode and command clock distribution

The implementation techniques for the minor-mode commands and the command clock are quite similar (see Fig. 4). These logic-level signals are fanned-out of the LIU using the same basic final decode gate that is used for the major-mode commands, the difference being that these signals are applied to one input of the and gates while an enable signal is applied to the other. In the case of the command clock, the enable signal is the sequencing command from the decoder which controls the execute time for the major-mode command. In the case of the minor-mode command, the second input to the and gate is an enable that is generated by the decoder (bracketing the time for transmission of minor-mode command) which is generated only when the data bus control bits in the real-time command indicate the presence of minor-mode data. For commands executed out of the command memory, an equivalent signal is sent to the decoder from the memory to permit the same sequence control operations as for the real-time commands.

*This conversion results in the glossary reduction from 512 to 496 available commands.
Programmer functional behavior

The programmer performs several command and control functions required by the spacecraft and the experiments. These include operation of the command memories, attitude control, generation of nadir pulses, and operation of the propulsion subsystem. The operating control of the attitude-control torquers and the orbit-adjust propulsion subsystem is accomplished by dedicated sequencers which are controlled by commands, either real time or stored in the memories. Thus, the memory programmer portion of the programmer is integral to most of the functions performed. There are two (redundant) programmers in the spacecraft, only one of which operates at a time. Either programmer can operate with either of the redundant command memories, only one of which is in use at any one time. As the memories are non-volatile, the maximum command storage capability can be accomplished if both devices are concatenated.

Memory programmer operations

The subsystem remote command capability is provided by executing commands which are stored in the command memory. The time of execution is determined by a bit-for-bit match of the 16 least significant bits of the spacecraft time code with a “time tag”, which is loaded into the memory with the command to be executed. If the command to be executed is a major-mode command, 32 bits of memory space are required, while 64 are required for a minor-mode command and its associated major-mode command. Each of the command memories has 32 kilobits of loadable storage; thus the total system capacity, using both memories in series, is 2048 major-mode commands.

The complete memory is scanned every four seconds (granularity of the spacecraft time code) with all time-tag information examined for time-tag match. In the event that multiple commands are tagged for the same time, they will be executed in the order in which they are found in the memory. The execution rate of commands is limited to one command every 64 ms, the same as real-time commands. The memory programmer, as noted above, executes on match only in comparing the time code to the time tags. The above-indicated command-execution rate permits approximately 64 commands to be executed in the four-second clock granularity (ignoring the scan time) until the spacecraft time code is updated. To preclude “missing” commands, either because too many are tagged for the same time event or because the real-time command link is in use and has taken priority over commands out of memory, the memory programmer contains logic, called the “slip-time counter.” Functionally, the slip-time counter keeps track of commands which are scheduled to be executed in a time scan and causes them to be executed as soon as the system is available.

PCMC and RTM functional behavior

A PCM controller and two associated remote telemetry modules, working in concert, provide the data-handling function of the C&DH subsystem of the spacecraft. The two RTM's, one located on each baseplate, serve as the commutators of the telemetry data generated on those baseplates, operating under the control of their associated PCMC. Once gathered, the telemetry data, either digital or analog, is processed by the controller to result in the single serial digital data stream of the spacecraft. This data stream may be recorded on either of two tape recorders or may be transmitted in real time either by the redundant S-band downlink or by the vhf beacon transmitter.

Telemetry format

The output of the PCM controller is a 16,384 b/s bit stream containing 8-bit words which represent the main frame format. The sync pattern in the first three words is assigned by NASA in accordance with the data standards for the combined networks which will service the program. The 16-bit command verification word following the sync pattern contains 11 bits of verification data and 5 bits of an error-correction code which is encoded in the spacecraft. Several sub-coms, of various lengths, are also located in the frame. Eight 4-word blocks are assigned to be overwritten when a command memory dump is in process. During the time that such a memory dump is in process, the command verification word will carry telemetry indicating that a dump is in process and the data appearing in these assigned 4-word blocks should be treated accordingly.

Format control

The various design requirements of the experiments would prefer the choice between digital multiplexing of output
telemetry data and/or reception of digital word-enable signals on either separate or common lines for the different telemetry words generated within a unit. The digital word-enables are logic-level sync signals which are transmitted to the various telemetry source units to indicate that the telemetry word is to be read out of the unit. These signals will either be 8-, 16-, or 24-bit periods long, depending on the length of the word to be interrogated.

The telemetry format is controlled by a read-only memory (ROM) in the PCM controller. The configuration of the controller is such that only the ROM must be changed from mission to mission.

In responding to the information contained in the ROM word, the PCM controller first identifies which remote telemetry module will be in use for the upcoming telemetry data and whether that data will be digital or analog when received. The digital contiguous indication is provided such that the basic generated “word enables” will be continuous over the N-multiple \((N=2, 3)\) 8-bit words. This facility is used to handle the 16- and 24-bit long digital telemetry words which appear in the telemetry format.

Additional information in the ROM controls subcom selection and also contains information for the generation of word-unique digital interrogation signals \((i.e.,\) the “word-enables”).

**RTM gate addressing and digital word-enable generation**

When a digital word-enable is transmitted to a source unit, the appropriate receiving gate must be opened in the remote telemetry module. This may be the same gate which is used for other digital telemetry words from the same source unit or may be a different gate, depending on the instrument design.

The digital word-enables are generated and fanned out in exactly the same manner as the major-mode commands described earlier. The PCM controller is the source of a 9-bit word which is transmitted to the logic interface unit for code conversion into a 32-bit word with exactly 2 logic ones and 30 logic zeros. The final decode from this register is accomplished in the same fashion as the major-mode command. The timing of the word-enable, while being decoded, is under control of the PCM controller in use.

For conditions where there is a “one-for-one” correspondence between use of a digital word-enable and gate from the remote telemetry module (RTM), the 9 bits in the ROM-stored word are used to generate both the word-enable and the RTM gate address. If the same RTM gate is to be used for different word-enables, information in the ROM word is used to address a separate small ROM in the PCM controller which generates the RTM gate address. Thus, the system design requirement to allow instrument choice is accomplished.

### Analog vs. digital telemetry

Two types of input gating circuits are used in the RTM’s in the system, one for digital information and the other for analog information (either positive or negative analogs). The digital input gate employs a “pull-up” resistor to the operating voltage in the RTM such that the digital interface technique requiring that the source unit supply only an open or saturated state to ground switch for logic zeros and ones respectively is maintained. The analog gate includes additional filtering capacity to provide satisfactory noise immunity for the analog signals.

### Time-base configuration

The spacecraft is configured with two redundant countdown chains in each logic interface unit (LIU), one each associated with one of the two redundant crystal oscillators in the system. The design of the system is such that one of the two oscillators (one located in each LIU) supplies the input to two countdown chains associated with it (one in each LIU) to provide the frequencies required on the LIU’s baseplate.

Only one of the redundant countdown chains in each LIU is in operation at any one time, although both crystal oscillators are continuously powered. Because of the large number of fanned-out frequencies required on each baseplate, this configuration was chosen to minimize the inter-baseplate harnessing of high frequencies. To ensure that the proper phase relationship is generated between like signals on the two baseplates, a closed-loop phaselock is maintained between the two parallel countdown chains.

The output configuration of the time-base generated signals is such that the same approach to output redundancy coupling is used as that for the command clock signals. In this case, the time base that has been selected applies an input to one of the two inputs to an and gate, while the particular frequency to be outputted is applied to the other. The output coupling configurations for all signals at or below 16,384 b/s are so coupled and utilize the ground or open output signal configuration with the pull-up resistor in the receiving circuit. For high-speed signals, either localized or from baseplate to baseplate, the high-speed interface is used.

### Sync generation

In the paragraphs dealing with the PCM controller, it was stated that the sync signals required by the experiments are primarily related to the generation and control of the digital telemetry. All of these signals are related back to the time base which forms the clock standard for all operations of the PCM controller. All of these operations are related to the spacecraft 24-bit time code, such that time correlation for operation of the experiments and synchronization of remote command execution are also available.

### Summary

This paper has touched briefly on the various functions performed in the C&DH subsystem for Atmosphere Explorer satellites. It should be noted that the design of the subsystem is such that it is readily amenable to addition and modification should new command, sinc, or telemetry allocation be required. The localized “station” concept for interfacing and the standard approach to interfaces provide adequate noise immunity and a high degree of versatility. While the satellite requires two sets of such stations, the subsystem is not limited to this number; additional stations may be added by using spare bits in the command-word structure and ROM word structure. As these data-source locations (PCMC, DEC) utilize high-speed interfaces to the local stations, much larger systems may utilize the same basic subsystem with no modification to the basic design.
Performance monitoring and fault diagnostics of command and control systems

T. Taylor | M. LeVarn | J. O’Connell

In this paper, the authors describe RCA techniques devised for automatically monitoring command and control system performance, diagnosing faults, and correcting system failures to assure a high degree of systems availability. Typical command and control system configurations and associated automatic test requirement are described. Maintenance requirements, fault detection, fault isolation, and diagnostic approaches are discussed. The test techniques for both on line and off line system testing are presented. The complete test system provides a configuration in which automatic and manual controls combine to provide a complete diagnostic program for C&C systems.

MISSION REQUIREMENTS of modern command and control systems, such as TIPI and 485L, demand a continuous 24 hours/day operational capability. Detection and correction of system failures within a mean-time-to-repair (MTTR) of 15 to 30 minutes, and an availability criteria of .998 or greater, are typical.

Such requirements are satisfied largely through the resources of the test and monitoring system employed. Special attention directed towards system design and development is mandatory, not only for the reason of meeting technical criteria, but also from the standpoint of economics.

Design objectives

Test and monitoring systems developed for TIPI and 485L command and control systems are described in this article as examples of approaches meeting such demands. The design of the test and monitoring system follows the results obtained by MIL-STD-499 functional analyses, and overall plans for an integrated maintenance and logistic subsystem.

Major design objectives include:

a) Maximum utilization of existing test and diagnostic software inherent in the selected central and satellite processors.

b) Maximum employment of existing fault detection and isolation capability within each hardware configuration item (CI) such as: Built-in test equipment (BITE), status reporting via data transfers and ready lines.

Reprint RE-19-5-24
Final manuscript received Jan 3, 1974
c) Recognition of existing maintenance test resources in the next higher echelon of organization structure of which the command and control system is a part, and how these resources may complement the subject test system.

Typical command and control system configuration

Before discussing test system design and diagnostic approaches, a brief description of a typical command and control system is in order. Principal command and control (C&C) missions provide automated assistance to such tasks as combat intelligence, mission planning, air-strike assessment, weapons deployment, and air traffic control. Most C&C systems are basically data processing oriented with heavy emphasis on peripherals. These devices include standard keyboard entry terminals, video tabular and graphic displays, digitized network input devices using modems and teletype cryptos, radar/digitized video, printer, and large-screen status displays. Such a system (Fig. 1) represents Air Force thinking for an integrated tactical air command and control system (TACC). The system provides a modular display subsystem of tabular, graphic and group displays; a data processing subsystem using the Univac AN/UYK-7 or comparable computer; and a communications subsystem. The system is packaged in deployable, expandable, environment-controlled shelters.

In this system any number of remote (up to 150 miles) data source terminals (DST's) provide digitally formatted data into a central processing and display center TACC (Fig. 1) where summarization and data formatting tasks are performed. Data traffic between remote and central locations is minimized by selective satellite (mini) processing of data at the DST's. The satellite processors may have memory capacity of 1/2 million bits each while the central processor has a memory of 6 to 8 million bits. Overall operation involves continuous 24-hour communications data transfers and operator exchanges at keyboard terminals while data are being processed.

Maintenance requirements

Complexity of typical installations involves as many as 120 of the major critical items and 4000 line replaceable units (LRU's) or boards. When combined with the maintainability requirements listed in Table 1 maintenance repair actions of 56 per month for TACC center and 8 per month for each remote DST site are predicted.

The MTTR time limits are generally totally inclusive and account for all of the following actions being successfully completed:

a) Detection and notification to operator that a failure exists.

b) Restructuring of C&C system to an alternate mode of operation with a consequence of reduced processing and throughput capability until failure is corrected.

c) Initiation of fault diagnostics and failure isolation to the specific location of a line replaceable unit (a group of one to five circuit boards).

d) Physical removal of failed unit and insertion of spare.

e) Retest and verification that repair was successfully made.

f) Restoration of overall C&C system to normal operation by returning to a full capability mode.

Because a major portion of the MTTR time limit is consumed in replacing a failed unit, as much automation as possible must be instituted in the failure detection and diagnosis process. However, whatever process is used, it must not compromise the overall C&C system mission and performance capability. The requirements on the C&C test system, therefore, are translated into a number of general objectives:

1) Failure detection and fault isolation processes must have minimum impact on the C&C data processing activity.

2) Software fault detection techniques must complement the hardware detection capabilities so that as high a percentage of system failures as possible is detected by the system in its normal operating modes.

3) Devices and major components such as redundant video terminals and digital data networks (not regularly scheduled by normal system traffic) should be exercised by pre-scheduled self-test programs to assure that no failures exist, confirming total system availability. Fault isolation must be responsive to the skill level of operating and maintenance personnel and to the overall maintainability (MTTR) requirement cited earlier.

Fault detection and diagnostics

An analysis of significant operating parameters and failure characteristics of C&C systems provides guidelines and definitions of the types of fault detection and diagnostics required to meet the maintainability requirements. Referred to as a test requirements analysis, it is performed at both the systems and configuration (major unit) levels, and resulted in the basic test system categories summarized below. A distinction between failure detection and fault isolation (diagnostics) is necessary because of the separation of priorities between the two categories. Failure detection receives foremost priority in the test system design because of the thoroughness and urgency under which it must function. Fault diagnostics, on the other hand (although important) can be a deferred action. A summary of the failure detection and fault isolation functions and definitions follows:

Failure detection

a) Operating System Self Test—Tests automatic functioning of hardware or operating system software, detects failures upon occurrence without supplementary or scheduled tests.

b) On-Line Self-Test—Scheduled tests interleaved with normal system operation on
non-interfering basis. Exercises devices not being continuously used by normal system traffic.

c) Off-Line Self-Test—Tests which require portions of the system to be temporarily removed from normal service. Fairly thorough, capable of fault detection beyond the on-line capabilities. May branch to fault isolation routines on the detection of failure.

d) Manual Observations—System abnormalities detected through operator observations.

Failure isolation (diagnostics)

a) Automatic Execution—Isolation of failed LRU or LRU groups without manual intervention.

b) Combined Automatic and Manual Execution—Diagnostic process requiring manual intervention for evaluation or control.

c) Built In Test Equipment (BITE)—Use of CI internal test capability to minimize dependence on computer programs and AGE.

d) Aerospace Ground Equipment (AGE)—External equipment utilized only for diagnostic purposes—both common and special purpose.

Table II summarizes the distribution of system test categories applied to the various C&C configuration items shown in Fig. 1. The table clearly illustrates, as would be expected, that no single test technique detects and isolates all failures.

The percentages of total faults detected by each category are estimates based on programs such as TIPI at ASD, Burlington, where similar equipments are employed. The LRU group sizes and totals are based on actual board counts. Conclusions from the test requirements analysis and data in Table II indicate that:

1) The highest percentage of failures are detected by the operating system.

2) On-line test techniques also detect a major quantity of failures.

3) Most of the failures are isolated by use of diagnostic programs with combined automatic and manual execution.

4) Nearly all failures are isolated to the printed card or module level.

Test system configuration

The definitions of failure detection and fault diagnostics may be understood more clearly from the flow diagram of Fig. 2.

Under the operating system self-test category of failure detection (Fig. 2) there is always a close association, if not a direct duplication, of processor operating system software with diagnostics. As an example, the error-detection part of diagnostics is run as an integral part of the computer operating system. The scheduled self-test programs are dispatched periodically by the operating system; and, the operating system must properly react to failures reported by these self-tests.

The on-line self-test programs have basically the lowest priority of all operational tasks, since they are intended to exercise devices going for periods of time without activity. The operating system must provide the means to schedule such programs. The interval for scheduling will be at the option of the System Administrator, variable from some fractions of a minute (for some devices), up to intervals of once per hour, or completely cancelled during times of peak system activity.

The operating system must report failure status to the operator. The failure identification may result from normal error detection, result from a scheduled self-test, or be the result of any further testing initiated by the system after the initial indication of a failure; for example, the "fault isolation to the CI" routine is...
One further association between the diagnostic programs and the basic software operating system lies in the reporting of re-try efforts. The operating system will maintain a performance monitor to provide some indication of system performance to the System Administrator/Operator. An input to the performance monitor may be the occurrence of successful re-try efforts, since this could be a measure of system degradation.

All failures which affect C&C systems as a whole are corrected at the organizational level by direct replacement of failed units; the failed unit, or assembly, is further repaired off-line at intermediate or higher levels. Failure detection and diagnostics techniques in C&C systems need only be structured to provide isolation to this level. A considerable reduction in quantity of software diagnostics and test equipment results.

Operating system self-tests

In this category, as stated above, fault detection computer programs are part of the operating systems for both the central processor and satellite mini-processors. The operating system also provides programs for re-try, performance monitoring, scheduling and dispatching of on-line self-test and fault isolation, and for reporting failures to the status panel and operator. The principal types of faults recognized are:

1) interrupts caused by hardware failures,
2) device status words generated by faults occurring during input/output data transfers, and
3) time-outs from an interval timer when an I/O data transfer fails to terminate in a specified time period.

It is logical for the operating system to assume this role because it has the responsibility for processing all interrupts and regulating or controlling all I/O data transfers. Therefore, it must be aware of any errors preventing task completion. Examples of faults detectable by the operating system self-tests are:

1) Central Processor—Machine Check Interrupts—Prime power failures; power supply failures in CPU, IOU, or MMU; parity errors in data and instructions; failures in processing and control hardware.
2) Programmable Control Unit (PCU) or Satellite Mini Processor—Parity error interrupt; time-out interrupt; and illegal instruction interrupt.
3) Peripheral Devices—Parity error status; not-ready status; character failure; interval timer time-out; and message acknowledge failure.
4) Operator Input messages or commands—Format errors.

Fig. 3 illustrates the sequence of events for one typical failure detected in normal operating mode and automatically isolated.

On-line self-tests

This category detects faults by computer
test programs which are periodically scheduled by the operating system. The function of these programs is to write a test message to a device, read it back, and compare the received message with the output message. If the two messages fail to compare, the normal re-try with failure reporting will occur. Since these self-test programs are executed on-line under the control of the operating system and perform normal data transfers, the most likely method for detecting errors results from exercise of the devices in the "operating system self-test" category.

The periodic scheduling of these routines by the operating system involves adjustment with respect to loads and priorities of other system tasks. It is recommended that these routines receive low priorities and infrequent scheduling during periods of peak system activity and, conversely, be scheduled approximately once per minute at other times. The latter will permit failure detection and repair of standby devices in sufficient time to ensure their availability for next peak load. Examples of fault detection by on-line self-test include:

1) Write/Read (Loop) tests which exercise the data processor, PCU's and satellite processors and their peripherals by sending a test message to a device, returning it and comparing the returned data to that transmitted. These test messages may also serve to execute the test functions of operating system self-tests above, thus permitting two classes of testing with a single test message. The data processor loop tests are tabulated below, loops (d)-(h) are illustrated by Fig. 4. Only one loop of each type is illustrated:

a) Data processor-to-drum

b) Data processor-to-interface controller
c) One main memory unit to another (comparison in CPU)
d) PCU-to-disk (Mass Memory)
e) PCU-to-PCU (inshelter)
f) PCU-to-PCU (intershelter)
g) PCU-to-display
h) PCU-to-magnetic tape

2) Operator observation of scheduled test messages to printers, displays, and remote teletypes.

3) Reading of prerecorded test data from storage devices and comparison with expected value. Only used if write operations are prohibited.

4) Intentional interrupt (normal or failure check) and response verification.

5) Periodic exercise of unused communication link by performing dial up and test message transfer.

Off-line self-tests

Fault detection in this category is achieved by removing the processor, controller, and device from normal C&C System processing functions (off-line) and executing a complete test of the controller and/or device. This test differs from on-line self-tests in that it is more thorough. There are certain failure modes which are not detected by on-line self-tests (such as performing data transfers at a slower than normal rate). Since these tests seriously impact normal system operation, they are used only when necessary.

All on-line self-tests have counterpart tests run off-line. The following are examples of unique off-line tests:

1) Write/Read (Loop) tests which exercise the data processor, PCU's and satellite processors and their peripherals by sending a test message to a device, returning it and comparing the returned data to that transmitted. These test messages may also serve to execute the test functions of operating system self-tests above, thus permitting two classes of testing with a single test message. The data processor loop tests are tabulated below, loops (d)-(h) are illustrated by Fig. 4. Only one loop of each type is illustrated:

a) Data processor-to-drum

b) Data processor-to-interface controller
c) One main memory unit to another (comparison in CPU)
d) PCU-to-disk (Mass Memory)
e) PCU-to-PCU (inshelter)
f) PCU-to-PCU (intershelter)
g) PCU-to-display
h) PCU-to-magnetic tape

2) Operator observation of scheduled test messages to printers, displays, and remote teletypes.

3) Reading of prerecorded test data from storage devices and comparison with expected value. Only used if write operations are prohibited.

4) Intentional interrupt (normal or failure check) and response verification.

5) Periodic exercise of unused communication link by performing dial up and test message transfer.
Automatic executed diagnostics (fault isolation)

Fault isolation in this category includes computer programs which automatically isolate the fault to a LRU or LRU group without operator assistance. Consequently, the fault isolation process is fast and reduces MTTR. Computer programs in this category are automatically or manually initiated, and run on-line or off-line. A CI fault isolation program which identifies the failure source to a configuration item (CI) is an example of automatic initiation and running on-line. A main memory fault isolation routine will likely be initiated manually and run off-line. In either case, the computer program isolates the source of the fault without manual intervention.

Since it is uneconomical to completely automate all fault isolation routines, the number of programs in this category is restricted to those considered essential in maintaining system availability and/or those fairly easy to implement such as the single-board controllers of the PCU. In many situations, the MTTR requirement can be met with cost-effective computer programs that require manual intervention.

Combined automatic and manual diagnostic execution

This category includes computer programs which assist the maintenance personnel in performing a fault isolation process. Quite often, these programs must be combined with BITE (built-in test equipment) and AGE (aerospace ground equipment) to isolate faults completely to LRU level. The amount of computer processing aid provided is a function of the complexity of the fault isolation branch.

This category is justified on the basis that some manual intervention is required in many cases where it is not technically feasible to perform completely automated fault isolation. The principal interface between the computer and operator is the teleprinter, providing a continuous interactive exchange of maintenance instructions.

Test results summation

The results of all fault-detection techniques (aside from certain manual observations) are reported to the processing unit within the programmable control unit (PCU); reporting takes place in the following ways:
1) Interrupt from the CE involved,
2) Status word from the CE involved,
3) Test Data comparison within the processing unit (as for a loop test), and
4) Status information entered via keyboard.

In those instances where fault detection alone does not locate a fault to the degree required to light only the proper status indicator, the results of automatic fault isolation routines controlled by the PCU are employed to refine the fault location prior to display generation. The reporting methods used by these routines are the same as 1), 2), and 3) above.

Each PCU has access to a record in memory, determining which peripherals under its control are in "standby" mode. Examples of standby mode are digital-network modems and cryptos which have not been directed to establish a circuit to a remote terminal.

Finally, each PCU receives a report of the "Power Off" status of all equipment under its control. The PCU data processing unit employs all the data to establish the status of each major item (CI or lower) of the subsystem which it controls, including the subdivisions of the PCU itself. Each major item is determined to be one of the four mutually exclusive states.

---

The diagram illustrates the sequence of events for one typical failure detected by an off-line self-check and isolated by a combination of semi-automatic and manual techniques. Off-line tests often branch to fault isolation routines automatically, where computer-directed fault isolation routines are a part of the maintenance plan.

---

Fig. 5 - Off-line self-tests.
Data presenting the established status of each device are transferred from the PCU into the status panel and stored in registers. Each register stage is associated with appropriate lamp drivers and lamps. Multicolored indicators may be used for status of each major item as follows:

- **Power off** (white)
- **Standby** (amber)
- **Failed** (red)
- **Operational** (green)

The panel layout in Fig. 6 is typical. Operating and/or maintenance personnel will be informed by status indications in the following manner:

1. **Power off**—Alerts the operator that a unit has been shut off for maintenance or other reasons, and is therefore inoperative but not necessarily failed.
2. **Failed**—Alerts the operator to a failure so that he will read the printer or CRT output to obtain the identity of the failed unit, or diagnostic instructions. He may also initiate reconfiguration based on a failure indication.
3. **Stand by**—Informs the operator what communication equipment is available but not presently in use. This information may be used for communications subsystem reconfiguration in case of failure.
4. **Operational**—Provides a continuous positive indication of equipment availability.

### Table III — Diagnostic overhead (typical).

<table>
<thead>
<tr>
<th>Test Category</th>
<th>Number of Programs</th>
<th>Total Size (Words)</th>
<th>Size of Largest Program</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Central Processor</strong></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Fault Detection</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>On-line Operating Sys.</td>
<td>3</td>
<td>4,230</td>
<td>1700</td>
</tr>
<tr>
<td>Off-line self-test</td>
<td>5</td>
<td>1,200</td>
<td>500</td>
</tr>
<tr>
<td>Fault Isolation</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Off-line</td>
<td>9</td>
<td>132,000</td>
<td>3000</td>
</tr>
<tr>
<td><strong>Satellite Processor(PCU)</strong></td>
<td>(16 bit)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Fault Detection</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>On-line Operating Sys.</td>
<td>17</td>
<td>3,185</td>
<td>400</td>
</tr>
<tr>
<td>Off-line self-test</td>
<td>18</td>
<td>17,300</td>
<td>2500</td>
</tr>
<tr>
<td>Fault Isolation</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Off-line</td>
<td>(combined with off-line self test)</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

**Software overhead**

The orientation of C&C systems towards data processing favors the use of software as the principal technique for failure detection and diagnostics. However, the overhead or diagnostic burden, which the software imposes, should be properly considered in sizing total C&C system processing and storage capacity. Where a choice or design objective can be maintained, as much diagnostics as possible should be relegated to off-line mass storage devices to be recalled by executive program only when the need arises. Scheduled test programs of outlying peripherals and communication links fit this category. By the same reasoning, as much test information as possible should be extracted between data exchanges in normal system operating modes to reduce special test program requirements.

**Diagnostics**

Table III lists diagnostic program sizing for the typical C&C System and test techniques described in this article. Data are shown for both the main central processor (RCA, R-200 in this case) and a typical programmable control unit or mini satellite processor (Rolm AN/UYK-12). Sizes are given in words, (32-bit for central and 16-bit for satellite). On-line operating system fault detection is the most critical in that it almost solely must remain in core. On the other hand, except for a small percentage of control words, on-line and off-line self-test detection and isolation categories can reside off-line in mass storage.
The use of computer-aided design in the manufacturing of welded-wire circuit boards

M.A. Eastwood and A.S. Baran

Back-plane, welded-wire technology is a packaging concept being used to replace conventional printed-circuit boards. Major reasons for pursuing welded-wire technology are 1) its greater density, which permits elimination of large multilayer printed-circuit boards, and 2) its adaptability to computer-aided design, which reduces drafting and manufacturing costs.

The welded-wire board consists of standard circuit board material with gold-plated pins inserted through it on a standard grid. The pins are made of nickel, Kovar, or rodar and are of two types: one shaped like a flat-headed nail, for accepting IC's; and the other basically similar, but with a post extending from the nail head to accept discrete components. The pins are inserted into the board and, in some cases, are connected to conductive planes by use of solder preforms and a hot-oil wave machine. Once the pins have been inserted into the board, the board is subjected to a welding process in which a teflon-insulated nickel wire is welded through the insulation onto the side of the inserted pins. Circuit networks ("nets") are developed by interconnecting a series of pins in a contiguous manner with one piece of wire. A single pin may contain as many as three individually welded wires that fan out to other pins as well as connectors.

Why AED uses welded wire

There are a number of advantages in using back-plane, welded-wire boards over multilayer printed-circuit boards:

1) The cost of board assemblies and testing is reduced. Drafting time is reduced due to the use of computer-aided design (CAD). All assembly and test operations are computer-controlled or numerically controlled, except for pin insertion and the welding of components to the boards.

2) All welded connections are visible, permitting 100% inspection.

3) Higher density packaging permits twice as many circuits to be placed on a welded-wire board than on a multilayer printed-circuit board of the same size.

4) The high-density circuit packaging permits reduced "black-box" volume of the end product.

5) The single board used with back-plane, welded-wire technology permits easy access for design changes and repairs, as no inner board layers are present.

6) The modular concept allows for easier growth.

7) Reliability is increased because wire-welded joints are typically equal to or greater than 85% of the tensile strength of the wire. High reliability is obtained because the entire board configuration is welded, including through-the-insulation, back-plane wiring, as well as the welding of IC's and discrete components into the board assemblies.

Manufacturing welded-wire boards

A typical flow for fabricating a welded-wire back-plane board is shown in Fig. 1. It starts with the drilling of hole patterns in the board to accept weld pins. This drill operation is accomplished on a numerically controlled drill machine, shown in Fig. 2. The raw-stock material is drilled to a desired pinfield pattern depending upon the inputs specified on the punched paper tape controlling the numerically controlled drill. Once this pattern has been completely drilled, the board is removed for routing of the outside dimensions to that prescribed on the drawing. If the board contains one or two conductive layers (voltage plane and ground plane), then the next step in the flow is to etch out those points on the conductive planes that are not to be connected to inserted pins. The third step in the flow is to mark the board pin and component identifications using a layout negative. The board is then cycled through a semi-automata pin insertion machine and the appropriate pins (either for IC's or discretes) are inserted into the board. Those boards having some pins connected to conductive planes are then fitted with solder preforms and passed through a hot-oil wave bath to reflow the solder, producing a connection between the pin and the plane.

The board is now ready to proceed to the
Anthony S. Baran, Mgr., Advanced Manufacturing Development facility, Astro-Electronics Division, Princeton, New Jersey, received the BSEE in 1958 from Pennsylvania State University where he was also employed for 2 years on a U.S. Army Quartermaster Corps program studying shock impact to electronics equipment. From 1964 to 1967, he performed graduate study in electrical engineering at Newark College of Engineering. Before joining RCA, Mr. Baran was employed as a Senior Engineer in instrumentation at Pratt and Whitney Aircraft's Research and Development Center, West Palm Beach, Florida, where he later managed a vibration and acoustic instrumentation group supporting experimental jet and rocket engine testing. At AED, he was assigned as Lead Engineer for Minuteman vibration testing and later as Instrumentation Engineer in the Environmental Simulation Group, where he participated in the testing of TIROS, Relay, Ranger, and other satellites. More recently, he has been involved in implementing automatic testing and computer-based test stations. He presently has responsibilities for long-range planning and technical guidance in advanced manufacturing technology and industrial engineering services. Mr. Baran is a member of the IEEE and the ARPL.

Margaret A. Eastwood,* Computer-Aided Design facility, Astro-Electronics Division, Princeton, New Jersey, joined the Scientific and Engineering Programming group at AED in 1966, after receiving the BS in systems analysis from Miami University, Oxford, Ohio. In 1973, she received the MS in computer science from Rutgers University. Previous to her work in computer-aided design, she was responsible for spacecraft dynamics and attitude control simulation programs. She also wrote minicomputer programs to perform real-time computer testing of the 2-inch RBV television camera and to do real-time ranging calculations for navigation satellites. As lead programmer for the welded-wire project, she was responsible for designing the current system. She continued in that capacity to begin the development of an expanded system, which will also support printed circuit boards and circuit analysis.

*Since this paper was written, Mrs. Eastwood has left RCA.

In the welding operation, teflon-coated nickel wire is fed by an operator and placed between the pincer electrodes and the pin to receive the wire. The operator controls the Z-axis motion, lowering the head over the pin through the use of a foot-control pedal. Once the head is in the down position and properly oriented with the pin and wire, the pincers are closed, producing a cold flow through the wire insulation and a resistance weld of the wire to the pin. When the operator again enables the system, the tape reader reads instructions from the next block of data and moves the table to the appropriate X and Y coordinates to produce the next weld. Again, with operator intervention, the head is lowered in the Z-axis and a through-the-insulation weld made. This process continues until an audible alarm alerts the operator that the net being welded is complete. At this point, the wire is cut, and a new net is started. To signify to the operator which color wire is to be used in the forthcoming net, a command is received from punched paper tape that numerically controlled welder system for welding the wires onto the pins in the board back-plane. The welder, shown in Fig. 3, is a Hughes Model 1020 wire termination system capable of performing through-the-insulation welding, wire wrap, and Termi-point type connections. The machine is controlled from punched paper tape prepared earlier through the CAD program.
lights the appropriate indicator in the wire bin. The X-Y table, under the control of punched paper tape, continues to move through the coordinates, picking up all of the pins specified in the weld net list. If a heavier terminal or weld pin is encountered in a net, an instruction is received from the punched paper tape that automatically changes the weld schedule for welding the new combination of wire to pin. The process continues until the back-plane is completely wired. The welded-wire back-plane is now ready to proceed to the next stage in the flow, which consists of testing the wired nets for continuity and discontinuity.

The equipment used for testing the back-plane wiring is a DIT-MCO Model 660B Circuit Continuity Tester (CCT), shown in Fig. 4. This equipment essentially provides 100% continuity-discontinuity testing of all wires (nets) on the board. Simply stated, the continuity test is one that checks for continuity between all pins in a given net, and the discontinuity test checks all other pins against a specified net to assure that accidental connections (short circuits) have not occurred. Basically the system consists of 5000 relays that are controlled from punched paper tape. Each pin on the board is interfaced with the CCT through a spring-loaded probe located in the test fixture that couples the CCT with the board being tested. The test item is placed over the spring-loaded test fixture and pneumatically engaged to produce contacts between all pins and corresponding test probes. Each test probe is connected to a separate relay that receives enable-disable instructions from punched paper tape.

Perhaps the single most significant factor using this testing technique is the ability to perform 100% continuity and discontinuity testing. To test a board with 1000 pins for full continuity and discontinuity manually would require nearly half a million independent measurements as approximated from the equation \( n(n-1)/2 \) where \( n \) is the number of independent measurements. Because of the high test rate of the automated method, it is also cost-effective to resubmit repaired boards to a complete continuity-discontinuity test to provide assurance that the welded back-plane is totally correct. After the board has cycled through the test and is validated as being 100% correct, it then is moved to the parallel gap welder stations to have components welded on the other side, followed by conformal coating of the welded-wire back-plane.

The final operation after assembly of the board is complete is performance of a functional and parametric test on the board, using a computer-based test system shown in Fig. 5. The Manufacturing Automated Test System (MATS) employs a mini-computer, a higher-level test language, and is capable of applying various stimuli to the unit under test as well as monitoring the output conditions of the board being tested.

Improvements to the manufacturing process are continually being sought. For example, a time-shared computer-controlled weld station to replace the numerically controlled welder (for welding wires to pins) is now under investigation. This computer-controlled welder uses a coaxial weld head that requires less operator intervention, is faster than the numerically controlled welder, can operate 10 other slave welders, and can perform 100% continuity-discontinuity testing.
stations on a time-shared basis, and produces a more favorable welded backplane that is less prone to accidental shorts resulting from cold flow of insulation. Other areas under surveillance are adoption of a standard pin grid to reduce test fixture costs; development of standard size boards and more connector types and special components such as relays; and the use of automated methods for insertion of pins.

Computer-aided design software

The production of welded wire boards at AED is supported by a series of computer programs. All these programs are written in Fortran IV and are executed on a UNIVAC Series 70/45 using only 230 kilobytes of memory. The programs were written to meet three objectives. Some programs produce the paper tapes to drive the numerically controlled equipment. Some produce error printouts and reports, assisting manual inspection or changes to the boards. All programs interact with a central program which monitors the progress of every board through the manufacturing process.

The computer programs are executed in a sequence paralleling the stages of production. A simplified flow diagram is shown in Fig. 6. The first operation in the flow is recording of the physical placement of the connectors, IC's, and discrete components. Each package type is associated with a physical code that indicates the number of pins and the pin spacing. Therefore, to record the physical layout, the X and Y coordinates of one pin (the "key pin") and a physical code are punched for each component.

Drill program

These key-pin list cards are the input to the drill program. The program calculates the coordinates of all the non-key pins, and punches the paper tape for the numerically controlled drill after separating the grounded from the non-grounded pins.

This one "drill" tape has several functions. The entire tape is used to control the drilling of the component board. The non-grounded pin portion of the tape is used to drill a photo-mask board for etching the copper from the ground plane. A third, heavier board is drilled to use as a guide during pin insertion. In addition, a wire-mask board is made, which will later protect the wiring from probes during continuity-discontinuity testing.

Network optimization

After the layout has been specified, the connectivity (specified interconnections) for each board must also be prepared manually and keypunched. This information is referred to as a wire net list. A sample list is shown in Fig. 7.

The input card deck, containing the nets and their signal names, is first run through a format check program. In addition to detecting certain key-punch errors and missing data, the format check program also gives two important printouts. The first listing indicates those nets connected by a common pin. Pins inadvertently wired to ground are often located in this listing. The more common use of this listing is to compress two connected nets into one longer net. After the individual nets are optimized, all the nets are sorted so that those with the longest wire segments will be welded last in the net. Nets that start and end on the same pin, common in some of our decoder logic, also require special handling.

After the individual nets are optimized, all the nets are sorted so that those with the longest wire segments will be welded first. By processing the nets in this order, long wire runs will be tied down by the shorter runs above them. The final optimized arrangement is saved for use by other programs.

The continuity-discontinuity test tape cannot be produced immediately, however. The fixture which will be used to interface the board and tester must be determined first. Since the fixtures are very expensive, it is desirable to use one fixture to test many different boards. The cumulative probe pattern corresponding to all the boards is prepared by the fixture.
composite program, using the individual component placement files as input. The composite program sorts the points, removes redundant points, and assigns a particular DIT-MCO CCT connector and pin to each probe location.

**Fixture drill and weld program**

The fixture drill and weld program is run next. It punches a paper tape to drill the cumulative hole pattern in the fixture probe plate. It also produces the paper tape to wire the specified probes to the specified CCT connectors and pins.

With the fixture specified, the paper tape to perform the actual continuity-discontinuity test can be made. Each net in the wire net list must be converted from unit-under-test pins to CCT pins. Each net is punched out on paper tape in ascending CCT pin order. Nets connected by a common pin are tested together. Nets formed by latching relays in both the harness board is also fabricated by using the welded-wire programs.

**Harness net list program**

After the harness board has been drilled (using the drill program previously described), the necessary wiring is determined automatically by the harness net list program. It scans each individual component board net list, searching for nets which terminate at a harness connector. The signal names and connector pin numbers of those nets are saved. At the completion of that search, the harness net list is compiled by assigning all connector pins with the same signal name to the same net.

Next the harness board is welded, using the regular weld program to prepare the paper tape.

Harness board continuity-discontinuity testing is different from component board testing, because prewired connector cables rather than a fixture are used as the interface to the tester. Consequently, a harness board test program determines the correspondence between unit-under-test pins and cable pins. It then punches the paper tape to drive the continuity-discontinuity tester.

**Master status program**

There is one more computer program in our welded wire system—the master status program. It monitors the progress of every board by maintaining a status file for it on the disk.

This status file is updated by each of the other computer programs every time they are executed. They log the run date, the board affected, the new status, and the revision letter. Other status changes, not involving a computer program, are logged by using punched cards. When the disk file is initialized, the master status program punches card decks to be given to the drill shop, weld shop, etc. As a board successfully passes through one of the production stations, its punched card is returned to the computer room and used to update the status file.

The master status program also prints various status reports. One of them is shown in Fig. 8. Another report lists the latest revision letters for all the numerically controlled paper tapes. All paper tapes begin and end with a block letter label (see Fig. 9) to facilitate identification of the proper tape.

The future

Additional improvements are planned for the software. In conjunction with other RCA divisions, the current programs are being interfaced to a common data base, and several new programs are anticipated. A digitizer is also being purchased to greatly reduce the amount of keypunching required for the input data.
Self-calibration of land-based pulse radar from satellite track

J. J. O’Connor | R. R. Rowe

C-band pulse radars measuring slant range, azimuth, and elevation have constituted the major missile and satellite tracking system on the Air Force Eastern Test Range (AFETR) since the beginning of the range. To serve their purpose, these radars must be constantly maintained in a highly accurate state of readiness—a condition which can be assured only if accurate and timely calibration methods are applied. This paper does not discuss all the current calibration techniques, which include mechanical, optical and electronic processes and also multiple tracker comparisons; rather, attention is focused on the one procedure primarily responsible for the present high accuracy level; self-calibration from satellite passive track by a single radar.

The present accuracy of the pulse radars represents a factor-of-four improvement over that existing before introduction of the self-calibration method. With such accuracy, the radars are now adequate for measuring certain drift characteristics in missile guidance systems and for determining detailed aerodynamic coefficients of reentering missiles.

With the self-calibration method, a single radar is calibrated without the use of any other tracking system or measuring device. The radar’s own passive tracking data from two successive passes of a satellite is used in conjunction with the theoretical equations of motion, known physical constants, and associated computer program. Although the computer used in radar calibration is a CDC-3600, it is believed that with careful programming, a computer with floating-point and double-precision capability and with only 16 k-word memory would be adequate.

To ensure proper application of the self-calibration exercise, a formal Radar Accuracy Monitoring Program (RAMP) has been instituted at the Air Force Eastern Test Range. In this program, a network of radars extending from Cape Kennedy to Ascension Island is served by a central computer facility at Cape Kennedy. With good operating procedures, calibration at about five-day intervals is sufficient to maintain accuracy, unless a hardware change is made. Calibration is scheduled immediately after a hardware change.

Satellite selection

Satellites for use in calibration are selected based on density, radar cross-section, altitude and inclination angle. Generally, any payload except a balloon is sufficiently dense to avoid radiation pressure effects and excess drag. Large objects are desirable to provide good signal strength. Altitudes over 200 n.m. are desirable to avoid excess drag, but much higher altitudes must be avoided to maintain good auto track from horizon to horizon. Orbital inclination angles slightly larger than the latitude of the radar site are desirable to obtain visibility and good geometrical conditions. Presently, either Pegasus I (Object 1085) or Pegasus II (Object 1381) is used, as convenient. Ephemerides for these satellites are maintained at the central computer facility so that calibration tests can be scheduled.

Reprint 19-4-12

Final manuscript received November 13, 1973

John J. O’Connor, Technical Analysis, Missile Test Project, Patrick Air Force Base, FL, received the BS from Western Kentucky University in 1939, the MS from Vanderbilt University in 1940, and did two years of graduate work at Ohio State University, all in physics. He was elected vice-president of Ohio State Chapter of the National Honorary Physics Society. Employed by the DuPont Company from 1942 to 1961 in the capacities of Research Physicist, Senior Research Physicist and Research Associates, he worked on the atomic bomb project and on accuracy of spinning projectiles and also represented DuPont on the technical committee of an inter-industry council. From 1961 to present he has been employed by RCA at Patrick Air Force Base with primary emphasis on problems associated with missile and satellite trajectory computations, and by Florida Institute of Technology as Adjunct Professor in Celestial Mechanics. He has several publications in Physical Review on topics in magneto-optics and nuclear physics and is recognized as the discoverer of two radioactive isotopes. He is a member of the American Physical Society and the American Astronomical Society.

Robert R. Rowe, Technical Analysis, Missile Test Project, Patrick Air Force Base, FL, received the BS in Mathematics from Delta State College in 1962 and joined Dover Corporation Elevator Division. He did graduate work at University of Mississippi before coming to RCA MTP Data Processing in 1963. In the latter organization, he was responsible for planning, reduction and analysis of both optical and electronic tracking data. In 1968 he joined Technical Analysis and has been active in the design of both orbital and reentry trajectory computational techniques and instrumentation calibration techniques. He is a past Regional Officer in The American Society of Photogrammetry and The Society of Photo-Optical Instrumentation Engineers.

| Table I — Percentage frequency distribution of measured zero-set errors. |
|-------------------|-------------------|
| Error interval     | Azimuth zero-set  |
| (mil)             |       | Elevation zero-set |
|                   |      |                   |
| -0.05 to +0.05    | 67    | 53                |
| -0.10 to +0.10    | 94    | 86                |
| -0.15 to +0.15    | 92    | 96                |
| -0.20 to +0.20    | 100   | 98                |
| -0.25 to +0.25    | 100   | 99                |

<table>
<thead>
<tr>
<th>Error interval (m)</th>
<th>Range zero-set</th>
</tr>
</thead>
<tbody>
<tr>
<td>-15 to +15</td>
<td>65</td>
</tr>
<tr>
<td>-30 to +30</td>
<td>88</td>
</tr>
<tr>
<td>-45 to +45</td>
<td>96</td>
</tr>
<tr>
<td>-60 to +60</td>
<td>98</td>
</tr>
<tr>
<td>-75 to +75</td>
<td>98</td>
</tr>
</tbody>
</table>
and acquisition information supplied to the radar site in advance. Each exercise requires tracking data from two successive revolutions of the satellite such that the passes are on opposite sides of the radar being calibrated and have maximum elevation angles between 20° and 85°. (Passes with maximum elevation below 20° contain insufficient information, and passes with maximum elevation above 85° may result in failure of the radar to maintain autotrack.)

Calibration data

Track data used in the calibration exercise consist of range, azimuth and elevation measurements at one point per second, corrected at the radar site for refraction, transit time, dynamic lag, and any other known errors from previous calibration runs. Data points with elevation angles below 5° are discarded because of possibility of excessive residual refraction errors and multipath. The resulting track data from both revolutions are forwarded to the central computer facility and fitted in a single arc in a standard orbital computer program, which includes least squares adjustment of the parameters describing the satellite orbit plus various radar-error coefficients. If a significant error is detected in the radar data, a revised coefficient is transmitted to the radar site to use in future on-site data correction. Of course, other action is taken not specifically part of the calibration exercise. The entire exercise from first track to computer output can be completed in a few hours elapsed time and less than one hour computer running time on a CDC-3600 computer.

The complete set of radar-error coefficients involved in calibration number around thirty, some of which are rarely or never estimated in a self-calibration exercise. The major radar error coefficients are the zero off-sets or biases in range, azimuth, and elevation; these are nearly always included in a self-calibration exercise. Table I shows a percentage frequency distribution representing a combined history of the zero off-sets for four radars over a recent six-month period. The tabular values represent the errors, determined at the central computer facility, in the "corrected" data as obtained from the

Definitions

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>Radar azimuth angle. Defined by two vectors lying in the local horizontal plane and originating at the radar gimbals. The reference vector points northward in the plane, and the other vector is a projection of the radius vector extending toward the tracked object. Angle is positive eastward from north.</td>
</tr>
<tr>
<td>C rer</td>
<td>Dimensionless spherical harmonic coefficient in geopotential expression having degree n and order m.</td>
</tr>
<tr>
<td>D</td>
<td>Drag parameter of satellite, equal to product of drag coefficient and cross-sectional area divided by twice the mass.</td>
</tr>
<tr>
<td>E</td>
<td>Radar elevation angle. Measured positively in vertical plane from local horizontal plane upward to radius vector to tracked object.</td>
</tr>
<tr>
<td>F</td>
<td>Column vector of corrections to adjusted parameters.</td>
</tr>
<tr>
<td>G</td>
<td>Column vector of measurement residuals, defined as actual minus computed, for the ( i ) time point.</td>
</tr>
<tr>
<td>H</td>
<td>Column vector of differences between current and ( a ) priori estimates of the parameters defined as current minus ( a ) priori.</td>
</tr>
<tr>
<td>[K] jxj</td>
<td>( a ) priori covariance matrix of parameters to be adjusted, usually diagonal.</td>
</tr>
<tr>
<td>[M] jx3</td>
<td>Covariance matrix of measurements at ( i ) time point, usually diagonal.</td>
</tr>
<tr>
<td>P m (sin ( \phi ))</td>
<td>Associated Legendre polynomial of degree ( n ) and order ( m ) with argument sin ( \phi ).</td>
</tr>
<tr>
<td>( \Omega ) jx3</td>
<td>Matrix of partial derivatives of the measurements with respect to the adjusted parameters for the ( i ) time point.</td>
</tr>
<tr>
<td>R</td>
<td>Radar range. Distance from radar gimbals to tracked object.</td>
</tr>
<tr>
<td>S nm</td>
<td>Dimensionless spherical harmonic coefficient in geopotential expression having degree ( n ) and order ( m ).</td>
</tr>
<tr>
<td>T</td>
<td>Used as superscript to denote matrix transpose.</td>
</tr>
<tr>
<td>U</td>
<td>Earth gravitation potential.</td>
</tr>
<tr>
<td>V</td>
<td>Satellite velocity (scalar).</td>
</tr>
<tr>
<td>X, Y, Z</td>
<td>Position coordinates of satellite in a geocentric, Earth-fixed, rectangular right-handed coordinate system with ( X ) and ( Y ) axes in equatorial plane, ( X ) extending through meridian of Greenwich and ( Y ) through 90° east longitude. ( Z ) points northward along spin axis.</td>
</tr>
<tr>
<td>X', Y', Z'</td>
<td>Velocity components of satellite in ( X, Y, Z ) coordinate system.</td>
</tr>
<tr>
<td>X'', Y'', Z''</td>
<td>Acceleration components of satellite in ( X, Y, Z ) coordinate system.</td>
</tr>
<tr>
<td>( a_{x}, a_{y}, a_{z}, a_{v}, e_{x}, e_{y}, e_{z} )</td>
<td>Radar azimuth-error coefficients defined by radar error equation.</td>
</tr>
<tr>
<td>( e_{i}, e_{j}, e_{k} )</td>
<td>Radar elevation-error coefficients defined by radar error equation.</td>
</tr>
<tr>
<td>( i, j, k )</td>
<td>Time point index.</td>
</tr>
<tr>
<td>( n, m )</td>
<td>Number of adjusted parameters including those which describe the trajectory and the radar error coefficients.</td>
</tr>
<tr>
<td>( p )</td>
<td>Degree and order respectively in spherical harmonic expression for geopotential.</td>
</tr>
<tr>
<td>( q )</td>
<td>Total number of time points at which radar measurements are taken.</td>
</tr>
<tr>
<td>( r_{x}, r_{z} )</td>
<td>Earth equatorial radius.</td>
</tr>
<tr>
<td>( \phi )</td>
<td>Radar range-error coefficients defined by radar error equation.</td>
</tr>
<tr>
<td>( s )</td>
<td>Distance from center of Earth to satellite.</td>
</tr>
<tr>
<td>( t )</td>
<td>Subscript signifying &quot;true.&quot;</td>
</tr>
<tr>
<td>( u, v )</td>
<td>Radar mislevel error coefficients; ( u ), north; ( v ), east.</td>
</tr>
<tr>
<td>( \lambda )</td>
<td>Random error.</td>
</tr>
<tr>
<td>( \mu )</td>
<td>East longitude from Greenwich.</td>
</tr>
<tr>
<td>( \rho )</td>
<td>Earth gravitational constant, equal to product of Newton's gravitational constant and the Earth mass.</td>
</tr>
<tr>
<td>( \lambda )</td>
<td>Atmosphere density.</td>
</tr>
<tr>
<td>( \phi )</td>
<td>Geocentric latitude.</td>
</tr>
<tr>
<td>( \omega )</td>
<td>Earth rotation rate.</td>
</tr>
<tr>
<td>( -1 )</td>
<td>Used as superscript to denote matrix inverse.</td>
</tr>
</tbody>
</table>
The gravitational model is defined by the experimentally determined values for \( C \) and \( S $. Any of the newer models issued by Naval Weapons Laboratory, Smithsonian Astrophysical Observatory, Applied Physics Laboratory, or NASA (Goddard) is satisfactory for this application and contributes no significant error to the calibration. Any of the newer atmospheric models issued by NASA, Smithsonian Astrophysical Observatory, or U.S. Air Force may be used to describe the air density \( p \). Consequently, given a vector \((X, Y, Z, X', Y', Z')\) and the various models and constants entering into the equations of motion, one can numerically integrate the equations of motion (Eq. 4) and thereby generate a satellite orbit for as many revolutions as necessary. For this exercise, never more than two are needed.

The only orbital parameters to be estimated in the calibration exercise are the components of the vector \((X, Y, Z, X', Y', Z')\) at an epoch corresponding to first track point. The constant \( D \) cannot be estimated from two passes of the satellite and hence must be determined beforehand using three or more passes of track data in an adjustment process similar to the one to be described here. In practice, there is always a value for \( D \) available from ephemeris-maintenance activity at the central computer facility.

The adjustment process

The object is to make a simultaneous estimate of the orbital parameters and the radar-error coefficients. The mathematical procedure is iterative and is based upon the criterion of minimization of the sum of squares of the weighted measurement residuals. A weighted measurement residual (dimensionless) is simply a measurement residual divided by the a priori estimate of standard deviation in random error in that particular measurement. Measurement residuals are obtained as follows:

An initial estimate of a vector \((X, Y, Z, X', Y', Z')\) at the first track point may be obtained by a simple curve fitting and numerical differentiation process on a short span of actual \( R, A, E, \) track data transformed to \( X, Y, Z \). An orbit is generated by numerical integration of equations of motion. Then by a coordinate transformation, the resulting \( X, Y, Z \)
Z time points are transformed to theoretical $R, A, E$ measurements for the radar site. These theoretical measurements are used in place of the true but unknown measurements on the right side of the radar measurement equations (Eqs. 1, 2, 3) to produce "computed" measurements. Random errors are never included in the computer measurements, and on the first iteration all other error coefficients are set to zero since they are unknown and have expectations of zero. After the first iteration there are nonzero estimates of the radar-error coefficients to use in producing computed measurements. The differences between the computed measurements and the actual measurements produce measurement residuals.

The measurement residuals from the first computation will generally be much larger than the random errors in the actual radar measurements and thus indicate that one or more of the adjusted parameters have been poorly estimated. Corrections are therefore computed and added to the original estimates of all the adjustable parameters. Corrections are computed using a basic least-squares iterative equation satisfying the original minimization criterion:

$$ F = \frac{1}{N} \{ [Q]^T [M]^T [Q] + [K]^{-1} \}^{-1} \{ [Q]^T [M]^T G \} - [K]^{-1} H \} \ (5) $$

$F$ is a column vector of corrections for the adjusted parameters, including orbital parameters and radar-error coefficients. Eq. 5 is used together with the equations of motion (Eq. 4) and the radar measurement equations (Eqs. 1, 2, 3) like any other set of iterative equations. After a number of iterations, reductions in the sum of squares of the weighted measurement residuals become negligible, at which time the adjustment process is said to have converged, producing final estimates of all parameters. To give some idea of the accuracy in the self-calibration method, the results of a simulation exercise are presented in the next section.

Simulation exercise

Selection of the particular coefficients to be estimated in an actual calibration exercise depends upon some practical considerations that are discussed in the next section. The following treatment illustrates a possible calibration exercise.

Simulated data at one point per second for a radar on Grand Turk Island were generated assuming a Pegasus trajectory and three types of errors:

- **Unadjusted parameters** were represented by conservative estimates of errors in geopotential, survey, and drag;
- **Adjustable parameters** were represented by zero-set, mislevel, nonorthogonality, collimation, velocity lag, acceleration lag, and droop;
- **Random errors** (i.e., standard deviations) of 21 ft in range, 0.1 mrad in azimuth and 0.12 mrad in elevation were also introduced.

The simulated error coefficients, the error coefficients estimated from the adjustment process and the differences are shown in Table II. *A priori* estimates for all error coefficients were zero.

```latex
\begin{table}[h]
\centering
\begin{tabular}{|c|c|c|c|}
\hline
| Error coefficients | Introduced value | Estimated value | Error in estimate |
\hline
$\tau$ (ft) & +21.0 & +38.5 & +17.5 \\
$a$ (mrad) & + 0.070 & + 0.006 & + 0.010 \\
$b$ (mrad) & + 0.070 & + 0.095 & + 0.011 \\
$c$ (mrad) & + 0.070 & + 0.079 & + 0.089 \\
$d$ (mrad) & + 0.00276 & + 0.00239 & - 0.00031 \\
$e$ (s/ft) & + 0.0590 & + 0.0613 & + 0.0023 \\
$f$ (mrad) & + 0.070 & + 0.0245 & - 0.0025 \\
$g$ (mrad) & - 0.0640 & - 0.0547 & + 0.0090 \\
$h$ (mrad) & + 0.070 & + 0.091 & + 0.021 \\
$i$ (mrad) & + 0.070 & + 0.053 & - 0.017 \\
\hline
\end{tabular}
\caption{Numerical results from simulation.}
\end{table}
```

Some practical considerations

In an ordinary satellite pass, the cosine $E$ and secant $E$ functions do not vary much, and therefore the corresponding droop and collimation coefficients can not be readily distinguished from the constant elevation and azimuth zero sets in the calibration exercise — hence very high maximum elevations are desirable for estimating droop and collimation. In the simulation just presented, one pass had a maximum elevation of 83°. When ideal conditions are not present, it is good practice to exclude droop and collimation from the calibration exercise and let their residual effects be absorbed in the estimates of elevation and azimuth zero sets. The lag coefficients $a$ and $as$ also cannot be estimated unless a high maximum elevation pass is available, because only in such a situation is the lag significant.

If the full set of error coefficients is adjusted, it will be found that even under good geometrical conditions high correlation will be observed between elevation bias and droop, and also between azimuth bias and collimation. Methods for treating nonorthogonal problems of this type have been widely discussed in the statistics literature and investigated at length at this Test Range. Refs. 2 and 3 give particularly valuable treatment. It is well known that adjustment errors caused by correlation among the adjusted parameters are associated with increased magnitude of the final value of the vector $H$. If the magnitude of this vector is properly restrained, the adjustments are frequently improved. A restraining influence on the size of this vector is the covariance matrix $[K]$. When correlation problems exist, it is frequently profitable to arrive at an approximately realistic estimate for $[K]$ and then multiply $[K]$ or a subset of $[K]$ by a fraction (e.g., 1/3 or 1/4) before beginning the adjustment process. Such a procedure results in a statistically biased but nevertheless more accurate estimate of the adjusted parameters. In the simulation exercise just discussed, a realistic value of $[K]$ was multiplied by 1/4 before proceeding with the adjustment. In a simple calibration exercise involving only zero-set estimations, there is no problem with correlation; and the inverse of the $[K]$ matrix may be defined to be a null matrix to simplify computations.

Finally, it should be pointed out that random errors in the radar data have a harmful influence on the accuracy of the adjustments. The random errors used in the simulation here represent average values at this Test Range. A 50% change in these errors is significant in terms of adjustment accuracy. Effects from random errors can sometimes be reduced by using input radar data at ten points per second instead of one point per second but only at considerable additional expense in computer time. Prefiltering ten-point-per-second data to obtain less noisy one-point-per-second data is also a possible source of improvement that may be exploited in the future.

References

Videovoice — an update

S.N. Friedman

Two-way television communication devices similar to telephones have long held the imagination of the public. The Videovoice system is the first currently available all-electronic system which can be used in conjunction with existing telephone lines and switching equipment. The commercial installations, both domestic and international described in this paper verify earlier expectations that the addition of video to voice over the telephone circuits provides a significant enhancement of communications capability. This paper is an update of the Videovoice article previously published in the RCA Engineer, Vol. 17, No. 2. A new section on present customer usage has been added.

In this day and age, with television impacting so dramatically on our personal lives, there is a growing awareness in the business and professional communities of the potentially very broad benefits to be derived from effectively utilizing video and imagery transmission technology in communications.

Because of the general unavailability of economical two-way video transmission facilities, conventional television has been accepted as a "one-way" communication system, where individuals have their own receivers, but transmission is restricted to major broadcasting facilities. The absence of a "two-way" capability restricts the individual to the role of a viewer and precludes the exchange of information which is vital to total communication. Although a picture is recognized to be "worth a thousand words," up to now we have been unable to use our vast television know-how and facilities to the full advantage of the communicator.

Videovoice defined

Videovoice (see Fig. 1) is a slow-scan TV system designed to include TV-type picture information as part of a duplex communications link, using standard voice-bandwidth facilities for transmission of video and standard TV monitors for viewing.

A brief comparison of Videovoice and conventional TV highlights some of the performance characteristics of the slow-scan version. In standard TV, 30 full pictures, or frames, are transmitted every second; that is, a full picture is displayed on the screen and then erased, and the next picture is displayed in its place. This sequential display of full frames at 30 frames per second produces the illusion of motion in the picture. The price we pay for this illusion is the amount of spectrum space required for transmission. Video signals typically fill a bandwidth of 4 MHz, more than 1000 times the bandwidth of a voice circuit and, in addition, pose problems of transmission and recording.

The visual telephone system presently under development by the telephone company appears, at first impression, to be a video system utilizing standard telephone lines. Actually it requires a transmission bandwidth of 1 MHz: more than 300 times the capacity of the standard telephone line. This means that new cables and switching system would be required domestically, and that international use of these systems could be prohibitive.

On the other hand, Videovoice is completely compatible with standard voice-grade circuits and, in addition, allows the user to hold a single desired frame of video; in regular TV, a moving sequence would pass on and be lost. It is this one-frame-at-a-time transmission capability that allows us to achieve video transmission by telephone line. We can, at the same time, realize additional benefits.

Fig. 1 — TV picture transmission over a telephone circuit by Videovoice.
from the conversion of video to the audio range: video information can be stored on standard audio tape; voice and video can be recorded on stereo tape; and video can be transmitted over standard voice-grade facilities alternating with voice, data, and teletype.

Videovoice functions

- Videovoice can transmit standard television signals over a telephone network or over a conditioned (FCC specifications for assigned voice-grade telephone lines) voice-grade circuit. The system can be used to photograph a subject with a television camera operating at the standard TV scanning rate, convert to a much lower rate for transmission over voice-bandwidth circuits, store the received picture at the slow rate, and then present it for display on a standard TV monitor screen.

- Videovoice signals are also compatible with closed-circuit TV; therefore they can be retransmitted at the "receive" terminal over a local closed-circuit TV system.

Videovoice does not present motion in the display. Many video requirements, however, can be met without the need for instantaneous transmission, and many users do not require motion in the received display. For these users, Videovoice fills a very important communications need: that of providing a video system that is compatible with standard television as well as standard telephone lines.

The Videovoice transmission process takes between 30 and 55 seconds for each frame of video, depending on the amount of detail desired. Still subjects such as equipment, objects, documents, charts, blackboard information, and scenic views, can be transmitted with a "live" TV camera; that is, the actual subject material, rather than the stored frame, is scanned during the transmission period.

- Subjects in motion must have the motion stopped for the duration of the transmission period. To accomplish this, a "frame-freeze" unit stops the action in a single TV frame period and holds that frame for transmission. When a new frame is desired, the "frame-freeze" unit stores another frame, erasing the preceding frame. This frame-freeze capability is made possible by the novel silicon target storage tube developed at the RCA Laboratories, Princeton, N.J. This tube stores the information electronically, allowing it to be read off at the specified speed. At the distant location, the receive terminal is automatically activated by the incoming video signal and the picture can be viewed for as long as desired or until the presentation of the succeeding frame.

Because Videovoice provides for long-time retention of the video at both the transmit and receive terminals, it is not necessary to keep the long-distance circuit "up". Once the transmission is completed, both telephones can "hang-up" and the subject information will still be retained for viewing at both terminals.

System description

The basic Videovoice system is shown in Fig. 2. The TV camera and monitor are mounted in a custom-designed desk-top unit which allows full rotation of the camera so the "executive" operator can select subjects in the nearby area. A special 90-degree reflecting device is incorporated in the camera assembly to allow focusing on documents and small objects placed on top of the desk, without tilting the camera. No familiar TV controls are required.

The same monitor is switched to display the received picture in a bi-directional half-duplex system. In a simplex system, the monitor is available as a separate unit, without camera, for receive only. Also available is an RCA large-screen TV set which allows for convenient viewing by a number of people; other sizes or types of standard television monitors can be used to suit individual requirements.

Each desk-top assembly contains a speaker (within the monitor) and a microphone (in the telephone) providing hands-free operation of the Videovoice system and independent use of the telephone handset. The balance of the electronic equipment is furnished in an electronics package (see Fig. 3) which can be remotely located or mounted in a desk-type console as an available option.

Other optional devices include the following:

1) A hard-copy photographic printer, which provides a picture of the video in 10 seconds; a 4''x5'' positive print and a negative transparency are produced at the same time, for projection, duplication, or enlargement.
2) A stereo cassette tape recorder which can be used for any of the following purposes:
System operation

An example illustrates how the system operates. A subscriber at a telephone terminal who wants to send video to a remote office picks up his telephone (Fig. 4) and places a voice call in the usual manner (using telephone, switchboard, or leased channel).

He then rotates his camera to the subject material, and views it on his monitor. The camera is essentially prefocused, so only minimal adjustment of the focus control is required in most cases. If the subject material is a document or other motionless object, the subscriber presses his TRANSMIT button and the video is sent out over the line. If the subject might move after the camera is set, the subscriber presses his FRAME FREEZE button and views the stopped-action frame which appears on his monitor. He then presses TRANSMIT and the picture is sent out.

At the receive end, the operator switches to Videovoice mode and receives the picture. He can adjust his monitor for contrast and brightness as he would his home TV. Voice communication is possible at all times except while the video material is being transmitted, but after the full picture has been received, the system automatically switches back to “voice” mode. The video remains on display as long as required (ten minutes or more), and is erased only for presentation of the succeeding frame. Full voice conversation can continue during the no-video-transmission period even though the video is still being viewed, or the circuit can be disconnected without interfering with the displayed video at either terminal.

Technical considerations

In processing the standard video signal for transmission over the narrower voice-grade circuits, several major problem areas had to be considered: frame storage; scan conversion; synchronization and transmission; and control and sequencing.

Frame storage

Various alternative approaches have been publicized for slow-scan TV frame storage and display:

1) High-persistence display tube. In this type of implementation, a very-low-resolution picture is transmitted in about eight seconds per frame. The transmission is continuous, and the persistence of the tube is relied upon to create a still-image effect. However, this approach required a dark viewing area.

2) Memory tube. Here a storage-type cathode-ray tube records the video at the transmission speed but stores it for display over extended periods without requiring continuous transmission. The display area is limited in size and the video signal is not compatible with standard TV.

3) Magnetic disc storage. A magnetic disc memory device uses a small disc driven at speeds in the order of 1800 r/min. The video information is laid down on the rotating disc in a prescribed pattern and an associated TV synchronizing signal is recorded directly on the disc itself. The signal can be played back and displayed continuously at standard TV rates.

4) Silicon target storage tube. The silicon storage tube is a small, low-cost, single-ended, non-destructive read-out storage device. It is particularly useful where video information is to be displayed on conventional TV monitors, or “written” on a target from conventional TV cameras. The video can be recorded at TV speeds and read off at slower rates or, conversely, written at slow speeds and read off to a TV monitor. A video frame can be written onto the target in 1/30 of a second (one TV frame); it can then be continuously read out for display for periods of 15 minutes or more. At the end of this time, the stored message can be erased in a fraction of a second. A limiting resolution capability of 1000 TV lines is typical on present one-inch-diameter storage tubes (see Fig. 5).

The most desirable system lies between the disc approach and the use of a silicon-storage device. The high-persistence display tube and the memory tube suffer from serious disadvantages: unacceptably low resolution, limited gray scale capability, requirement for motionless objects, and incompatibility with standard TV components. Neither offers any real advantages to compensate for these shortcomings.

The disc, on the other hand, has advantages and does present some attractive features such as, a) almost unlimited...
picture viewing duration, b) large frame storage capacity at little added cost or increased size, and c) simultaneous read and write capability. However, it has major disadvantages due to the mechanical aspects of the device: high maintenance cost, low reliability, need for synchronization of power sources at the terminals, and other associated problems. For RCA Global Communications applications in particular, where field maintenance is an overriding consideration, these shortcomings ruled out this approach.

On the other hand, the list of advantages for the silicon-target storage tube is conclusive: all electronic, high resolution, high reliability, full compatibility, stop-action capability, and versatility in a variety of applications.

Scan conversion and synchronization

The scan converter, shown in Fig. 6, at the transmit end provides for converting the composite video signal of the TV camera (or a repetitive version of the stored single frame outputted from the frame-freeze unit) to a slow-scan signal suitable for transmission.

In the transmit mode, the scan converter reproduces the repetitive single video frame, scanning vertically from line to line to produce a filtered audio signal suitable for transmission by an analog modem over voice bandwidth facilities.

The picture can comprise either of two formats: normal-resolution (frame-freeze mode) with 30-second transmission time, or high-resolution (live camera mode) with 55-second transmission time. The received picture in both formats exhibits seven to eight shades of gray and is compatible with standard 525-line TV frames.

The procedure is to sample vertically across each of the horizontal lines in a specified sequence, proceeding across the target horizontally so that the received picture builds up from left to right. The receive display is totally blanked during transmission, except for a thin line cursor which moves horizontally, indicating the instantaneous position of the scanner.

Synchronization is more complicated in the slow-scan method than in standard TV because the sampling signal is pulsed rather than video; and the sampling period is of very short duration. Since available circuit bandwidth limitations preclude the transmission of synchronization signals other than long duration codes, synchronization is maintained by independent, highly stable crystal oscillators at each terminal.

In the receive mode, the incoming audio signal is demodulated, fed through a low-pass filter and then sampled in the same manner as at the transmitter and deposited in the storage tube. Synchronizing signals are provided by the local crystal oscillator system. At the end of the transmission, the picture is read from the storage tube for display with a standard TV raster.

Signaling and transmission

System operation in Videovoice is controlled from the transmit terminal. This means that the receiver is in constant "ready" mode, awaiting receipt of control signals from the transmitter. Basically, these signals include a multitone burst and a start tone, transmitted just prior to each picture transmission, initiated when the TRANSMIT button is depressed. The multitone signal switches the receiver into either normal-resolution or high-resolution mode; the start tone, which is transmitted immediately following the multitone, starts the erase and videowrite sequence in the receiver.

The slow-scan video information is transmitted over the line by a specially designed analog modem using subcarrier frequency modulation techniques. The characteristics of the transmitted signal are such as to make the essential video information relatively immune to circuit noise or to bandwidth restrictions.

Control and operation

Because the Videovoice system is intended for use by a relatively unskilled operator, the design philosophy is toward a minimum of manual control and maximum automation. In the camera, for example, only two controls are accessible for manual adjustment: a focus control and a lens tilt control. All other operating controls are contained in the desk-mounted telephone control unit.

The MICROPHONE Switch on the Telephone Control Unit (Fig. 4) is a rocker type and is used for microphone on-off. The STAND-BY, VIDEOVOICE, LIVE CAMERA, FRAME FREEZE, and TRANSMIT switches are pushbutton types. The pushbuttons are illuminated when energized. The STAND-BY and VIDEOVOICE pushbuttons are interlocked so that depressing STAND-BY releases VIDEOVOICE, but not vice versa. The remaining pushbuttons are not interlocked, so that several may be illuminated at the same time; for example, the TRANSMIT and VIDEOVOICE buttons will be lit when video transmission is in progress, along with either the LIVE CAMERA or FRAME FREEZE button, depending on the type of picture being transmitted. The RECEIVE pushbutton lights up only; it is not an operating control since the receive function is essentially unattended.

Many functions are automatic, including switching from local mode into receive mode upon receipt of a start-of-frame signal; blocking the voice in the direction of video transmission for the duration of such transmission; switching the monitor...
between the camera and the frame freeze unit; switching between the telephone handset (if used) and the loudspeaker system; switching from standby power on the video equipment to full power when the VIDEOVOICE switch is actuated; camera enlargement and reduction of the subject matter; and automatic light compensation in the camera to adjust for varying scene brightness.

During the receive cycle, all functions are automated. When a terminal initiates a video transmission, it transmits a code which contains synchronization and picture resolution information. When detected at the other terminal, the code starts the receive cycle, which consists of receiver-setup in proper resolution mode, storage-tube "erase", "write" and "read" cycling, and "start-of-frame" synchronization.

Commercial applications

VIDEOVOICE is a particularly significant development; it opens the way for visual communications by telephone, a capability which was heretofore precluded by both technology and economics. The long-range applications are almost unlimited. For the near term, we can mention the following:

1) Pictures of parts and assemblies for manufacturing purposes, maintenance and repair instruction, and spare parts catalog.
2) Advertising and sales promotion.
4) Drawings, charts, maps, and schematics for conferences and discussions.
5) Educational and training material to support an illustrated lecture, a blackboard-by-wire demonstration, or a textbook-on-audio tape.
6) Purchasing and inventory control documents and supporting photographs.
7) Telephone circuit interface to closed circuit TV systems.
8) Cardiograms, x-rays, and other medical graphic data.
9) Airlines arrival-departure information.
10) Brokerage display boards.
11) Banking, administration, and accounting material.
12) VIDEOVOICE can be used to input video data into a computer for data processing and to graphically display computer outputs.

Present customer usage

Since the original article was written, a substantial number of VIDEOVOICE terminals have been installed on customer premises and are being put to practical use, operated completely by their people. It is interesting to review the previous assessment of "commercial possibilities" in light of actual customer applications.

- One large company has installed VIDEOVOICE terminals in a number of locations distributed around the country; the terminals are used solely for remote conferences. Tying in the video capability with telephone conferencing directly into conference rooms has, in the opinion of this company, resulted in significant savings in travel expense as well as a substantial improvement in company-wide communications.
- The VIDEOVOICE conference facility is used by all departments including personnel, finance, technical, etc., besides executive conferencing. Presentation material for the conference is prepared for use via the VIDEOVOICE format; and the demand for the facility is so great that a strict reservation schedule is maintained.
- A large electronics company uses its VIDEOVOICE terminal system for remote field quality control; that is, they have the terminal installed in a vendor’s plant, and it is used to transmit close-up details of fabricated items, like a printed circuit board or mask, back to the central quality control facility for inspection and consequent acceptance or rejection. Where extra fine detail is required, the VIDEOVOICE camera is focused through a microscope with 40X magnification.
- One medically oriented installation provides for transmission between the doctor’s office and the hospital for close-up examination and remote evaluation of x-rays, medical graphic data, and microscope slides viewed by the VIDEOVOICE camera through the microscope eyepiece.
- Another medical use is from a conference room on the obstetric floor of a hospital to a distant university medical school lecture hall, serving a simultaneous education function.
- In the automotive industry, it is used for auto parts design and end-product evaluation between offices and the main plant.
- Major oil company usage is for ship-to-shore video from oil freighters at sea via satellite. The nature of the information includes seismic and oceanographic data, offshore oil installations, and equipment information, etc.
- A major cable TV company uses VIDEOVOICE terminals in their own offices for executive conferences in addition to offering it as an added service to subscribers on their private voice grade circuits. They feel this provides them with a competitive advantage over the regular telephone companies who do not offer such service.
- A major construction planning and civil engineering firm finds VIDEOVOICE very helpful to review layouts of plumbing, heating and air-conditioning ducts and for remote discussion of structural details as described in their architectural and mechanical drawings.
- A leading cosmetic firm makes use of VIDEOVOICE for transmission of advertising material and packaging design artwork for discussion and correction.
- In industrial and business use, technical conferencing and review of graphic information is common practice with many of our subscribers, including numerous RCA locations.
- Executive conferencing by VIDEOVOICE is also very popular and top executives in some of the major companies have expressed their satisfaction with the system.
- Government and defense agencies currently cover a broad range of specialized and sophisticated applications, although details cannot be generally released.

Among the more exciting applications, we can include the following:

- Video conferencing from shore points in the United States to the ship Hope in various ports-of-call in South America, mainly, for exchange of medical information.
- In conjunction with Apollo 17, VIDEOVOICE was used to relay continually updated tracking information, generated from earth locations, of the lunar rover on the moon surface, and also to monitor the deployment of recovery fleet ships in preparation for splashdown.
- VIDEOVOICE will be used for continuously updating antenna positions from video on the orbiting Skylab unit and collecting other information.
- International acceptance has also been very gratifying. The People’s Republic of China was quick to recognize the benefits for conferencing and for long-distance telecommunications with emphasis on transmission of Chinese characters. Spain was next to join the VIDEOVOICE international user list which now includes Sweden, as well. We further expect to receive orders shortly from at least five other European countries and others in South America and the Far East.

The interest in all of these countries is at the very highest levels of government and industry, which makes the outlook so much more promising.

In general then, all of the projected uses for VIDEOVOICE have materialized, and in addition, new ones are being discovered as the usage of the terminals expands. And, in keeping with our earlier expectations, it is being recognized that, where communications are required, the use of VIDEOVOICE can enhance efficiency and intelligibility, and at the same time, contribute significantly to circuit utilization.
Simplifying microcomputer architecture

J. Weisbecker

In 20 years computer hardware has become increasingly complex, languages more devious, and operating systems less efficient. Now, microcomputers afford the opportunity to return to simpler systems. Inexpensive, LSI microcomputers could open up vast new markets. Unfortunately, development of these markets may be delayed by undue emphasis on performance levels which prohibit minimum cost. Already promised are more complex next generation microcomputers before the initial ones have been widely applied. This paper discusses these points and describes a simplified microcomputer architecture that offers maximum flexibility at minimum cost.

LARGE SCALE INTEGRATION (LSI) of semiconductor devices has finally become a practical reality. The long-awaited revolution in electronic products appears to be at hand. The basis of this revolution is the ability to provide complex electronics at greatly reduced prices. Major cost reduction opens up entirely new markets and is as significant a development as the invention of the vacuum tube or transistor.

The four-function electronic calculator represents the first wave of the revolution. Further new markets will emerge with the ability to provide complete stored program computers at a fraction of current minicomputer costs. A number of microcomputer chip sets have already been announced. Consumer, educational, small business, and communications markets are prime targets for truly low-cost microcomputer based products. The architecture described in this paper was developed to satisfy the requirements of these potential new markets. Practical, stand-alone systems (including input/output device control and memory) requiring as few as six LSI chips are feasible with this architecture. Such systems have been breadboarded and programmed. Based on this experience, the microcomputer described appears to satisfy the requirements of a much wider range of applications than originally intended. It also appears to be simpler than most existing microcomputers. It is estimated that this new architecture compares favorably with the complexity of current four-function calculator chips. This simplicity is expected to provide significant production advantages. Improved yields and decreased testing costs are anticipated.

Since LSI improvements are permitting ever larger numbers of devices per chip, there are definite long-term advantages in minimizing microcomputer complexity. If the microcomputer is prevented from growing in complexity as the device per chip ratio improves, more of the system can be pulled back into a single chip. For example, small systems in which both memory and microcomputer are provided on a single chip become feasible resulting in added, long-term cost-reduction potential. This approach is ruled out when increased device per chip ratios are used to provide more complex microprocessors.

Joseph A. Weisbecker, LSI Systems Design, Solid State Technology Center, RCA Laboratories, Princeton, New Jersey, graduated from Drexel University in 1956 with the BSEE. Prior to joining RCA Laboratories in 1970, he was active in various product planning positions within RCA's Computer Systems Division. He has made major contributions to RCA computer development since 1956. During a three-year sojourn with a smaller company, he developed a number of customized data terminals. Mr. Weisbecker holds 19 patents with others pending. He is a member of IEEE and Eta Kappa Nu, and he has received an RCA Laboratories Achievement Award. His children have had a computer at home for several years which reinforced his interest in low-cost, widely available microcomputer systems.

Reprint RE-19-5-10
Final manuscript received November 26, 1973.
Design philosophy

Minimum system cost is the primary goal. To achieve this goal, an architecture is required that is both simple and flexible. Simplified computer architecture has received relatively little attention in the literature. Prior approaches toward simplified computers appear to be incompatible with microcomputer application goals.4,5,9

The architecture that was finally developed evolved from examining proposed applications. Another approach would have started with a more or less conventional minicomputer architecture and instruction set. This latter approach was discarded due to fundamental differences in minicomputer and microcomputer applications. It was also felt that a minicomputer starting point would not yield the simplest architecture.

Since a single-chip microcomputer promises minimum cost, the architecture was constrained to a 40-pin interface. Smaller microcomputer interfaces tend to require extensive multiplexing of interface signals which adds demultiplexing logic external to the microcomputer chip. This increases system cost.

An 8-bit parallel (or byte) architecture was chosen. This yields maximum performance consistent with interface pin constraints and is compatible with input/output requirements. One and four-bit organizations unduly restrict the range of potential applications. Sixteen or more bits exceed single-chip pin constraints or impose the need for multiplexed word transfers.

Since continued memory cost reduction is anticipated, techniques using memory to reduce hardwired logic complexity are heavily relied on. It is also apparent that many microcomputer applications will fall into the intelligent buffer category. For these reasons, direct memory addressing capability of up to 64K bytes is provided. Random-access memory (RAM) and read-only memory (ROM) can be mixed in any combination via a common memory interface. This is a distinct simplification over an architecture that provides separate RAM and ROM interfaces. The common RAM/ROM interface also enhances flexibility. System simplicity results since a single LSI chip containing both ROM and RAM segments will suffice for many applications.

While low memory costs can be expected, very low-cost systems will result only from minimizing memory capacity requirements. A unique architecture was devised which uses an 8-bit instruction format. This permits compact programs and subroutines. Useful systems requiring 1024 bytes or less of memory are possible with this format.

Random control logic uses chip area less efficiently than register/memory arrays. For this reason a simple, fixed cycle, microinstruction set was developed to reduce required control logic. The user has the option of programming directly in microcode, using a set of subroutines stored in memory (ROM/RAM), or a combination of these approaches. Sets of subroutines stored in memory are equivalent to applications-oriented macroinstructions and can readily be provided where ease of programming is important. On the other hand, many systems will utilize the microcomputer as a substitute for special purpose logic and can be programmed directly and efficiently in microcode. This approach retains most of the advantages of a micro-programmed computer but eliminates much of the specialized, hardwired sequencing and control logic usually associated with microprogrammed systems.6 Simple, short-subroutine-calling byte sequences provide flexible macroinstruction definition.

Whether used as a component of larger systems or as a freestanding computer, the microcomputer architecture requires efficient, flexible, input/output capability. This is provided via programmed byte transfers and a built-in direct memory access (DMA) channel. Programmed input/output byte transfer instructions provide maximum flexibility for input/output selection, control, and data transfer. The DMA channel facilitates high-speed I/O block transfer, TV display refresh, and initial program loading with a minimum of external logic. While the inclusion of a DMA channel adds negligible complexity to the microcomputer architecture, it greatly simplifies system design, leading to reduced overall cost. In addition to the two basic 1/O modes, four uncommitted flag lines are provided for activation by external devices. These flags can be tested as required by program. A flexible program interrupt capability also exists. Individual reset and load lines minimize external initializing logic.

The overall design philosophy consisted of developing a simple, flexible, microcomputer architecture which satisfies a wide range of potential applications at minimum cost. Performance levels were chosen to satisfy large-volume applications without overkill. The resulting architecture can be implemented initially on one or two chips using slow MOS technology.

Instruction execution times in the range of 4 to 8 μs are anticipated with LSI technologies that approach current TTL speeds. Experimental work has demonstrated that this performance level is adequate for most anticipated applications.

Microcomputer architecture

Fig. 1 illustrates the microcomputer architecture. "R" represents an array of sixteen, 16-bit general purpose registers. (This is essentially a 16x16-bit RAM.)

Registers P, X, and N are three 4-bit registers. The contents of P, X, or N select one of the 16 R registers. Register R(N) will be used to denote the specific R
Register D is an 8-bit register that functions as an accumulator. The arithmetic logic unit (ALU) is an 8-bit logic network for performing binary add, subtract, logical "and," "or," and "exclusive or" on two 8-bit operands. One operand is the bus byte and the other is contained in the D register. The D register can also be shifted right by one bit position. Add, subtract, and shift operations set a 1-bit overflow register (not shown) which can be tested by branch instructions.

The I is a 4-bit instruction register. Four-bit operation codes are placed in I and logical location addressed by the contents of two 8-bit operands. The I register selected by N, R(N) denotes the high-order byte. The contents of a selected R register (2 bytes) can be transferred to the A register. The 16 bits in A are used to address an external memory byte via an 8-bit multiplexed memory address bus. The 16-bit word in A can be incremented or decremented by "I" and written back into a selected R register.

M(R(N)) refers to a one-byte memory location addressed by the contents of R(N). This indirect addressing system is basic to the simplicity and flexibility of the architecture.

<table>
<thead>
<tr>
<th>Hex digit</th>
<th>Operation</th>
</tr>
</thead>
<tbody>
<tr>
<td>[0]</td>
<td>Increment R(N) by 1</td>
</tr>
<tr>
<td>[1]</td>
<td>Decrement R(N) by 1</td>
</tr>
<tr>
<td>[2]</td>
<td>Transfer R0(N) to D</td>
</tr>
<tr>
<td>[3]</td>
<td>Transfer R(N) to D</td>
</tr>
<tr>
<td>[5]</td>
<td>Transfer D to R0(N)</td>
</tr>
</tbody>
</table>

The operation of the microcomputer is best described in terms of its instruction set. A one-byte instruction format is used as shown in Fig. 2.

Each instruction requires two machine cycles. The first cycle causes an 8-bit instruction to be fetched from external memory and placed in the I and N registers. This is written as M(R(P)) — I,N. The 4-bit digit in register P selects R. R(P) is then transferred to A and used to address memory. While waiting for the memory access, A is incremented by "I" and replaces the original contents of R(P). The most significant digit of M(R(P)) is placed in register I and the least significant digit is placed in register N. At the end of an instruction fetch cycle, I and N always contain the 8-bit instruction originally addressed by the current program counter [R(P)], and this program counter has been incremented to point to the next memory byte in sequence. Note that any R register can be selected as the current program counter by merely changing the digit in register P. Multiple program counter systems and simple branch and link operations are readily achieved with this approach.

The next machine cycle always causes the instruction contained in registers I and N to be executed. This fixed two-cycle, fetch-execute sequence simplifies control logic and permits program interruption or DMA cycle stealing to occur only between instructions. This results in even further control simplification. Because the operation code in register I is limited to 4 bits, only 16 instruction types need be decoded. The 16 possible operations specified by the hex digit in I are listed in Table I.

The first group of instructions permits selecting any 16-bit general purpose

<table>
<thead>
<tr>
<th>Memory Operations</th>
</tr>
</thead>
<tbody>
<tr>
<td>[4]</td>
</tr>
<tr>
<td>[5]</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Miscellaneous Operations</th>
</tr>
</thead>
<tbody>
<tr>
<td>[6]</td>
</tr>
<tr>
<td>[7]</td>
</tr>
<tr>
<td>[8]</td>
</tr>
<tr>
<td>[9]</td>
</tr>
<tr>
<td>[10]</td>
</tr>
<tr>
<td>[11]</td>
</tr>
<tr>
<td>[12]</td>
</tr>
</tbody>
</table>

The last group of operations provides a variety of functions. The idle state can be entered via program or a reset line provided in the microprocessor interface. The idle state waits for externally generated program interrupts or DMA requests. The branch instruction performs a test specified by the value in N. As a result of this test, the next byte in memory, as addressed by R(P), is either skipped or placed in the lower half of R(P). This latter action causes a branch within the current 256-byte memory segment. Tests specified by N include zero in D, the states of four externally activated flags, and the status of the ALU overflow register. Two instructions, "D" and "E", permit the current digit in the P or X register to be modified. The "D" instruction provides the ability to change program counters at any point in a program. For example, "D4" would immediately change the current program counter to R(4). Branch and link operations are thereby facilitated. The "E" instruction permits changing the ALU operand or input/output byte address pointer. Instruction "F" permits 8-bit ALU operations. N designates the specific ALU operation to be performed. One of the operands comprises the byte contained in D. The other operand comes from memory. The second operand can be addressed by either R(P) or R(X) as specified by N. The result of ALU operations always replaces the original byte in D.
Instruction "6" permits byte transfers between memory and input/output devices via the common byte bus. The value of N specifies the direction of the byte transfer. M(R(X)) can be sent to an input/output device or any input/output device byte stored at M(R(X)). In the former case R(X) is incremented permitting X to be set equal to the current P value. The digit in N is made available externally during execution of the input/output byte transfer instruction. This digit code can be used by external I/O device logic to interpret the common byte bus. For example, specific N codes might specify that an output byte be interpreted as an I/O device selection code, a control code, or a data byte. Other N codes might cause status or data bytes to be supplied by an I/O device. In small systems the N code can directly select and control I/O devices.

Four flag lines that can be activated by I/O devices are provided. These can be used as general purpose I/O device status indicators (byte ready, error, etc.) These flag lines are tested by the branch instruction. Two interface lines control the built-in DMA channel. An I/O device can activate either an input or an output byte request line. At the end of the execution of the current instruction, a DMA channel cycle will occur causing the requested memory I/O device byte transfer to occur. R(0) is used for addressing memory during DMA cycles and is automatically incremented by one, following each byte transfer. Once initiated, blocks of data can be efficiently transferred between an I/O device and memory, independent from normal program execution.

A program interrupt line can be activated at any time by external means. At the end of the current instruction, an interrupt cycle will occur. During this cycle, X and P are placed in an 8-bit temporary storage register (T); P is then set to 1 and X is set to 2. Normal fetching and execution is then resumed. Activation of the interrupt line therefore causes a branch to the instruction stream addressed by R(1). R(2) should point to a memory area used by the interrupt routine to store the state of the machine for subsequent return from interrupt. Instruction "7" with N equal to 8 stores the contents of T in the memory location specified by R(X). It is a "save state" type instruction. If N is O, instruction "7" causes M(R(X)) to be placed in P and X. R(X) is incremented and an interrupt mask bit is reset. This instruction provides a "return after interrupt" function. The interrupt mask bit inhibits further responses to external activation of the interrupt line. This mask is always set by an interrupt permitting multiple interrupts to be queued under program control.

Programming considerations

Since the instruction set of this microcomputer differs considerably from that normally encountered, some comments relative to programming are in order.

A major difference between this architecture and more conventional organizations lies in the complete separation of operation codes and memory addresses. Conventional instructions have one or more addresses associated with each operation code. This system utilizes a limited table of memory addresses contained in a set of general purpose registers. These registers may also be used for program counters and data storage. Their use is entirely controlled by program, with the exception of R(0), R(1), and R(2). This structure is basic to the simplicity and flexibility of the architecture. It also permits the use of a short, 8-bit instruction format resulting in compact coding.

It has long been realized that storing a full memory address with each operation code is inefficient since a small number of unique memory addresses are repeated many times throughout a program. Various abbreviated addressing schemes have been used to permit more compact programs. These are almost always offered as optional alternatives to providing a full address in each instruction. Here we must always obtain a memory address from the limited, current set in the 16 general-purpose registers. We might intuitively suspect that this would be an unduly restrictive approach. Surprisingly, it turns out to be relatively easy to write programs and is highly efficient relative to the amount of memory used. A variety of programmers, from those who have only used high-level languages to engineers with limited programming experience, have had little difficulty in adapting to this architecture.

A number of programs have been written for a breadboard version of the microcomputer with a variety of input/output device attachments. This experience has validated the flexibility and efficiency of the architecture. For example, a four-function calculator program was found to require only 1024 bytes of memory, most of which could be ROM. This included provision for keyboard input; operands up to 30 digits; TV display refresh storage; 5x7 digit font conversion table; push-down stack; and algorithms for signed, n-digit decimal add, subtract, multiply, and divide. An interpreter for a simple, decimal, tutorial language was written in less than 600 bytes. A number of game and/or educational programs require well under 1000 bytes of memory. Many small business and communications systems programs are possible with 2000 to 4000 bytes of memory.

While the instruction set initially appears quite limited, it should be kept in mind that each operation requires only one byte of storage (ROM or RAM). Short sequences of these microinstructions readily provide macro-operations.

Apparent weaknesses in the architecture are the limited branch capability (within 256 bytes) and the lack of a hardwired program stack for multilevel nested subroutines. These apparent oversights are the result of a deliberate design philosophy which eliminates special purpose logic for those functions which are performed easily by subroutines. The architecture permits a flexible subroutine "call" and "return" system requiring less than 70 bytes of memory. This includes a push down stack for nested subroutines. By providing this system in software (or firmware) it can be tailored to individual applications.

Where extensive programming effort is required, a set of applications-oriented subroutines is easily developed. These subroutines constitute a user-oriented macroinstruction set for minimum effort programming. This technique has proved extremely useful in an experimental small
business system.

For microcode programming, an assembly language has been developed. This approach simplifies machine language coding considerably. An interactive simulator greatly facilitates initial program debugging. Both of these microcomputer software support systems are readily modified to run on existing time-sharing systems.

In general, the simplified microcomputer presents no difficulty in programming. It provides a simple set of short, easy to understand microinstructions that do not require high skill levels to use. For specific applications, tailored macroinstructions are readily provided via a flexible subroutine handling system.

Typical systems

Several systems using the microcomputer can be outlined. Many others are, of course, possible.

Fig. 3 indicates a possible microcomputer-based calculator. ROM and RAM might be provided on one chip resulting in a basic three-chip calculator. Functions could easily be added with ROM increments. This type of system could also provide a programmable calculator.

Fig. 4 illustrates a stand-alone system which might require only six LSI chips total.

It is assumed that 4x1024-bit memory chips will be available within the next several years. Subsequent LSI improvements could further reduce the chip count. Use of a small keyboard, audio cassette, and CRT display might reduce system cost to a few hundred dollars. Such a system could have wide application in consumer and educational markets. This system, with more memory, hardcopy output, and low-cost floppy disk (or magnetic bubble bulk storage), would provide the basis for a wide range of inexpensive, turnkey, small business systems.

Fig. 5 illustrates a large computer system in which each I/O device is controlled by a dedicated microcomputer providing an intelligent buffer, as well as a replacement for special purpose logic. RAM, ROM, microcomputer, I/O device and central computer interface circuits could readily be provided on a small set of LSI chips. The microcomputer DMA feature is extremely useful for high-speed block transfers in this type of system. Downline loading of the microcomputer memory can immediately change its mode of operation. Off-line editing and maintenance is provided free. This type of large-scale system approach will become more popular in the future as microcomputer costs decrease.

The performance level of the simplified microcomputer described is more than adequate for the above types of systems as well as many others.

Conclusions

Much current microcomputer development effort appears to be directed toward improved performance. There is, however, a need for simple, minimum cost structures that will satisfy large-volume applications which do not require minicomputer performance levels. These microcomputers must also be organized to reduce total system cost. One such microcomputer architecture has been developed. It promises low cost, together with minimum external memory and system logic requirements. Hopefully, microcomputers of this class will accelerate the development of major new markets.

Currently high input/output device costs might be used as an argument against minimizing microcomputer cost. This is extremely short-sighted. The availability of ten-dollar microcomputer chips will, by itself, exert considerable pressure on the development of compatible low-cost I/O and bulk storage devices. Even now there are many potential new products that demand minimum cost microcomputers of the type described.

Because of its flexibility and potential for low-cost systems, RCA is currently developing a COS/MOS-LSI version of this microcomputer — COSMACH, COS, versions are also being investigated for applications requiring higher instruction execution rates. Both implementations are expected to find wide application in a variety of future products.

Acknowledgments

The following people have devoted considerable effort toward evaluating and developing software for the microcomputer described here: S. Heiss, A.R. Marcantonio, J.T. O'Neil, A.D. Robbi, P.M. Russo, R.O. Winder, and C.T. Wu. Without this effort, validation of the architecture would have been impossible.

References

Recent developments in photomultipliers for liquid scintillation counting

D. E. Persyk | T. T. Lewis

This paper discusses a new photomultiplier for liquid scintillation counting, the 450IV4. It is contrasted with the generic type 450IV3 in terms of E/B in the tritium and carbon windows, cross-talk, background and accidental count rates. A discussion of the effects of tube materials upon photomultiplier performance is given. Future trends in photomultipliers are also discussed.

Tom T. Lewis, Engineering Leader, Phototube Applications Engineering, Electronic Components, Lancaster, Pa., received the BA in Physics from Wilkes College in 1964 and the MS in Engineering in 1971 from Pennsylvania State University. Employed by R.C.A. in 1964, Mr. Lewis was first assigned to the Photomultiplier Product Development group, engaged in the design and development of photomultipliers for various space applications. From 1967 to 1968, he served as a manufacturing engineer working on production problems associated with photomultipliers for liquid scintillation counting. In 1969, Mr. Lewis was transferred to Applications Engineering, providing technical assistance to photomultiplier customers. In 1972 he was promoted to his present position. Mr. Lewis is the recipient of an RCA Team Achievement Award for his contributions in the successful transition from lab to factory of high gain gallium phosphide dynodes for photomultipliers. Mr. Lewis is a member of the American Optical Society.

Dennis E. Persyk, Engineering Leader, Phototube Product Development, received his B.A. in Physics in 1963, and the M.A. in Physics in 1964 from the University of Wisconsin. He joined RCA in June 1966 as an engineer in the Photomultiplier Product Development Activity at Lancaster. Since that time, he has been involved with computer-aided photomultiplier tube design and development of photomultipliers for use in nuclear medicine. He has written several papers in the field of photomultiplier design and applications. He was promoted to his present position in 1972. Mr. Persyk is a member of the American Vacuum Society, and the Nuclear Instruments and Detectors Committee of the IEEE Group on Nuclear Science.

New photomultiplier

The new 450I photomultiplier tubes employ the Matheson-type photocathode-to-first dynode (“front end”) geometry coupled to a 12-stage, BeO dynode, focussed electron-multiplier structure. The anode-output structure consists of a grid-type anode and 3-element, parallel-plane transmission line; the main features are shown diagrammatically in Fig. 1.

The two-inch diameter spherical-section front end affords two advantages over alternative planar faceplate designs:

1) The photocathode trajectories are nearly isochronous, and the electric field strength near the photocathode is higher, affording a faster time response, and thus narrower time coincidence window for improved noise rejection, and

2) The spherical section permits a reduced mass of faceplate material which in turn provides a lower background from the photomultiplier.

A bialkali (K-Cs-Sb) photocathode is used which has a quantum efficiency of approximately 31% at 385 nm. The twelve BeO dynodes provide typical gains of 1.4 \times 10^7 at 2 kV applied voltage. Anode pulse risetime for delta-function excitation of the photocathode is approximately 2.5 ns, affording time response that exceeds the requirements of most present-day liquid scintillation counting systems.

The focused electron-multiplier structure has the following advantages over venetian-blind type multiplier structures which are also used in photomultiplier detectors:

CONSIDERABLE new interest is being expressed in the field of liquid scintillation counting. This interest exists due to increased efforts in the clinical, research, and environmental monitoring areas. The latter is of principal concern because of the increase in tritium-producing nuclear power facilities. Attendant with this interest there has been an introduction of new photomultipliers for use in liquid scintillation counters.

Reprint RE·19-5-2
This paper was presented at the Symposium on Liquid Scintillation Counting held in Brighton, England, September 3-6, 1973 and is included with the permission of Heyden & Son, Ltd from "Liquid Scintillation Counting," Volume 3, edited by M. Crook and P. Johnson, in press.
Materials considerations

The choice of materials for the envelope and internal tube parts of photomultipliers used in liquid scintillation counting applications is very critical. One obvious requirement is that all tube materials shall exhibit negligible radioactivity. Of prime concern in this regard is the glass used in the vacuum envelope housing the photomultiplier.

The glasses most commonly used in photomultiplier construction contain some $^{40}$K and thorium daughter-products. These radioactive contaminants are found in the raw materials comprising the glass, such as sand, and are also introduced into the glass when it is melted in thoria refractories.$^3$ The former problem can be overcome by judicious monitoring of the basic constituents prior to melting, and the latter problem can be solved by proper choice of melting vessel, platinum being an ideal choice.

Less well understood, but of equal importance, is the fluorescent activity of the faceplate glass used in the photomultiplier. All glasses scintillate to some extent when excited by X-rays or gamma-rays, and such scintillations lead to undesired background counts in the case of liquid scintillation counters. A pronounced scintillation-like effect due to cosmic ray bombardment is commonly seen in photomultipliers.$^{14}$ Unpublished measurements at RCA Lancaster indicate that commonly used photomultiplier window materials vary by several orders of magnitude in their scintillation efficiencies. The problem of scintillating windows can be approached by at least two avenues:

1) Employ special window materials with very few luminescent centers (<1%$^{14}$), or
2) Add dopant to common window materials to quench the luminescence (or shift it to sufficiently long wavelengths where the cathode is not sensitive.

Both approaches have been tried in photomultipliers with good success.

It is clear that considerable improvements can be made in photomultipliers by incorporating the best possible choice of materials in their construction. This approach has been followed in the 4501V3 and 4501V4 photomultipliers, as will be described in the next section.

Tube experiments

The experiments were performed in a commercially-available liquid scintillation spectrometer.$^7$ Data are reported on a large (>200) tube sample, ensuring statistically valid average values. In all experiments the tubes were paired in the sense that tubes with comparable gains were used together. This requirement stemmed from the limited gain-adjust capability of the instrument (a characteristic of nearly all commercially available spectrometers) and was therefore an instrument-imposed constraint.

Unless otherwise stated, all measurements were made in the tritium energy window and with a tritium counting efficiency of 60%. The carbon counting efficiencies for carbon-above-tritium measurements averaged approximately 70%. The figure-of-merit ($E^2/B$) or, efficiency in percent, squared, and divided by background in counts/minute $(c/m)$ is stated in keeping with common usage.

The crosstalk measurements were made with the spectrometer elevator down and with an empty counting chamber. The background measurements were made with a Packard background standard. The accidental count rate was measured with the elevator shaft blocking the tubes.

To prove the magnitude of the cosmic ray/Cerenkov effect, certain measurements were repeated with the entire spectrometer rotated 90° so that the photomultiplier axis was vertical instead of horizontal.

Results and interpretation of experiments

The results of the experiments are tabulated in Table II. In brief, the $E^2/B$ has been increased from 160 to 210 in the 4501V4 as contrasted to the 4501V3. It should be noted that this increase was
achieved through a reduction in background (22.5 c/m to 17.1 c/m) and crosstalk (12 c/m to 7 c/m). From these data one may conclude that the material-related improvement in the 4501V4 was due largely to a reduction in scintillation efficiency of the window, as opposed to a reduction in radioactivity of the envelope, as the background standard is an efficient scintillator.

To assess the magnitude of cosmic ray interactions, the crosstalk and accidental rates were measured with the spectrometer in the horizontal (normal operating) position and in the vertical position. This experiment was performed on only one pair of tubes owing to the mechanical difficulties of rotating the spectrometer. The data are summarized in Table III.

The crosstalk in the tritium window increased from 5.7 to 8.6 c/m in going from horizontal to vertical, a ratio of 1.5. The accidental rate increased from 0.34 to 0.95 c/m, horizontal to vertical, a ratio of 2.8. The higher ratio for the accidental rate is ascribed to the increased probability of an energetic particle interacting with both tubes with the system vertical.

The accidental rate may be calculated from the relationship \( A = 2 N_1 N_2 \lambda \) where \( N_1 \) and \( N_2 \) are the dark count rates of the two individual photomultipliers and \( \lambda \) is the resolving time of the coincidence system. Typical values in our experiments were \( N_1 N_2 \approx 10 \times 10^3 \) counts per minute and \( \lambda = 20 \) ns, giving \( A = 0.07 \) c/m. The order-of-magnitude increase in measured accidental rate over calculated rate was attributed to cosmic ray events, as demonstrated by the greatly increased accidental rate for vertical operation of the system.

Future improvements

Further improvements can be made to photomultipliers to reduce their background and fluorescent properties, and research is being carried out in this area. Additional refinements can also be made to systems to improve performance with existing photomultipliers. A future possibility would be a single tube liquid scintillation counter utilizing pulse shape and pulse height discrimination to eliminate the need for a two-tube coincidence system. Such a system with a suitable photomultiplier could distinguish between tube dark pulses (single-electron in origin), tube Cerenkov pulses (fast pulses of multiple-electron origin) and scintillator signal pulses (relatively slower pulses of multiple-electron origin) on the basis of pulse height and pulse shape (time) information. Photomultipliers such as the developmental type C31024 may have the basic time response and electron resolution capabilities to render such a system feasible in today's state-of-the-art.

Summary

Materials used in photomultiplier construction have been investigated for radioactive and fluorescent properties. The results of this investigation have led to the introduction of the 4501V4 photomultiplier which offers a tritium \( E^2 / B \) of 210 as opposed to 160 for the 4501V3 present type. Further improvements in these photomultipliers are being researched, and possibilities of a single tube liquid scintillation spectrometer utilizing pulse shape and pulse height discrimination are being studied.

Acknowledgments

The authors are grateful to R. M. Shaffer and J. D. Doering for providing data, and to R. C. Pontz for encouragement in writing this paper.

References

3) Private communication with D. C. Smith, Vineland, N. J.
4) Data provided by A. W. Conantman and D. D. Crawshaw, RCA Corp.
5) Young, Andrew T., "Cosmic Ray Induced Dark Current in Photomultipliers", RSF 37, p. 1472 (1966)
220 Energy and Power Sources

batteries, solar cells, generators, power supplies.

SOLAR CELLS, Temperature, illumination and fluorescence dependence of current and voltage in electron-irradiated...T.J. Fein. A.F. Oberchain (AED,Pr) IEEE Photovoltaic Specialists Conf.; Palo Alto, Calif; 11/15/73

240 Lasers, Electro-Optical and Optical Devices

design and characteristics of lasers, components used with lasers, electro-optical systems, lenses, etc. (excludes: masers).

BUBBLE DOMAINS, Lensless optical detection of...R.H. Shashbander, J. Drugan (Labs,Pr) RCA Review, Vol. 34, No. 3, 336-343; 10/73

CO2 laser technology at RCA Limited...H.A. Crane (Labs,Pr) IEEE Circuits and Systems Society’s Computer-Aided Network Design Comm. Workshop, Carefree, Ariz; 9/24-25/73

ELECTRO-OPTIC WAVEGUIDES...R.J. Hammer (Labs,Pr) ONR Program Review, Washington, D.C.; 10/1-4/73

INJECTION LASERS, Experimental properties of...S.H. Sommers (Labs,Pr) J. Appl. Phys., Vol. 44, No. 8, pp. 3661-3668; 8/73

LASER DIODES, Heterojunction...H. Kressel (Labs,Pr) NERIEM Seminar, Boston, Mass.; 11/73

LASER DIODES, Red-light-emitting...Al. GaAs, heterojunction...H. Kressel, F.Z. Haskel (Labs,Pr) J. Appl. Phys., Vol. 44, No. 9, pp. 4227-4230; 9/73


LASER SCANNERS, Frequency response of...R.W. Cohn, I. Goron (Labs,Pr) J. Optical Soc., Vol. 63, No. 9, pp. 1071-1079; 9/73

METAL-VAPOR LASERS, Noblest of...K. Hamang (Labs,Pr) Laser Focus, pp. 39-40; 9/73

MICRO FRENSNEL ZONE PLATES, Properties and fabrication of...A.H. Firester (Labs,Pr) Appl. Opt., Vol. 12, No. 7, pp. 1869-1702; 8/73

MODULATION OF lasers, A comparative evaluation between internal and external...A.L. Walsenberg (Ldn,Can) IEEE J. of Quantum Electronics, Vol. 9, No. 11; 11/73

MODULATOR, Fast electro-optic waveguide diode...J.M. Hamer, D.J. Chinn (Mtn, Mt) Dutfy (Labs,Pr) Appl. Phys. Lett., Vol. 23, No. 4, pp. 176-177; 5/15/73

OPTICAL FILTERS, Selection of objective functions for optimization of...R.L. Craine (Labs,Pr) IEEE J. of Quantum Electronics, Vol. 9, No. 3, pp. 401-407; 8/73

VOLUME HOLOGRAPHY, science and applications...D.L. Staebler (Labs,Pr) Second Int'l Conf. on Electroholography, Washington, D.C.; 10/24-27/73

VOLUME PHASE HOLOGRAM recording medium...A. new...A. Bloom, R.A. Bantolin (Labs,Pr) Optical Soc. of America Annual Mtg.; Rochester, NY; 10/73

245 Displays

equipment for the display of graphic, alphanumeric, and other data in communications, computer, military, and other systems. CRT devices, solid state displays, holographic displays, etc.

DISPLAY, Blue-green numeric using electro-luminescent GaN...J.J. Pankow (Labs,Pr) RCA Review, Vol. 34, No. 2, pp. 336-343; 8/73

DISPLAYS, Liquid crystal...L.A. Goodman (Labs,Pr) J. Vac. Sci. & Tech., Vol. 10, No. 5, pp. 804-823; 9/10/73

DISPLAYS, Light-emitting diodes and semiconductor materials for...C.J. Nuese, H. Kressel, J. Ladany (Labs,Pr) J. Vac. Sci. & Tech., Vol. 10, No. 5, pp. 772-788; 9-10/73

FOCUSED-IMAGE HOLOGRAMS for a microfluidic system, Formation and replication of...D. Meierhofer (Labs,Pr) Soc. of photographic Scientists & Engineers Mtg., Washington, D.C.; 10/24-27/73

VIEWING SCREENS, Holographic and in-terferometric...P.F. Murachofer (Labs,Pr) Appl. Opt., Vol. 12, No. 9, pp. 2180-2184; 10/73

250 Recording Components and Equipment

disk, drum, tape, film. holographic and other assemblies for audio, image, and data systems.

DIGITAL RECORDERS, Ultra high data rate...J.S. Griffin (ASD,Can) Proc. of SPIE, Vol. 36, page 85; 73


VIDEO TAPE, A basic highly maintainable airborne...P.F. Muraco (CSD,Can) Proc SPIE, Vol. 36, page 66

280 Thermal Components and Equipment

heating and cooling components and equipment, thermal measurement devices, heat sinks.

COLD PLATE design analysis for electronic equipment...E.D. Veilleux (ASD,Can) J. of Spacecraft and Rockets; 10/73

MW CONVERTER, Performance models for...R.R. Lorenzen (EC,Har) IEEE JECCEG Mtg., Philadelphia, Pa.; 8/12-17/73

SiGe CONVERTER, Thermal performance of foil insulated...R.R. Lorenzen (EC,Har) IEEE JECCEG Mtg., Philadelphia, Pa.; 8/12-17/73

SiGe thermoelectric devices, Dynamic performance of foil insulated...R.D. McLauingh (EC,Har) IEEE JECCEG Mtg., Philadelphia, Pa.; 8/12-17/73

SiGe Thermoelectric properties, Updated...A. Amith (EC,Har) IEEE JECCEG Mtg., Philadelphia, PA; 8/12-17/73

310 Spacecraft and Ground Support

spacecraft and satellite design, launch vehicles, payloads, space missions, space navigation.

NUTRITION DAMPING utilizing spacecraft mass properties, Active...K. Phillips, (AED,Pr) IEEE Trans. on Aerospace and Electronic Systems 9/73

PROPOSITION MODULE, Ascent phase and orbital correction...D. Balzer, A.C. Beilsey (AED,Pr) AIAA 9th propulsion conf., Las Vegas, Nev.; 11/71/73

312 Land and Water Transportation

navigation systems, automotive electronics, etc.

VEHICLE EMISSIONS...L.R. Hulls (ASK,Can) IEEE Merrimack Valley Mtg.; 10/73

VEHICLE IDENTIFICATION system, A microwave automatic...H. Stara, P.I. Kiensch, J. Rosen (Labs,Pr) IEEE Electrical/Electronic Insulation Conf., Chicago, IL; 10/25/73

315 Military Weapons and Logistics

missiles, command and control.

FUTURE MILITARY SYSTEMS, The impact of surface wave technology on...E.J. Schmitt (ATL,Can) RCA Review, Vol. 34, No. 3; pp. 442-458; 9/73

320 Radar, Sonar, and Tracking

radar, optical, and other systems for detection, acquisition, tracking, and position indication.

NONCOHERENT RADAR integration gain, cross-correlation form expressions for...H. Urickowitz (Lnl,Can) IEEE Trans. on Aerospace and Electronic Systems, Vol. AES-9, pp. 781-783; 10/73

OPTICAL DOPPLER RADAR SYSTEM, Corrections to FM laser noise effects on...A. Waksberg (Ldn,Can) IEEE J. Aerospace and Electronic Systems, Vol. 9, No. 4; 7/73

SURVEILLANCE of satellites...some techniques and systems aspects...J.D. Proehl (ASD,Can) MIT Aeronautical Dept. Seminar; 10/15/73

WEATHER RADAR, Modern...G.C. Jung (EASD,Van Nuys) Brazilian Air Show, Sao Paulo, Brazil; 9/17/73

325 Checkout, Maintenance, and User Support

automatic test equipment, (ATE), maintenance and repair methods.


340 Communications

industrial, military, commercial systems, telephony, telegraphy and telemetry, (includes: television, and broadcast radio).


Record, Vol. Pate 10A-1, Section 10

MICROWAVE NETWORK optimization program...C.M. Kudliski, V.K. Jha, J.W. Bandier, J.R. Popovich (Labs,Pr) Microwave Symposium, Boulder, Colorado; 6/73

MICROWAVE REPEATER for a 24-channel domestic satellite system, Design of a...M.V. O’Donovan, C.M. Kudliski, L.A. Keys (Labs,Can) RCA Review, Vol. 34, No. 3, 9/73

MOBILE RADIO performance in urban hilly terrain, 900 MHz and 450 MHz...F.A. Barton, G.A. Wagner (CS,Can,lands) 1973 10/15/73

SATELLITE COMMUNICATION system for remote areas, A multi-purpose uhf...O.S. Roscoe, P.J. Osborne (Ldn,Can) Int’l Elec- trical Electronics Conf. and Exposition; Toronto; 1973

SHF high power airborne communications antenna...J.P. Grabowskki, E.L. Laphean (MSRd,Max) 95th Mtg. of the Avionics Panel on “Antennas for Avionics,” Munich, Germany; 11/26/73

SWITCHING approach for local and remote subscribers, Modular...P.J. Bird (CSD,Can) National Telecommunications Conf., Atlantic City, NJ; 11/28-30/73; Conf. Proc.

TELEX, Computer controlled...E.G. Tindall
345 Television and Broadcast

Television and radio broadcasting, receivers, transmitters, and systems; television cameras, recorders, studio equipment.


RVS CAMERA performance characteristics, ERTS two-inch — B.P.Miller, G.A. Beck, J.M. Barletta (AED.Pr) AIAA J. Spacecraft and Rockets; Oct/73


360 Computer Equipment

processors, memories, and peripherals.


COMPUTERS, Application of — M.S. Corrington, (ATL.Cam) Over-65 club; Haddonfield, NJ: 11/29/73

370 Computer Programs (Scientific)

specific programs and techniques for scientific use, computation, simulation, computer-aided design.

HIGHER-ORDER LANGUAGE considerations for radar systems application — Dr. J.A. Stieie/Dr. I.J. Galbrai (MSRD,Mrstn) Asilmar Conf. on Circuits, Systems, and Computers, Monterey, Calif. 11/27-29/73

380 Graphic Arts and Documentation

printing, photography, and typesetting; writing, editing, and publishing; information storage, retrieval, and library science.

WORD PROCESSING, a new approach to internal profit — B. Piscopo (ASD.Burl) Technical Communications (JSTC) Vol. 20, No. 4, (4th quarter 1973)

88
As an industry leader, RCA must be well represented in major professional conferences...to display its skills and abilities to both commercial and government interests.

How can you and your manager, leader, or chief-engineer do this for RCA?

Plan ahead! Watch these columns every issue for advance notices of upcoming meetings and "calls for papers". Formulate plans at staff meetings—and select pertinent topics to represent you and your group professionally. Every engineer and scientist is urged to scan these columns; call attention of important meetings to your Technical Publications Administrator (TPA) or your manager. Always work closely with your TPA who can help with scheduling and supplement contacts between engineers and professional societies. Inform your TPA whenever you present or publish a paper. These professional accomplishments will be cited in the "Pen and Podium" section of the RCA Engineer, as reported by your TPA.

Ed. note: Meetings are listed chronologically. Listed after the meeting title (in bold type) are the sponsor(s), the location, and the person to contact for more information.

**Dates of upcoming meetings—plan ahead**

Ed. note: Calls are listed chronologically by meeting date. Listed after the meeting title (in bold type) are the sponsor(s), the location, and the deadline information for submittals.

**Dates and Deadlines**


April 2-4, 1974 — JI. IFIP Fall COMPCON Appliance Int'l Sources


April 29-MAY 1, 1974 — SOUTHEASTCON, Region 3, ISHM, Dutch Inn, Orlando, FL Prog Info: B. E. Mathews, Florida Tech., Box 20000, Orlando, FL 32816.

April 28-22, 1974 — 46th Annual Gas and Diesel Engine Power Conference and Exhibit, ASME, Astroworld Hotel, Houston, TX Prog Info: Don Belenger, Staff Assistant, Information Services, ASME, United Engineering Center, 345 E. 47th Street, New York, NY 10017.

April 29-MAY 2, 1974 — Annual Meeting, Electronics Division, American Ceramic Society, Conrad Hilton Hotel, Chicago Prog Info: F. P. Reidel, Executive Director, American Ceramic Society, Inc., 65 Ceramic Drive, Columbus, OH 43214.


May 7-8, 1974 — Appliance Tech. Conference, S-AIA, Columbus, OH Prog Info: Robert E. Kind, Ranco Box 8187, Columbus, OH 43201.


May 15-17, 1974 — Earth Environment & the 9th Annual Meeting, American Ceramic Society, Conrad Hilton Hotel, Chicago Prog Info: F. P. Reidel, Executive Director, American Ceramic Society, Inc., 65 Ceramic Drive, Columbus, OH 43214.


Aug. 5-10, 1974 — IFIP Congress 74 and Exhibition, International Federation for Information Processing (IFIP), Stockholm, Sweden Prog Info: Paul E. Welch, Public Information Officer for the U.S. Committee for IFIP Congress 74 at Time Inc., Time & Life Bldg., Rockefeller Center, New York, NY 10020.

Oct. 22-13, 1974 — Linear Electric Machines, Power Division of the Institution of Electrical Engineers in association with the IEEE (UK and Rep. of Ireland Sect.) Prog Info: IEEE Conference Department, Savory Place, Lon­don WC2R 0BL.


Calls for papers—be sure deadlines are met.


SEPT. 10-13, 1974 — Western Electronic Show & Convention WESCON (region 6) Los Angeles, Calif. Deadline Info: (abstract) 4/7/4 to WESCON Office, 3600 Wilshire Blvd., Los Angeles, Calif. 90010.

SEPT. 18-20, 1974 — Fall Meeting Electronics Division - American Ceramic Society, Denver Hilton Hotel, Denver Colorado Deadline Info: (Titles and Authors) 5/17/4 to Program Chair.
Thermally-controlled crystalline lasers — P. F. Joy, Jr. and D. G. Herzog (ATL, Cam.) U.S. Pat. 3784529, January 8, 1974

Electronic Components

In-line electron gun — R. H. Hughes (EC, Lanc.) U.S. Pat. 3772545, November 13, 1973

Non-thrombin electron emissive tube comprising a ceramic heater substrate — A. F. McDonie, R. D. Faulkner, J. L. Rhoads (EC, Lanc.) U.S. Pat. 3777209, December 4, 1973


Color television picture tube screening method — F. Van Heikken (EC, Lanc.) U.S. Pat. 3782653, January 1, 1974

Magnetically-focused cathode-ray tube comprising a tilted and skewed off-axis electron gun — E. Luedicke (EC, Prj.) U.S. Pat. 3783326, January 1, 1974

Filamentary cathode mount and mounting method — B. B. Adams, Jr. (EC, Lanc.) U.S. Pat. 3783327, January 1, 1974

Solid State Division


Method for depositing refractory metals — W. A. Grill (SDS, Som.) U.S. Pat. 3785802, January 15, 1974

Semiconductor amplifier protection — C. F. Wheatley, Jr. (SDS, Som.) U.S. Pat. 3786364, January 15, 1974

Consumer Electronics


Modulator system — B. A. Hjortzberg (CE, Indpls.) U.S. Pat. 3775554, November 13, 1973

Modulator system — D. J. Carlson (CE, Indpls.) U.S. Pat. 3775557, November 13, 1973

Television deflection circuit with low power requirement — T. J. Christensen (CE, Indpls.) U.S. Pat. 3784857, January 8, 1974

Leakage detector for determining possible shock hazards to humans — L. P. Thomas (CE, Indpls.) U.S. Pat. 3784903, January 8, 1974

Instant-on-circuit for a television receiver offering independent filament voltage control — R. J. Gries (CE, Indpls.) U.S. Pat. 3783335

Dynamic convergence circuits — M. W. Hill (CE, Indpls) U.S. Pat. 3786800, January 15, 1974

Palm Beach Division

Storage circuits — C. J. Fassbender (PBD, Palm Beach Gardens) U.S. Pat. 3784916, January 8, 1974

Laboratories

Light modulator — R. A. Gange (Labs, Prj) U.S. Pat. 3772912, November 13, 1973

Apparatus for and method of reproducing an electrostatic charge pattern — R. Williams (Labs, Prj) U.S. Pat. 3778634, December 4, 1973

Stabilized photoresist composition — E. B. Davidson (Labs, Prj) U.S. Pat. 3776366, December 4, 1973

Method of making a magnetic recording head — R. A. Shabtender (Labs, Prj) U.S. Pat. 3777369, December 11, 1973

Damping means for ultrasonic transmitters — A. G. Lazzery (Labs, Prj) U.S. Pat. 3777770, December 11, 1973

Gated holographic coding system for reducing alignment requirements — W. J. Hanner (Labs, Prj) U.S. Pat. 3778128, December 11, 1973

Storage circuit using multiple condition storage elements — E. C. Ross (Labs, Prj) U.S. Pat. 3781570, December 25, 1973

Holographic memory with high intensity compensation means — D. H. Villkomerson (Labs, Prj) U.S. Pat. 3781630, December 25, 1973

Harmonic radar detecting and ranging system for automotive vehicles — H. Staras, J. Shefer (Labs, Prj) U.S. Pat. 3781879, December 25, 1973

Protection circuit therein — D. P. Dorsey (Labs, Prj) U.S. Pat. 3784870, January 8, 1974

Broadband apparatus using high efficiency avalanche diodes operative in the anomalous mode — K.K. Chang, H. J. Prager, S. Westbrod (Labs, Prj) U.S. Pat. 3784925, 1973

Photosensitive charge storage electrode having a selectively conducting protective layer of matching valve band on its surface — C. R. Wronski, W. M. Yim, J. Dresner (Labs, Prj) U.S. Pat. 3783324, January 1, 1974

Encapsulated semiconductor device assembly — G. A. Swartz, R. E. Chamberlain (Labs, Prj) U.S. Pat. 3782348, January 1, 1974

Electroluminescent semiconductor device capable of emitting light of three different wavelengths — J. J. Piatak (Labs, Prj) U.S. Pat. 3783353, January 1, 1974

Mullitlayer heat sink — A. F. Arnold (STTC, Som.) U.S. Pat. 3780796, December 25, 1973

High-density capacitive information records and playback apparatus therefor — T. O. Stanley (Labs, Prj) U.S. Pat. 3783196, January 1, 1974

Constant pulse width generator — H. M. Fox (GCS, Cam.) U.S. Pat. 3783304, January 1, 1974

Apparatus for automatic color balancing of television camera signals — L. J. Rabin (GCS, Cam.) U.S. Pat. 3786177, January 15, 1974

Microwave balanced mixer — R. L. Ernst, S. Yahn (GCS, Som.) U.S. Pat. 3772899, November 13, 1973

Patents

Keyboard encoder — C. M. Wright (P&L, CH) U.S. Pat. 3786115, December 11, 1973

RCA Limited

Automatic optical bias control for light modulators — A. Watsberg, J. I. Wood (Ld, Montreal) U.S. Pat. 3780296, December 18, 1973

Gas discharge devices and electrode for use therein — R. A. Crane (Ld Montreal) U.S. Pat. 3784926, January 8, 1974

Frequency comparator system — L. R. Avery (Ld, England) U.S. Pat. 3783934, January 1, 1974

Services

Control unit for an antenna rotator — R. Kaysen, F. R. Dimeo (E&S, Deptford) U.S. Pat. 3729942, January 15, 1974

Special

Cathode-ray tube-yoke platform-yoke combination and method of assembling the combination — T. M. Shrade (Special Contract Inv.) U.S. Pat. 3786185, January 15, 1974
EC realigns management operations

A major realignment of management operations within the Electronic Components activity of RCA was announced recently by John B. Fareese, Executive Vice President, Electronic Components.

The major changes in the realignment include the appointment of George W. Duckworth as Division Vice President and General Manager, Industrial Tube Division; addition of new responsibilities to Joseph H. Colgrove, Division Vice President and General Manager, Entertainment Tube Division; and the consolidation of the original equipment manufacturer (OEM) and distributor marketing operations under common marketing groups according to divisional product line.

Other members of Mr. Fareese's staff will continue in their present assignments as follows: Fred M. Bauer, Division Vice President, Finance; Charles E. Burnett, Division Vice President; Joseph H. Colgrove, Division Vice President and General Manager, Entertainment Tube Division; Arnold M. Durham, Manager, and News and Information; Joseph A. Haines, Division Vice President, Distributor Products; Clarence A. Kameen, Division Vice President, Industrial Relations; and Clifford H. Lane, Division Vice President, Technical Planning. Messrs. Bauer, Burnett, Colgrove, Duckworth, Durham, Haines, Kameen and Lane will continue to be headquartered in the Electronic Components executive offices in Harrison, N.J.

The organization of the Entertainment Tube Division was announced by Mr. Colgrove as follows: William G. Hartzell, Division Vice President, Engineering; Robert B. Means, Division Vice President, Marketing and Distribution; William B. Miller, Manager, International Operations; David O. Price, Director, Materials; Herbert Taber, Director, Financial Controls and Operations Planning; Charles W. Thierfelder, Division Vice President, Manufacturing.

The organization of the Industrial Tube Division was announced by Mr. Duckworth as follows: William E. Bradley, Manager, Quality & Reliability Assurance; William E. Cirec, Manager, Industrial Relations; Joseph M. Cleary, Director, Marketing; Joseph A. Galaccione, Manager, Material Systems; Arthur W. Hoek, Manager, Financial Controls and Planning; Victor C. Houk, Manager, Business Planning; C. Price Smith, Director, Industrial Tube Operations Department.

The RCA Electronic Components organization now consists of two divisions: Entertainment Tube Division that produces and markets receiving and television picture tubes, NUMITRON digital display devices, and tube parts and materials; and the Industrial Tube Division that produces and markets a broad variety of industrial tubes such as camera, intensifier, photomultiplier, gas laser, power, microwave, and display - plus tube cavities, solid-state microwave devices, solid-state silicon photodetectors, solid-state infrared emitters and closed circuit video equipment. In addition, a line of RCA batteries, audio magnetic blank tape cassette, reel-to-reel and eight-track tape cartridges, SK replacement semiconductors and kits, and over 40 electronic test instruments are sold through the RCA Distributor Products activity.

RCA Electronic Components operates domestic electronic tube manufacturing facilities in Harrison and Woodbridge, N.J.; Lancaster and Scranton, Pa.; Marion, Ind.; Los Angeles, Calif. and Circleville, Ohio and also manufactures electron tubes and tube parts at all subsidiary operations in Puerto Rico, Canada, Mexico and Brazil; in addition to participating in affiliate manufacturing operations in the United Kingdom, Italy and Taiwan.

Degree Granted

Deepak Chopra of the Power Design activity, Solid State Division, Somerville, N.J., received the MSEE from the University of Massachusetts in September, 1973.

Licensed engineers

When you receive a professional license, send your name, PE number (and state in which registered), RCA division, location, and telephone number to: RCA Engineer, Bldg. 204-2, RCA, Cherry Hill, N.J. As new inputs are received they will be published.

National Broadcasting Company


Advanced Technology Laboratories

R. D. Larabee, ATL, Camden, N.J.; PE-20934, New Jersey
E. Hermann, ATL, Camden, N.J.; PE-20961, New Jersey.

Government systems operations in Camden, N.J., Burlington, Mass consolidated

A new RCA division — Government Communications and Automated Systems Division — was formed recently with headquarters in Camden, N.J., to consolidate activities of the former Aerospace Systems organization.

The new Division will be headed by James M. Osborne as Division Vice President and General Manager and will have plant facilities in Camden and in Burlington, Mass., headquarters of the former Aerospace Systems Division.

The consolidation brings together RCA's Camden business in communications equipment for space and for military land, sea and air forces with Burlington-based businesses in automatic monitoring and control systems, automatic test equipment, electro-optical systems and intelligence processing systems.

In announcing the new organization, Irving K. Kessler, RCA Executive Vice President, Government and Commercial Systems, said the company's businesses in broadcast equipment and mobile communications systems will constitute the newly-named Commercial Communications Systems Division.

This Division will continue to be headed by Andrew F. Inglis, Division Vice President and General Manager, and will have plant facilities in Camden and Meadow Lands, Pa. Mr. Inglis continues to have landlord responsibilities for the Camden plant complex.

Prior to his new assignment, Mr. Osborne was Division Vice President, Government Communications Systems, a post he had held since September 1970. He joined the RCA organization in 1956 as Project Manager of the North Atlantic Tropospheric Scatter Program.

Stanley S. Kolodkin, who had been responsible for the Aerospace Systems Division at the time of the consolidation, will continue as Division Vice President in charge of the new Division's operations at the Burlington plant.
Beverage and Tunnel receive VWOA awards

Dr. Harold H. Beverage and G. William Tunnel recently were honored by the Veteran Wireless Operators Association. Dr. Beverage received the Marconi Medal of Honor and Mr. Tunnel received the Marconi Medal of Achievement.

G. (Bill) Tunnel joined RCA in 1940 in the Test Equipment Maintenance Section, where for several years he participated in the construction, repair and installation of extensive test facilities for items ranging from radio receivers to complex military radar, tv and special-purpose electronic equipment. Supervisory and administrative responsibilities during this period helped prepare him for the post-war years and the opportunity to become active in the rapidly developing tv broadcasting industry. He held responsibilities relating to merchandising, product planning and specialized selling, which later developed into product-line management responsibilities in the RCA Broadcast Equipment Division. He was quite closely associated with the development of tv studio equipment and related items.

In the mid-1950’s, he spent a short period developing a nationwide organization of Manufacturer’s Representatives for the distribution of specialized line of test instruments. This was cut short by an opportunity to join the RCA Service Company, where he became Manager of the Sales and Merchandising Section for the Technical Products Service organization and later Service Manager of the Mid-Eastern Region.

In 1968, Mr. Tunnel was named Vice President and General Manager of RCA Service Division in Canada.

In 1971 he returned to the United States to become Division Vice President, Technical Services — the position he currently holds — with responsibility for the rapidly expanding Data Communications Lease and Service business plus rf communications, entertainment and business services.

Dr. Harold Beverage was Vice President, Research and Development, RCA Communications, Inc., and Director, Radio Research Laboratory, RCA Laboratories at the time of his retirement in 1958.

Dr. Beverage received the BSEE from the University of Maine in 1915 and the PhD in Engineering in 1938. From 1915 to 1916, he was employed by the General Electric Company in Schenectady, N. Y. During the next four years he was a radio laboratory assistant to Dr. E. F. W. Alexander, at General Electric. He joined RCA in 1920.

Dr. Beverage was elected a Vice-President of RCA Communications in December, 1940. Previously he was Chief Research Engineer of RCA Communications since 1929, when the company was formed as a wholly-owned RCA subsidiary. From 1920 to 1929, he served as Research Engineer in charge of Communication Receiver Development for RCA.

Dr. Beverage holds a number of patents in the field of radio. He is also the author of many papers on radio which have been published in technical journals.

During World War II, Dr. Beverage was a Consultant to the Secretary of War on Communications Problems, for which he received a Certificate of Appreciation from the Chief Signal Officer. He was awarded the Morris Liebmann Memorial Prize of the Institute of Radio Engineers in 1923 and the Medal of Honor of the I.R.E. in 1945. He also received the President’s Certificate of Merit in 1948 and the Armstrong Medal from the Radio Club of America in 1938.

In June, 1957, the American Institute of Electrical Engineers (now IEEE) awarded him the Lamme Gold Medal “for his pioneering and outstanding engineering achievements in the conception and application of principles basic to progress in national and world-wide radio communications.” He is a Fellow and past President of the IRE. He is also a Fellow of the Radio Club of America, the American Institute of Electrical Engineers, and the Association for the Advancement of Science. In 1955, he was elected an Eminent Member of Eta Kappa Nu.

Awards

Missile and Surface Radar Division

David D. Freedman, of the Signal Processing Activity has been selected for the 1973 Annual Technical Excellence Award.

In winning this top honor, Mr. Freedman was cited for his outstanding technical achievements in ultra-high-speed circuitry and logic development applicable to radar systems, culminating in the conception, development, and delivery of a fully digital system performing pulse-space-timing of 2.3 in. or 1/2560 microsecond. This High-Speed Synchronizer (HISYN) system has been delivered to the Air Force’s Floyd Site radar where it has been operating successfully for several months.

James J. Campbell received a technical excellence award for the third quarter of 1973 for major contributions to the successful demonstration of the AEGIS missile mid-course command data link and guidance policy.

Samuel D. Gross received a technical excellence award for the second quarter of 1973 for outstanding contribution to the successful testing of the AFAR system with the Land Based Test Site.

Thomas H. Mehling received a second quarter 1973 technical excellence award for his role in lead radar system design in the successful integration and formal system tests of the first Engineering Development Model of the AEGIS system at the Land Based Test Site.

Dr. Ralph J. Pschunder received a technical excellence award for the second quarter of 1973 for outstanding technical achievement in his development of full state-of-the-art capability in the area of structural and dynamic analysis.

RCA Laboratories

Fifty-two scientists on the RCA Laboratories staff have received Achievement Awards for outstanding contributions to electronics research and engineering during 1973.

Recipients of the awards and brief descriptions of the work for which they were honored are:

Aline L. Akselrad for fundamental contributions to the fields of uniaxial anisotropy and magneto-optics in garnets.

Francis J. Campbell for development of a user-oriented computer program that takes into account the effects of space charge in calculating electron trajectories.

David E. Carlson for contributions to the understanding of the relationship between ionic motion and the modification of the properties of glasses.
Missile and Surface Radar Division

Merrill W. Buckley, Jr., accepted the Chapter of the Year Award on behalf of the Engineering Management Group of the Philadelphia Section. Mr. Buckley was Chairman of the Group (1972-1973).

MSRD authors' reception

Over a hundred engineers were honored recently at MSRD's annual authors' reception. This reception, hosted by Dudley Cotter, Chief Engineer, was the seventeenth in a series that started in 1966 to honor engineers who have presented or published papers, received patents, or earned doctoral degrees.

Mr. Cotter addressed the group briefly: "You can't put a price tag on the value of this kind of professional activity, either for an individual or for the reputation of the company. I won't even try. But I do want to say that it's important—very important—to all of us.

"In many ways it is the only yardstick the technical community has to measure the level of our capabilities. It is certainly true that a large part of our reputation around the country is a direct result of the papers you've given or published as well as the patents you've been granted."

Research and Engineering

Frank W. Widmann, Staff Engineer, Engineering is a member of the IEEE Manpower Committee. Recently, the 1973 Manpower Committee released a report entitled Career Outlook in Engineering (described separately).

IEEE releases manpower report

The Manpower Committee of the Institute of Electrical and Electronics Engineers has released its 1973 report, Career Outlook in Engineering (Regions 1 through 6, USA). The report states that generally "career opportunities in electronics will continue to grow [however] ... it will be necessary to remain flexible and ever ready to transfer to another area as activity changes in focus and grows in turn to meet market demands."

The report urges engineers to keep abreast of new developments so that it is possible to adjust with the market. It also warns of over-specializing your education during your college career. The Committee members observed that "employment problems are seldom encountered where the individual has maintained a regular program of study to update and broaden his background."

This book is planned as the first of an annual series of reports that will provide accurate employment information about the electrical/electronic engineering field. It is intended to aid in career planning for those currently employed in the field as well as aiding students at the high school and college levels. The 225-page report is divided into four sections: the industry picture; the manpower picture; careers in engineering; and the engineering challenges. The cost for IEEE members is $10; non-members $15. [Write to Martin Gitten, IEEE, 345 East 47th Street, New York, N.Y., 10017.]

The Committee studied a Carnegie Commission Report which called on professional schools and professional societies to provide careful studies of manpower supply and demand. "We [the Manpower Committee] heartily agree, and sincerely hope these activities of IEEE, of which this Report is the beginning effort, will fulfill just such a responsibility."

Professional activities

Aerospace Systems Division

Norm Laschever is General Chairman of the 1974 NEREM Convention to be held in Boston in late October. He was Chairman of the Special Events Committee in 1972 and Vice Chairman of the General Committee in 1973. NEREM is one of the largest IEEE shows in the country, and last year attracted almost 10,000 attendees.

Newton Teixeira is a past chairman of the Boston Chapter of the IEEE Engineering Management Society. He is a secretary program of the Boston Area Committee on the Technology Forecasting and Assessment project being conducted by IEEE as its contribution to the definition of national technical goals (an objective of the National Science Foundation). He has recently been elected a member of the Ad Com (National Administrative Committee) of the Engineering Management Society of IEEE. He is also a member of the American Association for the Advancement of Science and the American Management Association.

Pat Toscano has been associated with IEEE for twenty-five years and is currently a senior member. His most recent IEEE activity is as Colloquium organizer for "Software for the Engineer": a six-session tutorial review of software engineering currently in progress on Thursday evenings. He has served as author, moderator, and session chairman at several Automated Support Systems Conferences from 1966 to 1973 in St. Louis, Philadelphia, and Dallas. He is a member of the National Society of Professional Engineers, the Massachusetts Society of Professional Engineers and is a registered Professional Engineer in Massachusetts.

Advanced Technology Laboratories

John E. Friedman, Leader, Engineering Communications, has been elected to a
three-year term on the Administrative Committee of the IEEE Group on Professional Communication.

Murlan S. Corrington recently received the IEEE Philadelphia Section Award "for contributions to mathematical analysis of circuits and systems, and for service to the IEEE.

Staff announcements

Manufacturing Services and Materials

Howard W. Johnson, Staff Vice President, Reliability and Quality has announced the appointment of David W. Reynolds as Director, Reliability and Quality.

Engineering

Howard Rosenthal, Staff Vice President, Engineering, Research and Engineering has appointed Henry Ball, Staff Engineer, Engineering.

Laboratories

William M. Webster, Vice President, Laboratories has announced the appointment of Gerald B. Herzog as Staff Vice President, Technology Centers.

Nathan L. Gordon, Director, Systems Research Laboratory has announced the appointment of Alfred H. Tegar as Head, Advanced Systems Research.

Consumer Electronics

Roy H. Pollack, Division Vice President and General Manager, Color and Black and White Television Division, has announced the organization of the Color and Black & White Television Division as follows: Harry Anderson, Director, Manufacturing Operations; David E. Daly, Division Vice President, Advanced Product Planning; Loren R. Kirkwood, Director, Color TV Engineering and Strategic Planning; Robert J. Lewis, Chief Engineer, Black and White TV Engineering; William S. Lowry, Division Vice President, Product Management—Color TV; Tucker P. Madawick, Division Vice President, Industrial Design; Richard Mentzinger, Director, Product Management—Black & White TV; and James R. Smith, Director, Quality and Reliability.

Solid State Division

Ben A. Jacoby, Division Vice President, Solid State Power, has announced the appointment of Donald Watson as Manager, Power Operations Control.

Joseph W. Karoly, Division Vice President, Solid State—Europe has announced the appointment of Herbert B. Shannon as Manager, Power Operations.

Herbert B. Shannon, Manager, Power Operations—Europe, has announced the appointment of Barry B. J. Charles as Manager, Discrete Power Transistors.

Harry Weisberg, Director, MOS IC Product Operations has announced the appointment of William E. Wagner as Manager, Market Planning—MOS IC Products.

William E. Wagner, Manager, Market Planning—MOS IC Products has announced the appointment of George A. Riley as Manager, Product Planning.

Donald R. Carley, Manager, MOS IC Engineering, announced the MOS IC Technology organization as follows: Terry G. Athanas as Manager, MOS IC Technology and Acting Leader, SOS Product Technology; Martin A. Blumenfeld, Leader, MOS IC Process Technology; Walter F. Lawrence, Project Leader, Assembly Technology and Package Development; and Bernard B. Levin, Leader, MOS IC Model Shop & Manufacturing Support.

RCA Global Communications, Inc.

Philip Schneider, Executive Vice President, Leased Facilities and Engineering has announced the appointment of A. William Brook as Director, Satellite System Development. Mr. Brook has also announced the organization of Satellite System Development as follows: Armand L. Dill Pare, Manager, Launch Vehicles and Mission Requirements; Don L. Lundgren, Manager, Project Administration; Charles V. Lundstedt, Manager, Earth Stations and Data Sub-Systems; Lloyd A. Ottenberg, Manager, Telecommunications Systems; and Jack L. Ray, Manager, Applications Engineering.

James M. Walsh, Director, Engineering, Leased Facilities and Engineering, has announced the appointment of John Christopher as Manager, Spacecraft Engineering.

James H. Muller, Manager, Network Management and Joint User Service, Operations, has announced the appointment of John Golden as Manager, Joint User Service Technical Operations.

Edwin W. Peterson, Executive Vice President, Finance and Business Development has announced the organization of Finance and Business Development as follows: Thomas J. Brady, Vice President and Controller; Frederick J. Sager, Vice President and Treasurer; Armand W. Aymong, Director, Financial Planning and Capital Budgets; Dennis Elliott, Director, Business Development; Alexander MacGregor III, Director, Profit Planning; and Lewis Hoffman, Manager, Management Information Systems.

Morris Pincus, Vice President, Regulatory and Affairs and Operating Arrangements, Operations, has announced the appointment of Allan E. Schwamberger as Director, Regulatory Affairs and Tariffs.

Palm Beach Division

William J. Hannan, Chief Engineer, Engineering, has announced the Engineering organization as follows: Stanley A. Basara, Manager, Product Development Engineering; Charles M. Breder, Manager, Engineering Services; Lester J. Limbaugh, Manager, Data Interconnect Systems Engineering; Rudolf J. Spoelstra, Manager, Product Technology Engineering.

Lester J. Limbaugh, Manager, Data Interconnect Systems Engineering, has announced the Data Interconnect Systems Engineering organization as follows: James H. Bragdon, as Manager, Systems Planning; Emory C. Manager, Data Interconnect Design Engineering; Thomas B. Rhodes as Manager, Data Interconnect Programming; and Larry E. Thompson as Manager, Data Interconnect Design Engineering.

Stanley E. Basara, Manager, Product Development Engineering, has announced the Product Development Engineering organization as follows: Orville A. Gwinn as Manager, Special Products Engineering and Bruce E. Smith as Manager, Communication Systems Engineering.

Rudolf J. Spoelstra, Manager, Product Technology Engineering has announced the Product Technology Engineering organization as follows: Edwin M. Fulcher, Manager, Circuits and Design Automation Engineering; Karl H. Hoffman, Manager, Circuits and Design Automation Engineering; Karl H. Hoffman, Manager, Power Supply and Packaging Engineering and Robert G. Saenz, Manager, Components Engineering.

Charles M. Breder, Manager, Engineering Services, has announced the Engineering Services organization as follows: Bruce B. Ballard, Manager, Administrative Services; William A. Ostmann, Manager, Technical Services; Richard J. Simone, as Manager, Technical Publications and Louis B. Wolf as Manager, Experimental Shop, Engineering Services.

James Vollmer, General Manager, Palm Beach Division has announced the appointment of Stanley E. Basara as Manager, Ford Program.

Promotions

Astro-Electronics Division

W. W. Metzger from Mgr., Mechanical Analysis to Mgr., Mechanical Engineering (W. Manger, Hightstown)

B. Stewart, from Mgr., AE Systems Engineering to Mgr., Preliminary Design (W. Manger, Hightstown)
Electromagnetic and Aviation Systems Division


R. C. Hayes from Mgr., Design Engineering to Mgr., EW Engineering (F. C. Corey, Van Nuys)

H. Hite from Mgr., Electrical Develop. & Des. to Mgr., Electrical Engineering (F. C. Corey, Van Nuys)


R. L. Reed from Ldr., Logics to Mgr., Government Programs Support (F. C. Corey, Van Nuys)


Electronic Components

G. T. Rose from Engr., Manufacturing, Television Picture Tube to Mgr., Production Engineering (N. Meena, Marion)

Global Communications

A. W. Brook from Mgr., Satellite System Development to Dir., Satellite System Development (P. Schneider, 60 Broad Street, New York)

J. Christopher Mgr., Spacecraft Engineering (J. M. Walsh, 60 Broad Street, New York)

A. C. Cowan from RCA Educational System to Mgr., Iran Operations (A. A. Avanesians, Iran Operations)

Dr. A. Dil Pare Mgr., Launch Vehicle & Mission Reports (A. W. Brook, 60 Broad Street, New York)

J. Gleitman from Staff Engr., to Mgr., Computer Systems Engineering (A. A. Avanesians, Computer Engineering Projects)

Charles V. Lundstedt Mgr., Earth Station and Data Subsystems (A. W. Brook, 60 Broad Street, New York)

F. Micara from Tech. Illustrator to Ldr., Drafting (J. M. Walsh, 60 Broad Street, New York)

M. Rosenthal from Ldr., Documentation to Mgr., Engineering Services (J. Walsh, 60 Broad Street, New York)

R. Ruben from Engineer to Ldr., (D. Mandato, Technical Control and Special Equipment Engineering)

Record Division

J. E. Grein from Engineering Assoc. to Mgr., Mfg. Services Coordination (J. R. Mason, Indpls.)

J. Gunter from Sr. Membr. to Ldr., Engineering Staff (J. E. Lang, Indpls.)

J. Heller from Sr. Membr. to Ldr., Engineering Staff (J. E. Lang, Indpls.)

M. McNeely from Member to Ldr., Engineering Staff (J. R. Mason, Indpls.)

R. Nyman from Member to Ldr., Engineering Staff (J. R. Mason, Indpls.)

Eric M. Leyton — 1916-1974

Eric M. Leyton, 58, Staff Engineer, Engineering, for RCA at the David Sarnoff Research Center in Princeton, N.J., died February 26 in Geneva, Switzerland, where he was attending a meeting of the International Radio Consultative Committee (CCIR) at the request of the State Department.

Mr. Leyton, who was associated with RCA for more than twenty years, made major contributions to the development of television transmitters and tape recording. He was actively concerned with color television systems for many years and frequently represented both the United States and RCA at international conferences on television standards.

A native of London, England, he received his professional EE degree from Faraday House College of London University in 1938. He joined RCA Laboratories in 1953 after six years at the Research Laboratories of Electrical & Musical Industries, Hayes, near London.

Mr. Leyton was awarded twenty two patents. He was a Fellow of the Institute of Electrical and Electronics Engineers and a member of the British Institution of Electrical Engineers.

Chester M. Sinnett — 1900-1974

Mack Sinnett died early this year. He was widely known for his distinguished professional engineering work at RCA and for his encouragement, guidance, and training of young engineers. He was the first Editor of TREND and a pioneer in establishing the RCA Engineer magazine.

He retired from RCA as Director, Product Engineering Professional Development, in January 1965 after more than 40 years of service with RCA and one of its predecessor companies (Westinghouse).

Mr. Sinnett went to work for the Engineering Department of the Westinghouse Electric and Manufacturing Company in Pittsburgh in 1924 after he received the BSEE from the University of Maine. When RCA Victor was organized in 1929, he was moved to Camden as Manager, Phonograph Design and Development. Later, he was named Manager, Loudspeaker and Phonograph Section. In connection with his work in the audio and acoustic field, Mr. Sinnett was granted 30 patents.

In 1945, he was appointed Manager, Home Instruments Advanced Development Section. He held this post until his appointment in 1959 as Director of Product Engineering Professional Development. In this position, his responsibilities included programs concerned with continuing engineering education.

A Fellow of the IEEE, Mr. Sinnett served as Chairman of the Philadelphia Section. He was a Registered Professional Engineer in New Jersey and a member of the National Society of Professional Engineers and Tau Beta Pi.

A recognized authority on creativity in engineering, Mr. Sinnett lectured widely on the subject within and outside of RCA. At the time of his death, he was Head of Instrumentation at the Pepper Laboratory of the University of Pennsylvania. Mr. Sinnett was also a Charter Member of the Electronic Club, an RCA organization in Camden.
Three RCA men elected IEEE Fellows

The three RCA men cited herein have been honored for their professional achievements by being elected Fellows of the Institute of Electrical and Electronics Engineers. This recognition is extended each year by the IEEE to those who have made outstanding contributions to the field of electronics.

Editor's note: In addition to these three recipients, one past RCA employee attained the grade of Fellow from IEEE: Arthur D. Beard, formerly with Computer Systems Division, Cherry Hill, NJ, "for contributions to the development of magnetic storage devices and solid-state techniques for data processing"

Dr. Hershenov received the BS in Physics in 1950, the MS in Mathematics in 1952, and the PhD in Electrical Engineering in 1959, all from the University of Michigan. He joined RCA Laboratories in Princeton, N.J., in 1960. He became Head of the Microwave Integrated Circuits Group in 1968 and four years later was appointed Director of Research of RCA Research Laboratories, Inc., in Tokyo.

He has been active in the IEEE Professional Groups on Magnetics and on Microwave Theory and Techniques. From 1964 to 1966, Dr. Hershenov was a coadjutant professor in the Mathematics Department of Rutgers University.

Dr. Hershenov is a member of Phi Kappa Phi, the American Physical Society, Sigma Xi, and is listed in "American Men of Science."

Dr. Kressel received the BA in Physics from Yeshiva University in 1955, the MS in Applied Physics from Harvard University in 1956, and the MBA in 1959 and PhD in Metallurgy in 1965, both from the University of Pennsylvania.

Dr. Kressel joined the RCA Semiconductor Division in Somerville, N.J., in 1959 and transferred to RCA Laboratories in Princeton in 1966. He was appointed Head of the Semiconductor Devices Research Group in 1969.

He is the author or coauthor of more than 80 technical articles and has contributed chapters to several published books dealing with semiconductors.

He is a member of Sigma Xi and the American Physical Society as well as the IEEE.

Mr. Luther received the BSEE from M.I.T. in 1950. Subsequently, he joined RCA and was engaged in circuit design of color television cameras and monitors for the Communications Systems Division.

During the sixties, Mr. Luther's principal accomplishments included: development of the TR-22 recorder, the industry's first all-solid-state video recorder; a complete line of products based on the TR-22; the TR-70 highband video recorder; the development of long-life video headwheels; and development and introduction of the TCR-100 automatic video cartridge recorder system.

Mr. Luther is presently Chief Engineer, Broadcast Systems, Camden, N.J. He holds 29 U.S. Patents, is the author of numerous papers, and is a member of Eta Kappa Nu.

Mr. Luther is presently Chief Engineer, Broadcast Systems, Camden, N.J. He holds 29 U.S. Patents, is the author of numerous papers, and is a member of Eta Kappa Nu.
Editorial Representatives

The Editorial Representative in your group is the one you should contact in scheduling technical papers and announcements of your professional activities.

Government and Commercial Systems

Aerospace Systems Division
K. E. PALM Engineering, Burlington, Mass.

Electromagnetic and Aviation Systems Division
C. S. METCHETTE Engineering, Van Nuys, Calif.
J. MCDONOUGH Aviation Engineering, Van Nuys, Calif.

Astro-Electronics Division
L. M. SEIDEMANN Engineering, Princeton, N.J.
S. WEISBERGER Advanced Development and Research, Princeton, N.J.

Missile & Surface Radar Division
D. R. HIGGS Engineering, Moorestown, N.J.

Government Engineering
M. G. PIETZ Advanced Technology Laboratories, Camden, N.J.
J. E. FRIEDMAN Advanced Technology Laboratories, Camden, N.J.
J. L. KRAGER Central Engineering, Camden, N.J.

Government Plans and Systems Development
E. J. GODDELL Eng. Information and Communications, Camden, N.J.

Communications Systems Division

Broadcast Systems
R. N. HURST Broadcast Systems Advanced Development, Camden, N.J.
A. C. BILLIE Broadcast Engineering, Meadowlands, Pa.
A. LIGUORI Engineering, Camden, N.J.

Mobile Communications Systems

Government Communications Systems
H. J. RAULIS Chairman, Editorial Board, Harrison, N.J.
S. M. SHERMAN Solid State Technology Center, Somerville, N.J.

Palm Beach Division
B. BALLARD Palm Beach Gardens, Fla.

Research and Engineering Laboratories
C. W. SALL*
Research, Princeton, N.J.
J. H. KAUFMAN Solid State Technology Center, Somerville, N.J.
M. R. SHARON Solid State Technology Center, Somerville, N.J.

Electronic Components
C. A. MEYER Chairman, Editorial Board, Harrison, N.J.

Entertainment Tube Division
J. ROYF Receiving Tube Operations, Woodbridge, N.J.
J. R. LIPSETTE Television Picture Tube Operations, Marion, Ind.

Industrial Tube Division

Consumer and Solid State Electronics

Solid State Division
E. M. McELWEE* Chairman, Editorial Board, Somerville, N.J.
I. J. SCHOEN Solid State Division, Somerville, N.J.
J. DIMARIO Solid State Division, Mountaintop, Pa.
S. SILBERSTEIN Power Transistor, Somerville, N.J.
H. A. UHL Integrated Circuits, Somerville, N.J.
J. O. YOUNG Solid State Division, Findlay, Ohio

Consumer Electronics
C. W. HOYT* Chairman, Editorial Board, Indianapolis, Ind.
R. BUST Engineering, Indianapolis, Ind.
R. C. GRAHAM Audio Products Engineering, Indianapolis, Ind.
F. HOLT Advanced Development, Indianapolis, Ind.
E. E. JANSEN Black and White TV Engineering, Indianapolis, Ind.
J. STARK Color TV Engineering, Indianapolis, Ind.

Services
M. G. GANDER Consumer Services Administration, Cherry Hill, N.J.
W. W. COOK Consumer Services Field Operations, Cherry Hill, N.J.
R. M. DONBROSKY Technical Support, Cherry Hill, N.J.
R. J. CgooHill Missile Test Project, Cape Kennedy, Fla.

Parts and Accessories
C. C. REAUCH Product Development Engineering, Deptford, N.J.

RCA Global Communications, Inc.
W. S. LEIB* RCA Global Communications, Inc., New York, N.Y.
G. P. ROBERTS RCA Alaska Communications, Inc., Anchorage, Alaska

National Broadcasting Company, Inc.
W. A. HOWARD Staff Eng., Technical Development, New York, N.Y.

RCA Records
M. L. WHITEHURST Record Eng., Indianapolis, Ind.

Corporate Development
C. A. PASSAVANT International Planning, New York, N.Y.
W. A. CHISHOLM Research & Eng., Montreal, Canada

RCA Ltd
M. S. WINTERS Patent Plans and Services, Princeton, N.J.

*Technical Publications Administrators (asterisked * above) are responsible for review and approval of papers and presentations